

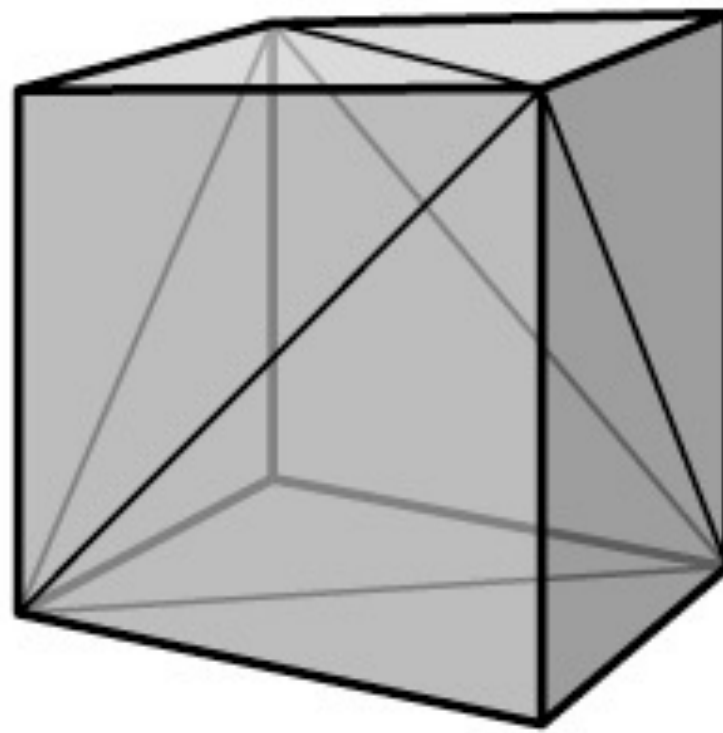
**Lecture 8:**

# **Mesh Representations & Geometry Processing**

---

**Computer Graphics and Imaging  
UC Berkeley CS184/284A**

# A Small Triangle Mesh



**8 vertices, 12 triangles**

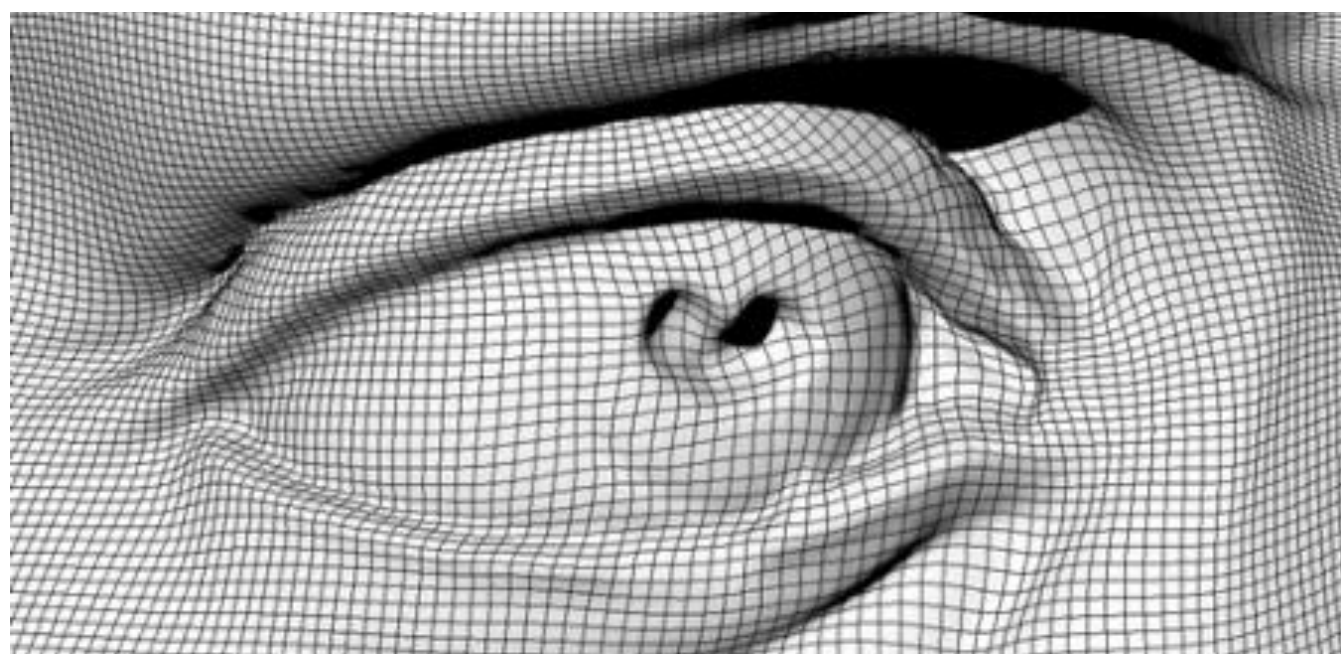
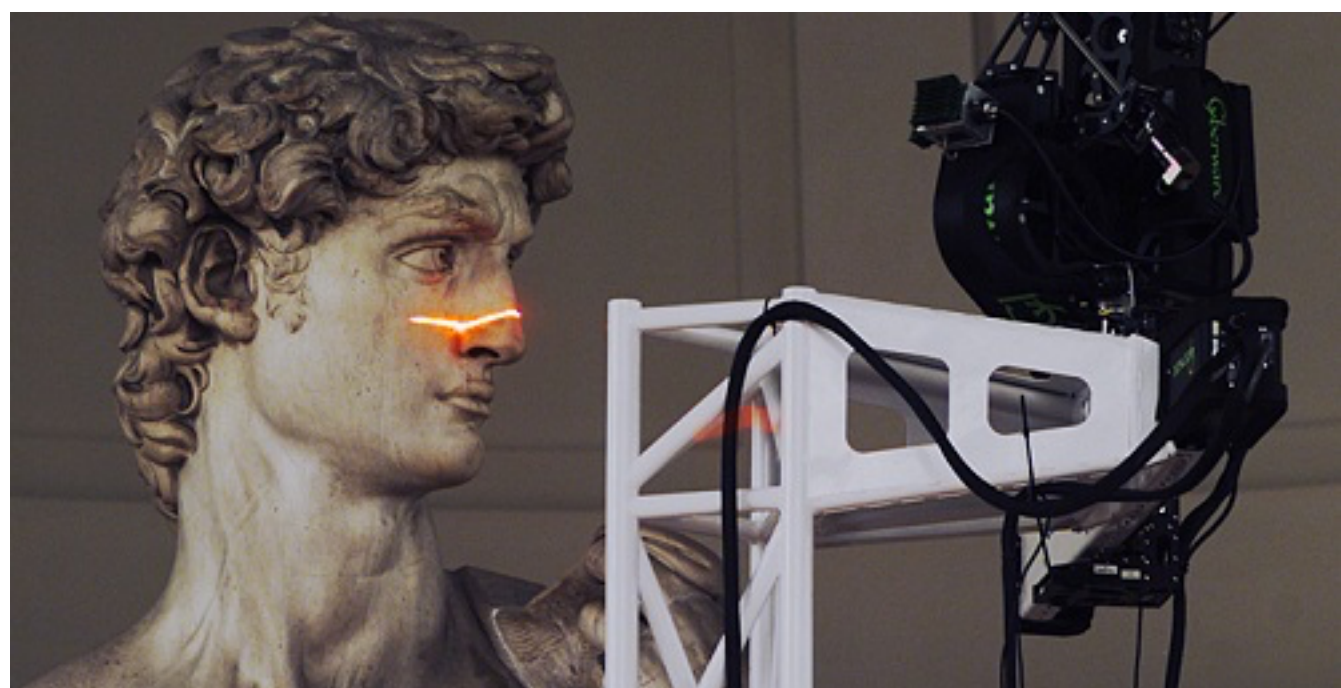
# A Large Triangle Mesh

David

Digital Michelangelo Project

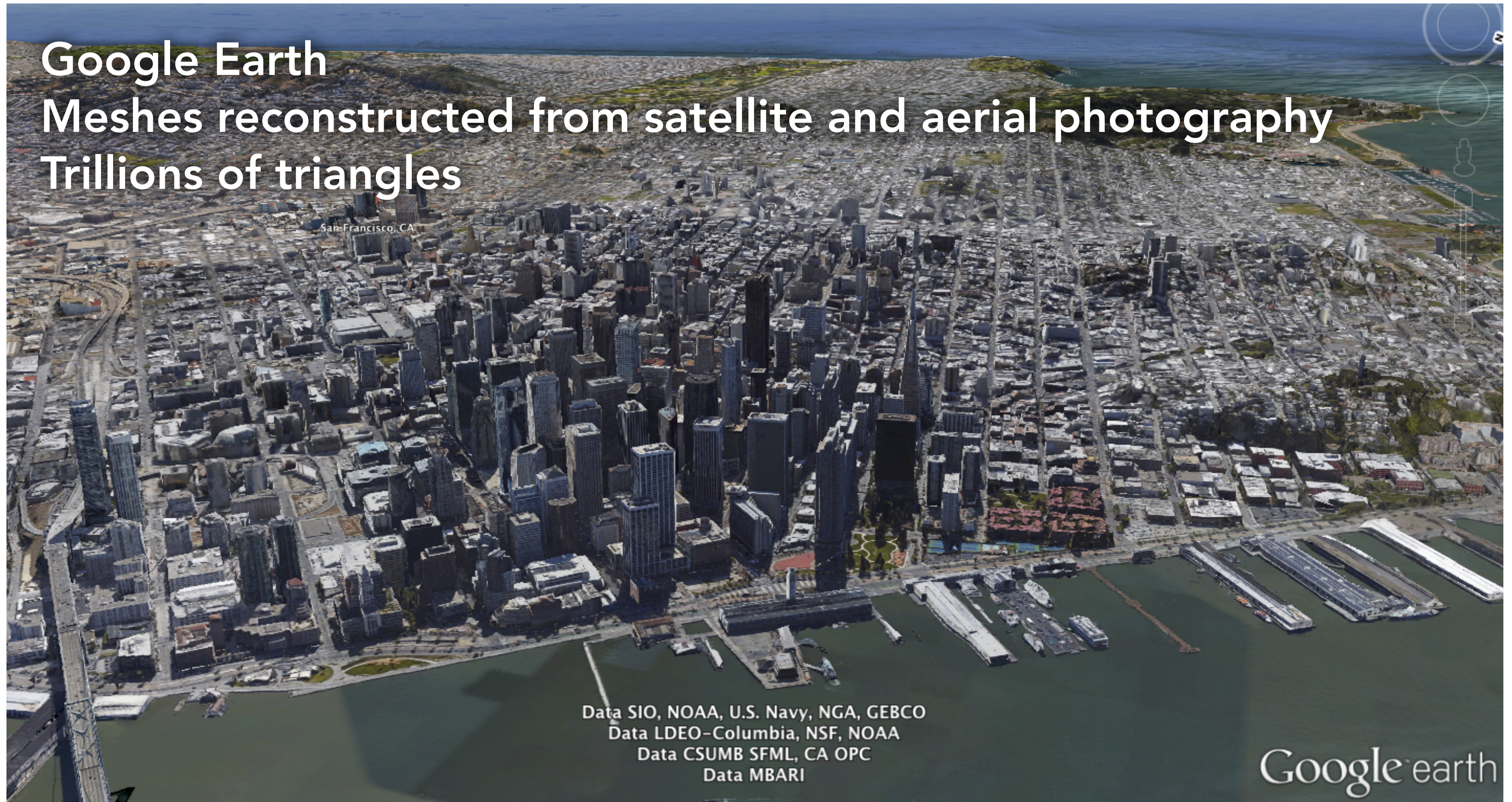
28,184,526 vertices

56,230,343 triangles



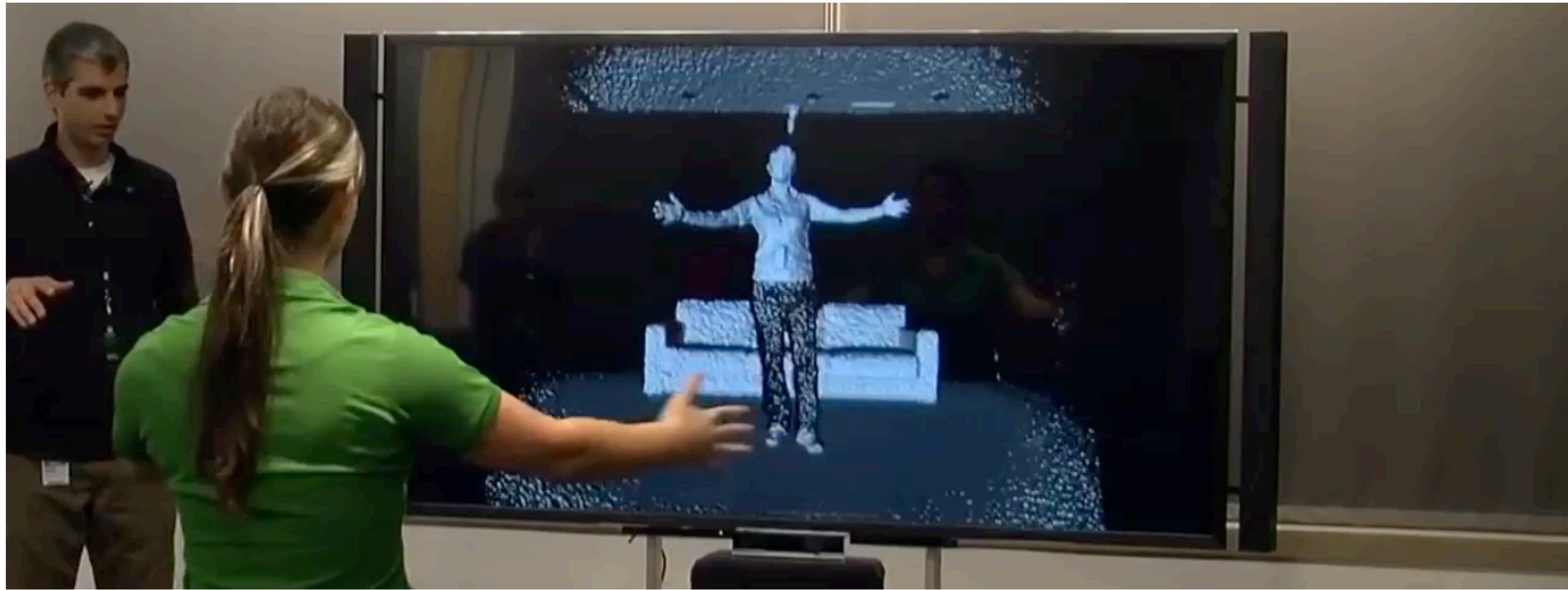


# A Very Large Triangle Mesh

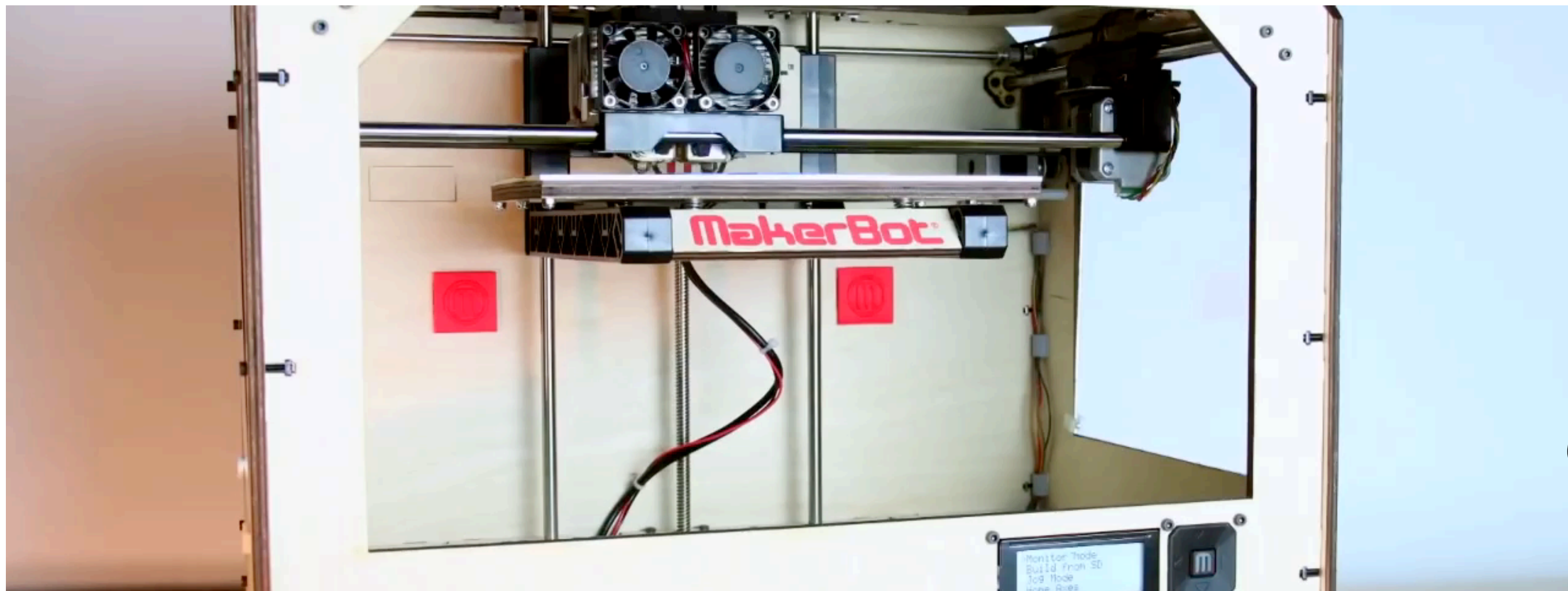




# Digital Geometry Processing



3D Scanning



3D Printing



# Geometry Processing Pipeline



**Scan**



**Process**



**Print**

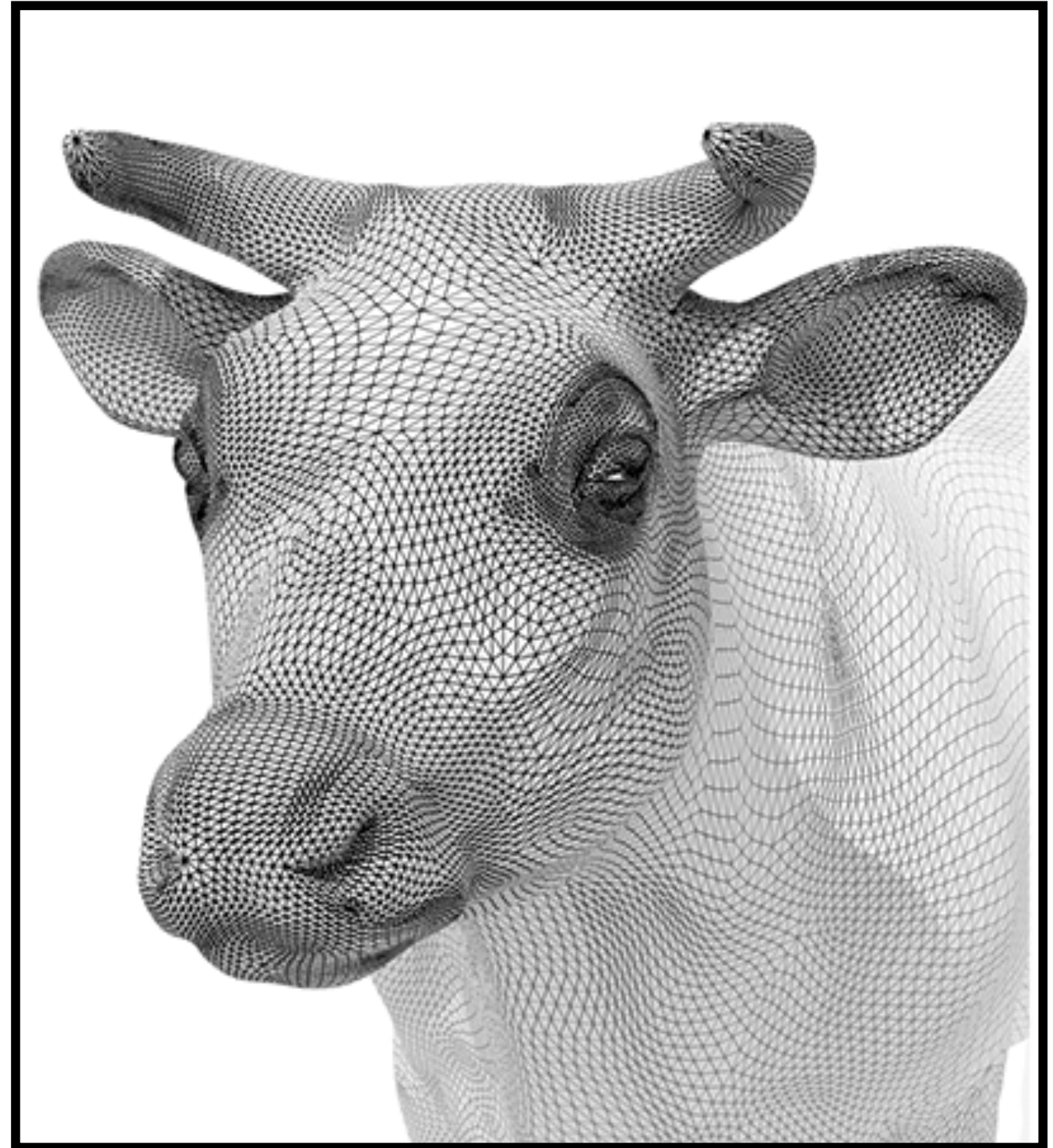
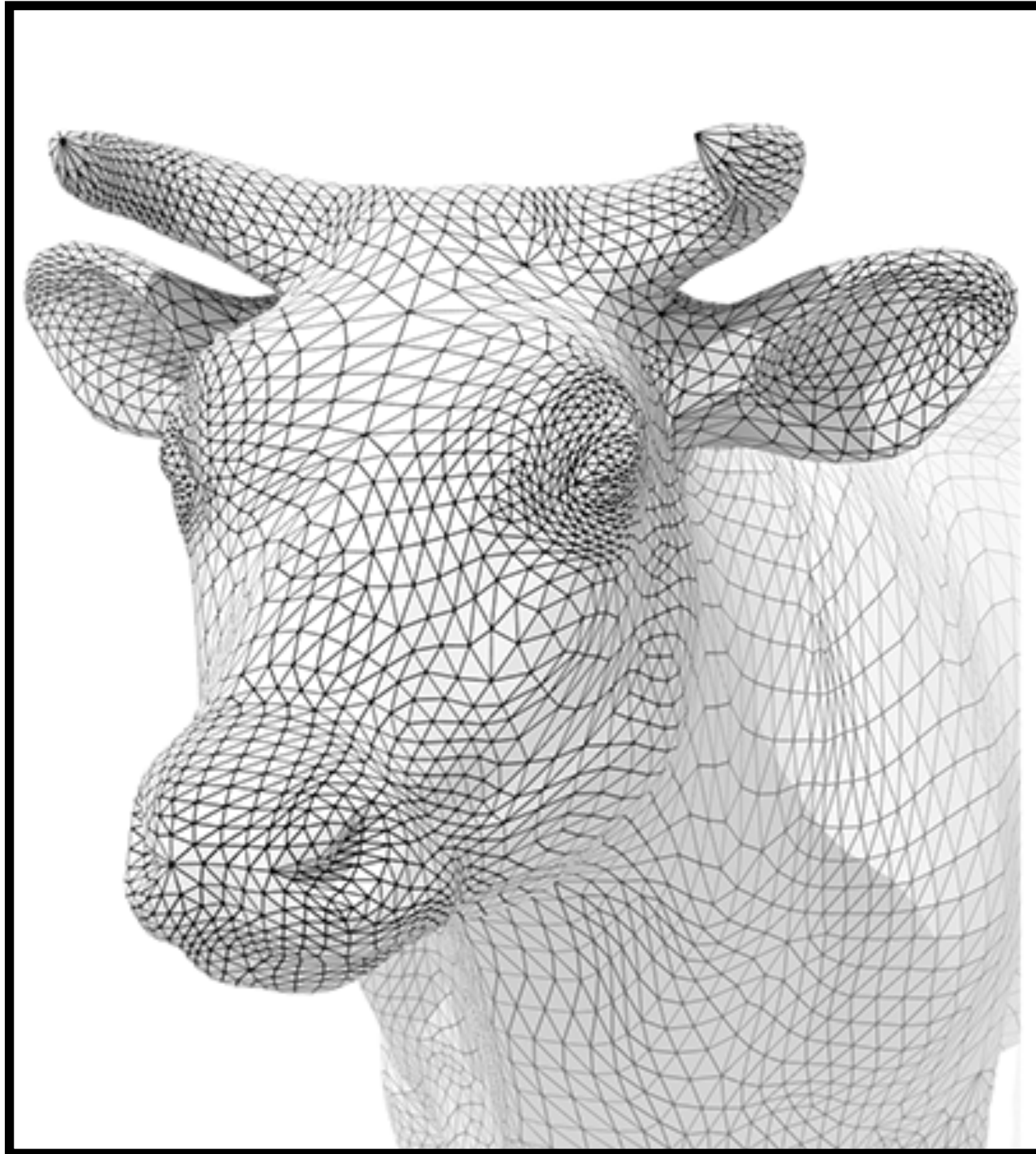


# **Geometry Processing**

## **Tasks: 3 Examples**



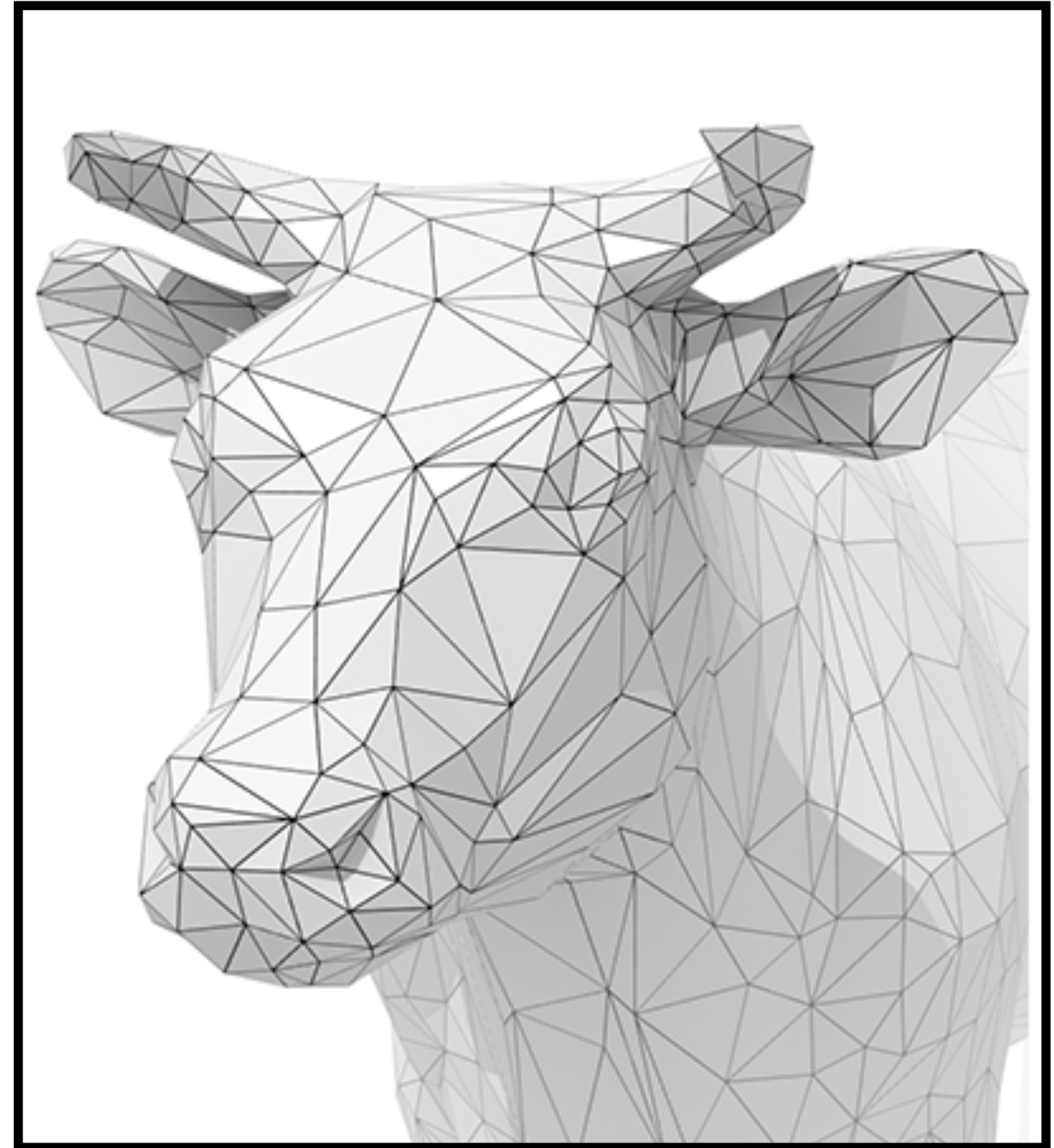
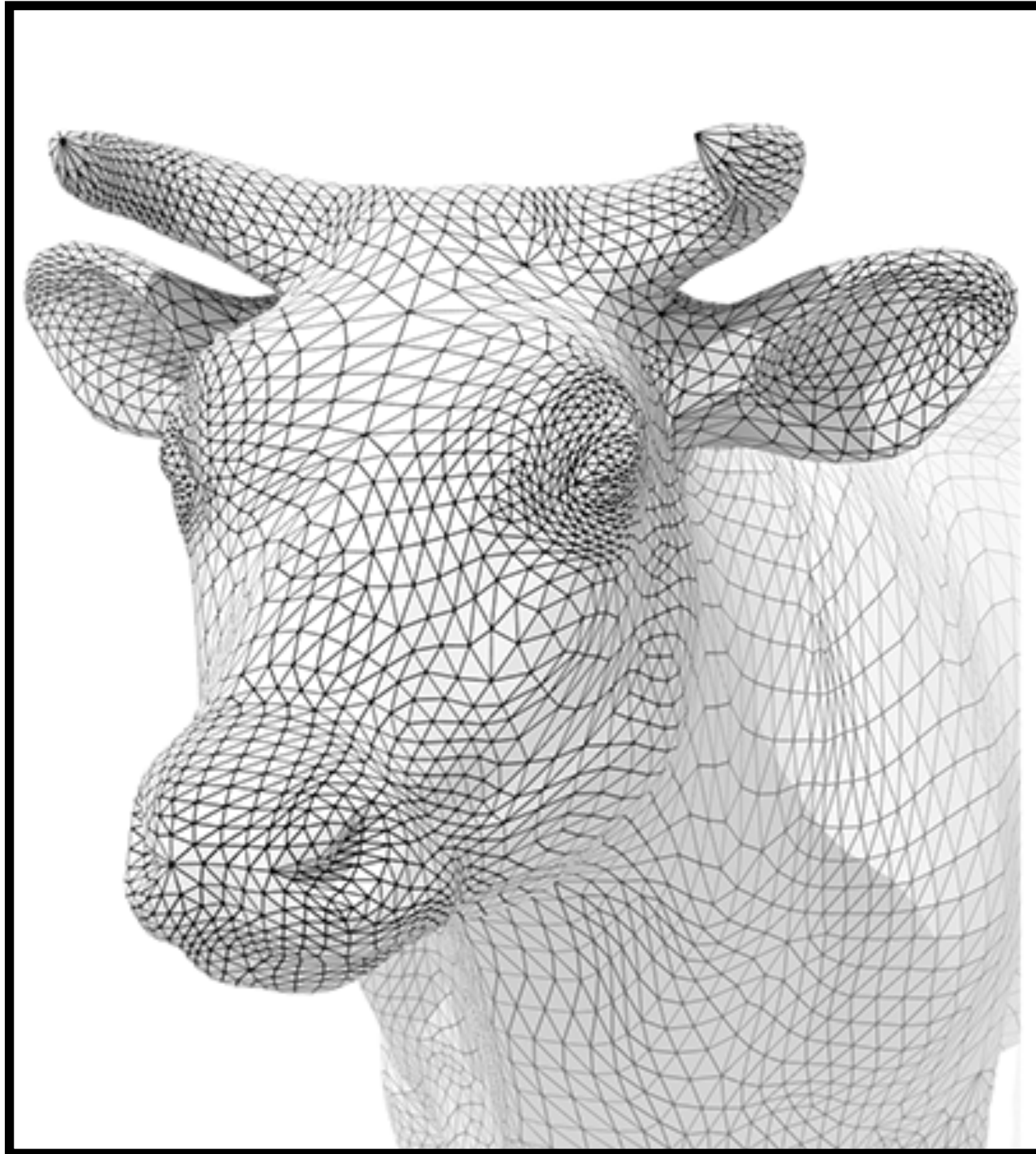
# Mesh Upsampling – Subdivision



Increase resolution via interpolation



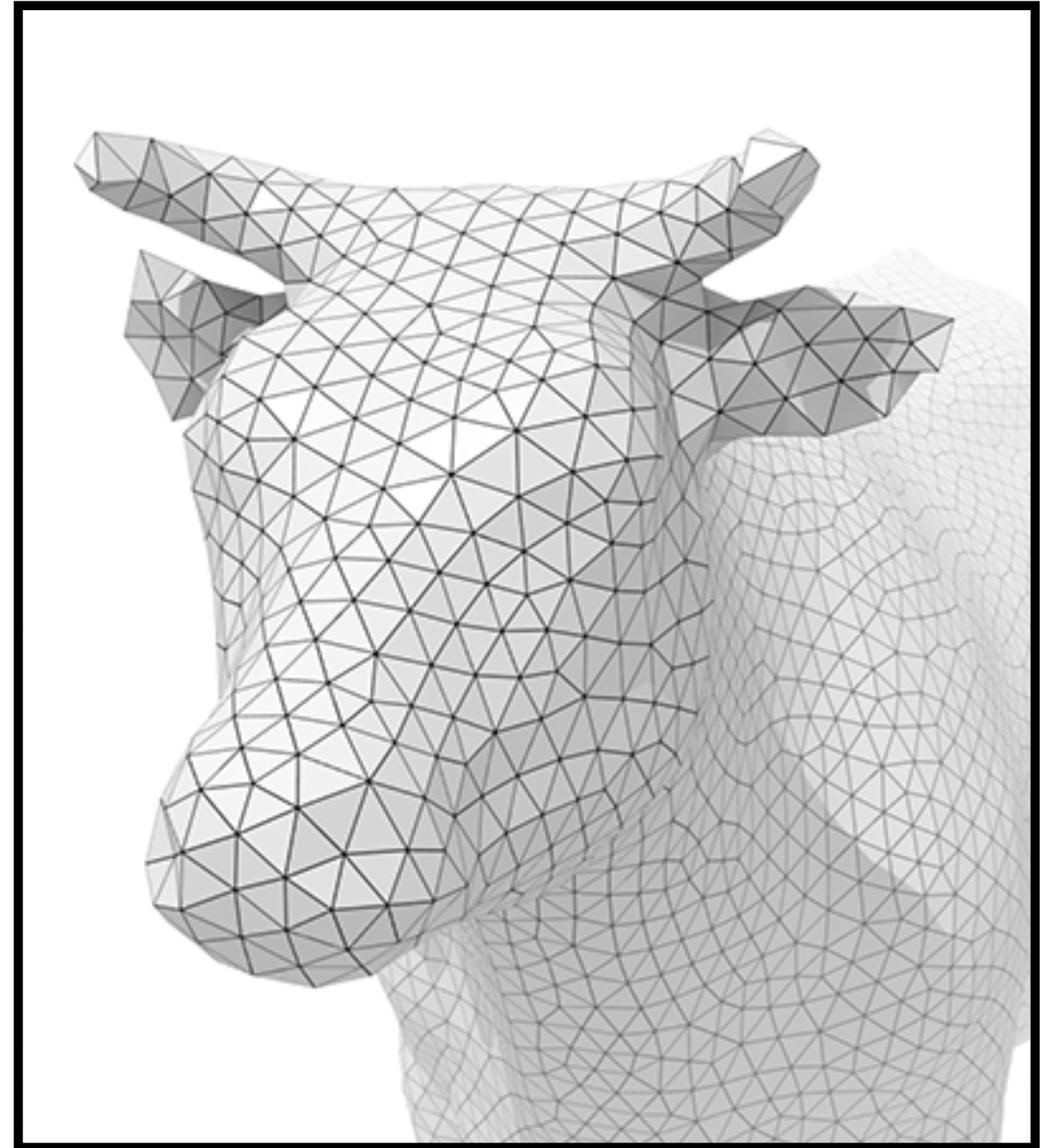
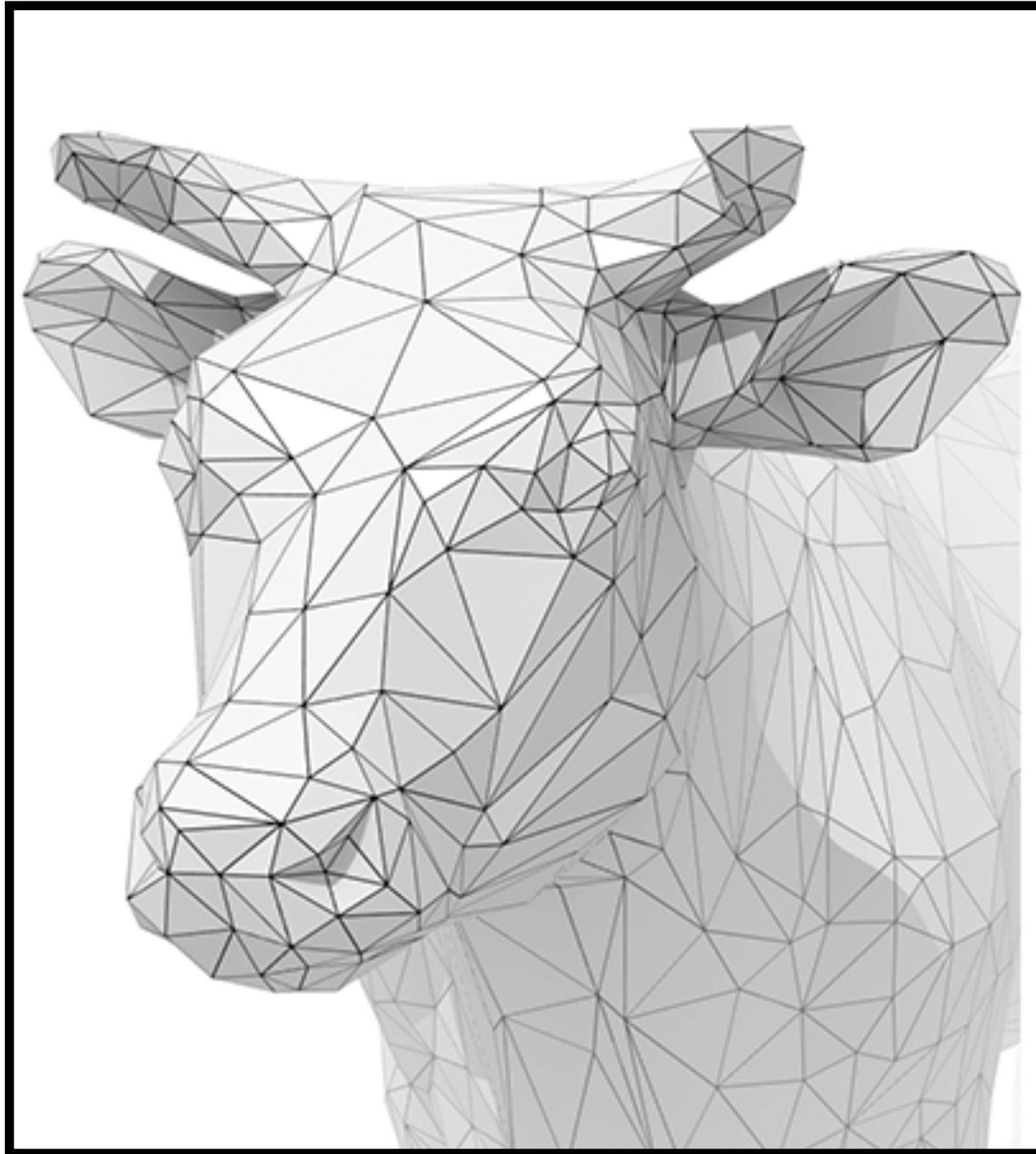
# Mesh Downsampling – Simplification



**Decrease resolution; try to preserve shape/appearance**



# Mesh Regularization



**Modify sample distribution to improve quality**



# **This Lecture**

**Study how to represent meshes (data structures)**

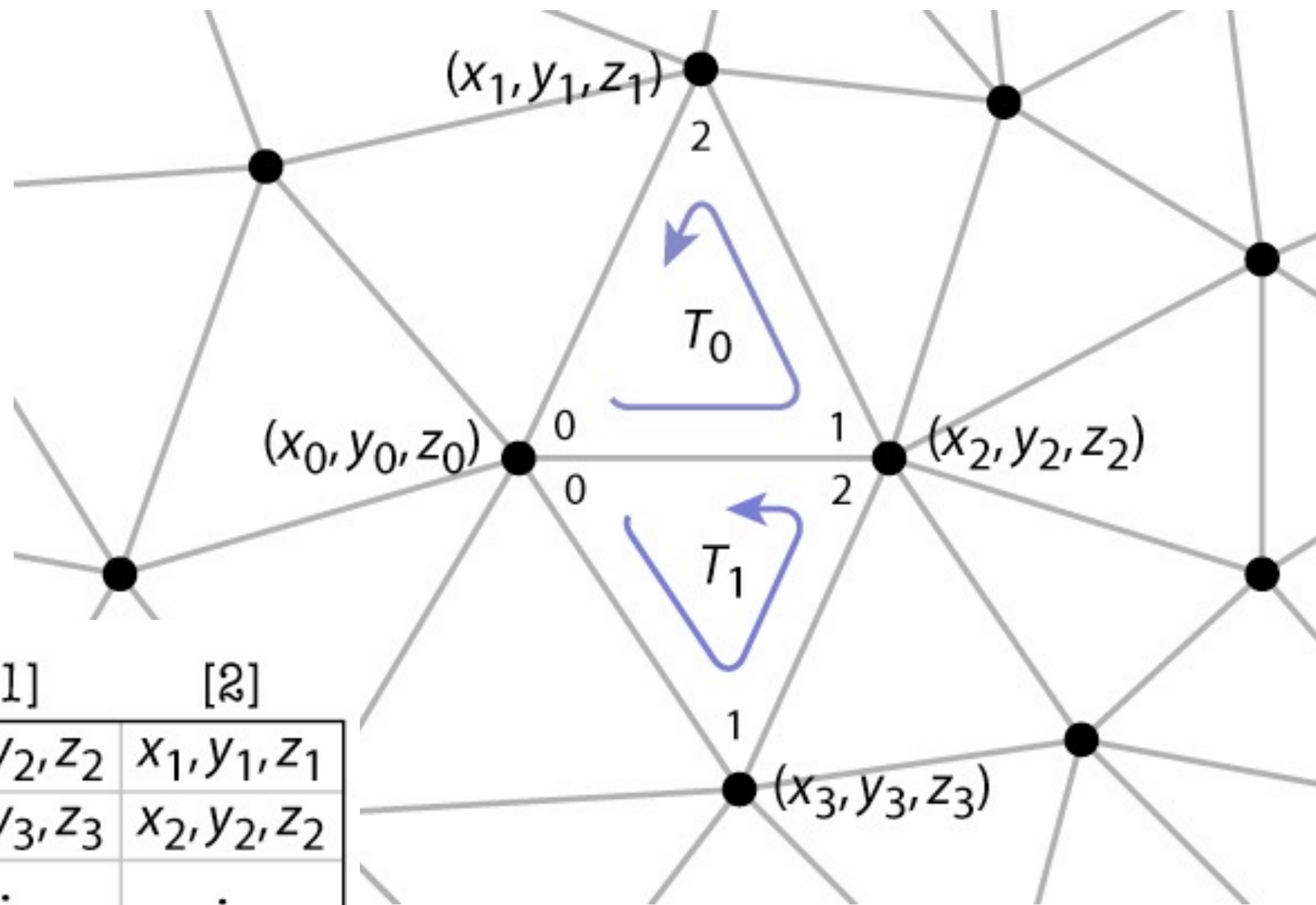
**Study how to process meshes (geometry processing)**



# **Mesh Representations**



# List of Triangles



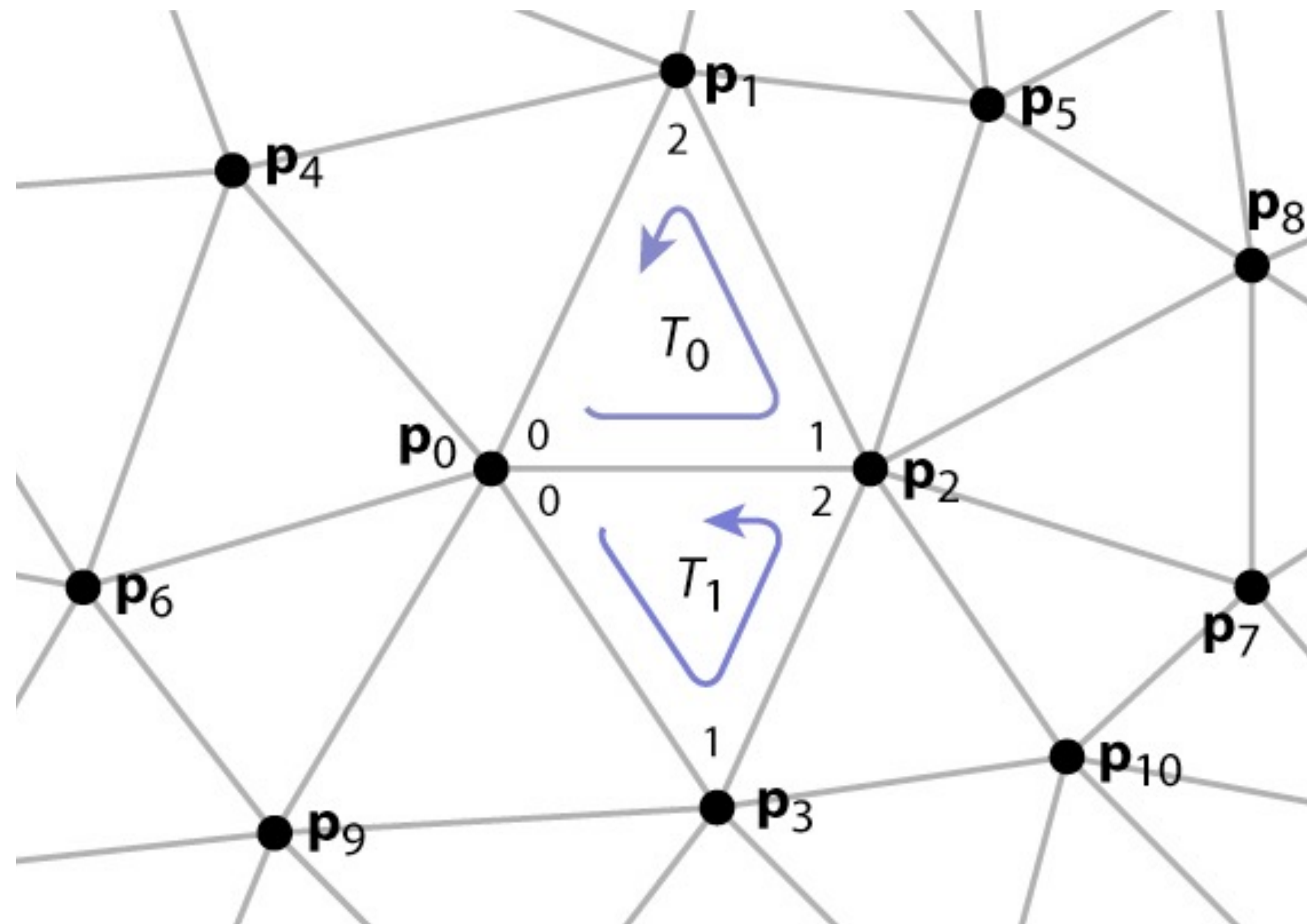
	[0]	[1]	[2]
tris[0]	$x_0, y_0, z_0$	$x_2, y_2, z_2$	$x_1, y_1, z_1$
tris[1]	$x_0, y_0, z_0$	$x_3, y_3, z_3$	$x_2, y_2, z_2$
	$\vdots$	$\vdots$	$\vdots$



# Lists of Points / Indexed Triangle

verts[0]	$x_0, y_0, z_0$
verts[1]	$x_1, y_1, z_1$
	$x_2, y_2, z_2$
	$x_3, y_3, z_3$
	$\vdots$

tInd[0]	0, 2, 1
tInd[1]	0, 3, 2
	$\vdots$



# Comparison

## Triangles

- + Simple
- Redundant information

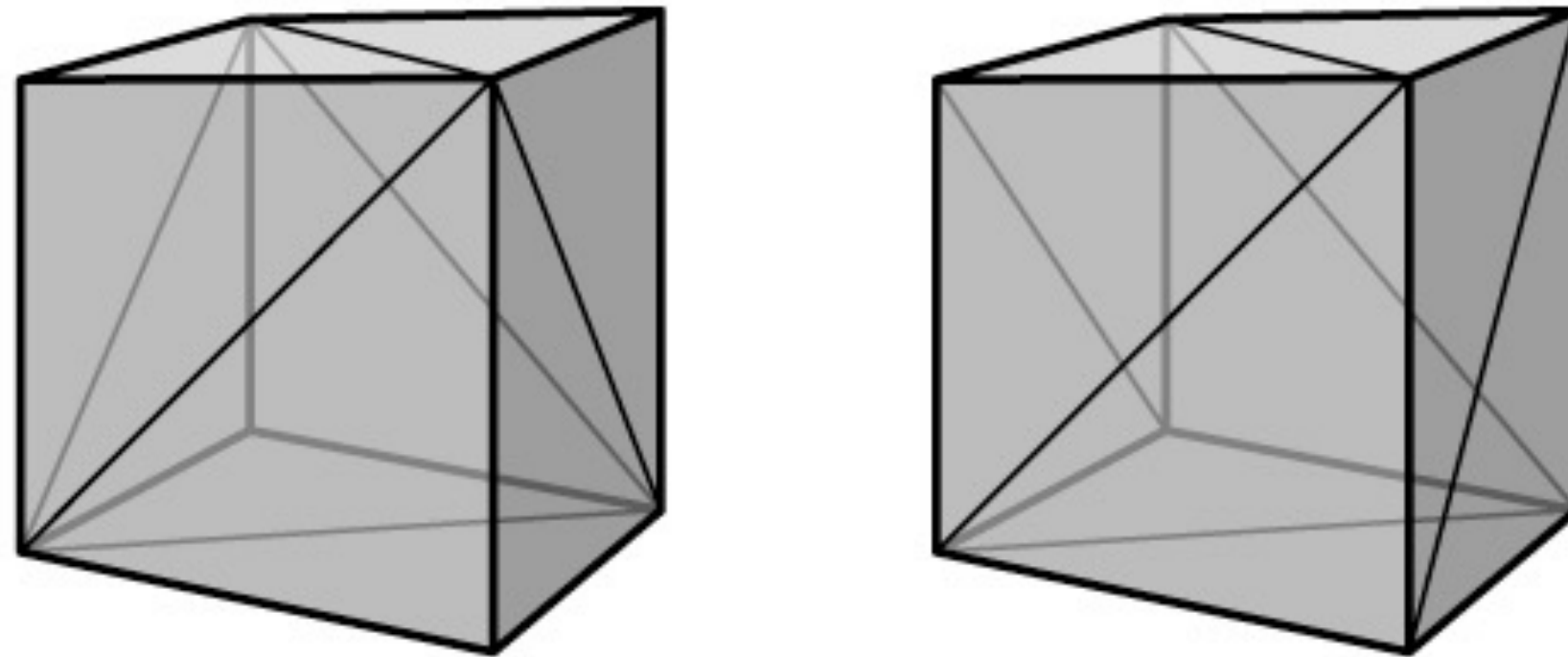
## Points + Triangles

- + Sharing vertices reduces memory usage
- + Ensure integrity of the mesh (moving a vertex causes that vertex in all the polygons to move)

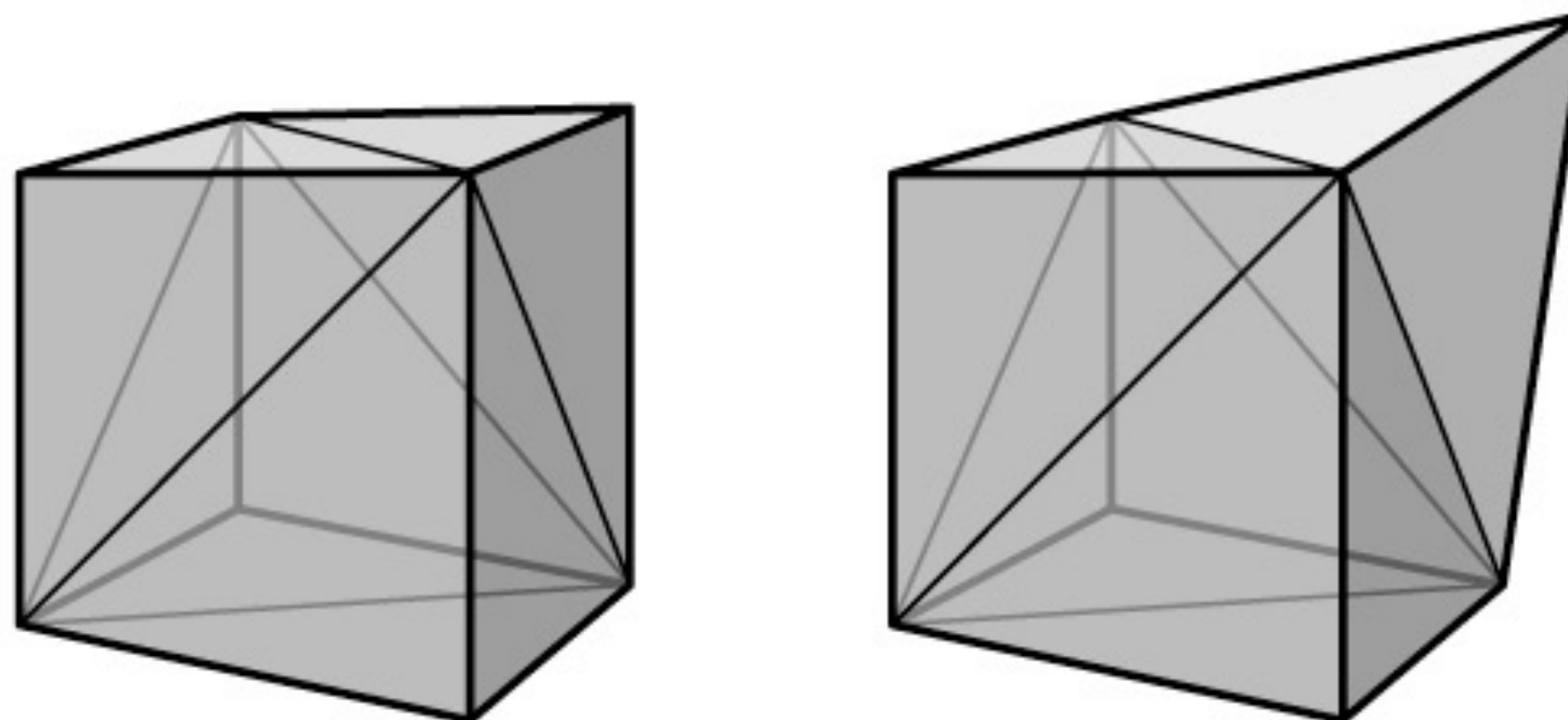


# Topology vs Geometry

Same geometry, different mesh topology



Same mesh topology, different geometry



# Topological Mesh Information

## Applications:

- **Constant time access to neighbors**  
e.g. surface normal calculation, subdivision
- **Editing the geometry**  
e.g. adding/removing vertices, faces, edges, etc.

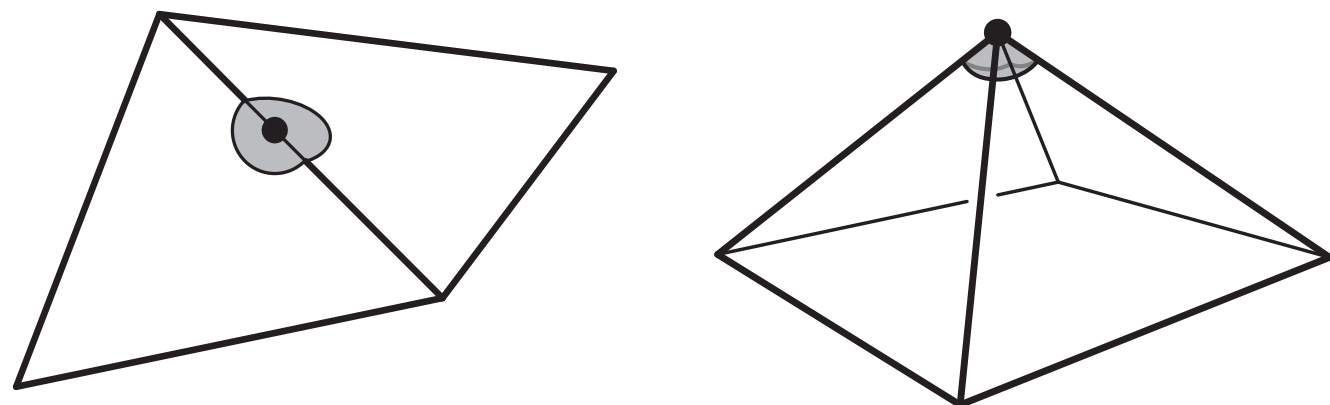
**Solution: Topological data structures**



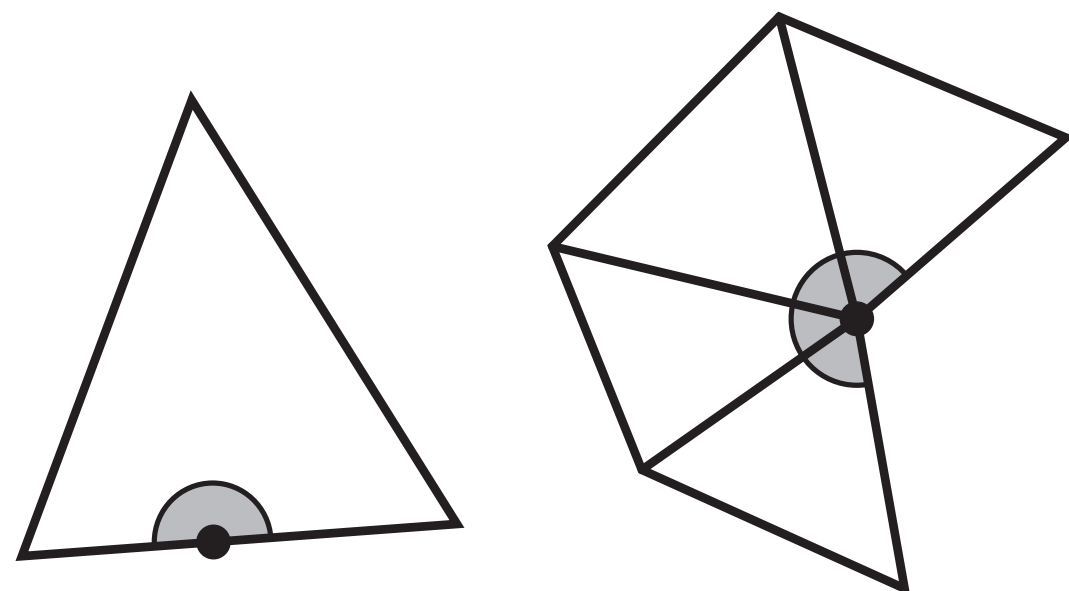
# Topological Validity: Manifold

Definition: a 2D manifold is a surface that when cut with a small sphere always yields a disk.

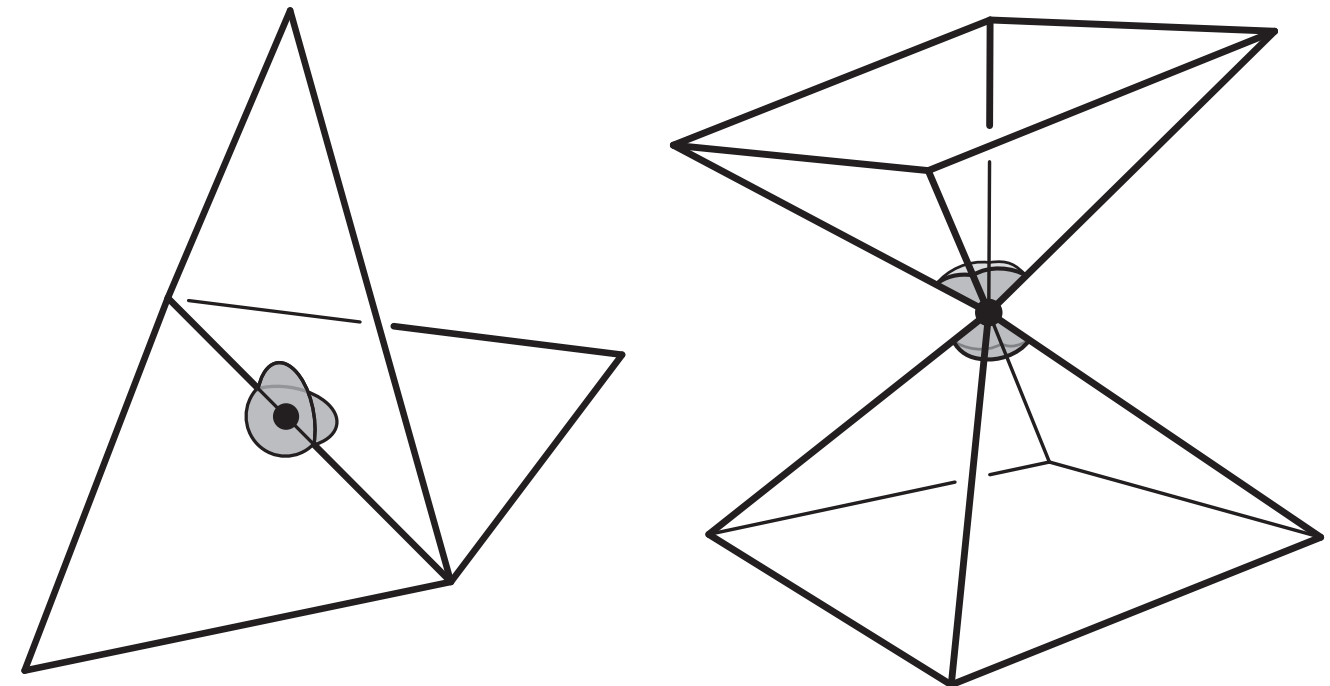
Manifold



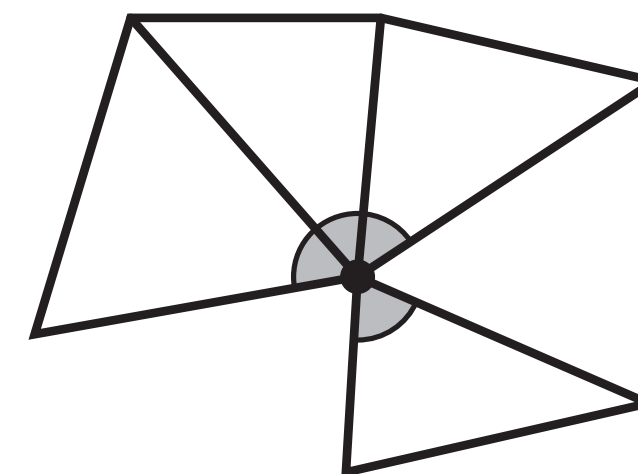
With border



Not manifold



With border



# Topological Validity: Manifold

**Definition:** a 2D manifold is a surface that when cut with a small sphere always yields a disk.

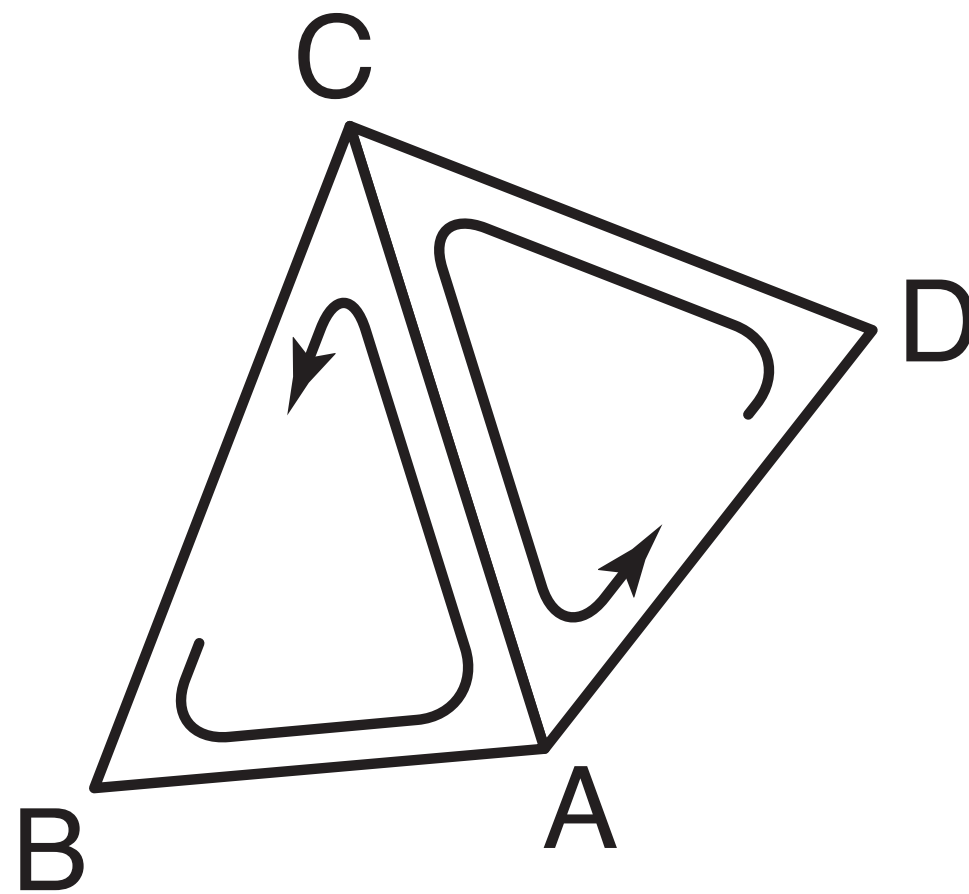
**If a mesh is manifold we can rely on these useful properties:**

- An edge connects exactly two faces
- An edge connects exactly two vertices
- A face consists of a ring of edges and vertices
- A vertex consists of a ring of edges and faces
- Euler's polyhedron formula holds:  $\#f - \#e + \#v = 2$   
(for a surface topologically equivalent to a sphere)  
(Check for a cube:  $6 - 12 + 8 = 2$ )

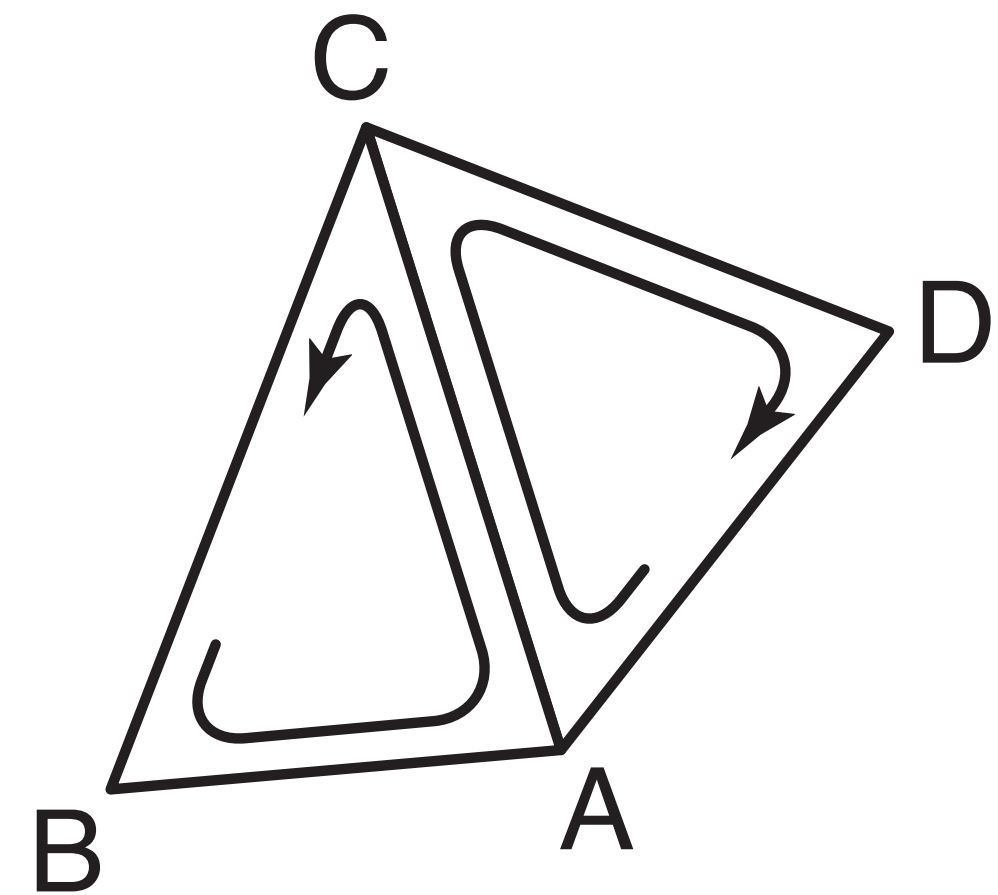


# Topological Validity: Orientation Consistency

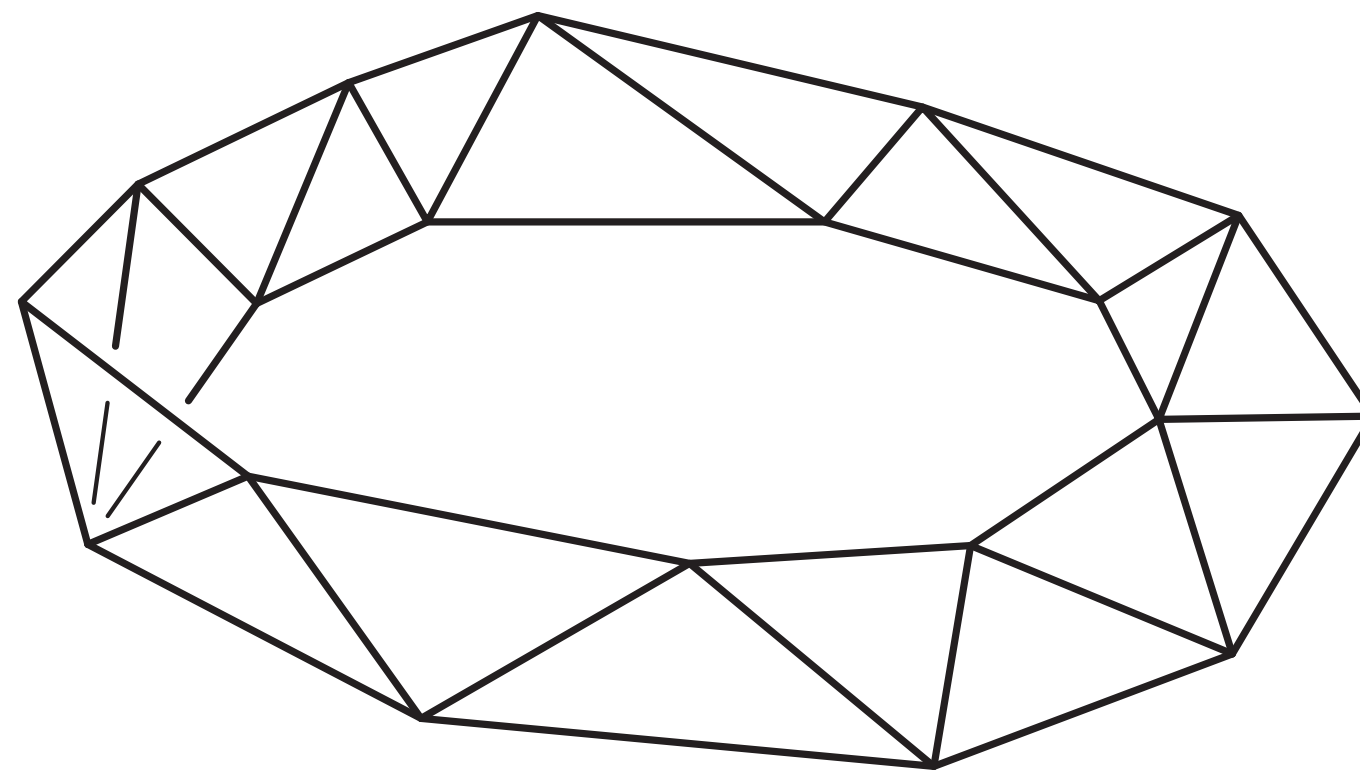
Both facing front



Inconsistent orientations



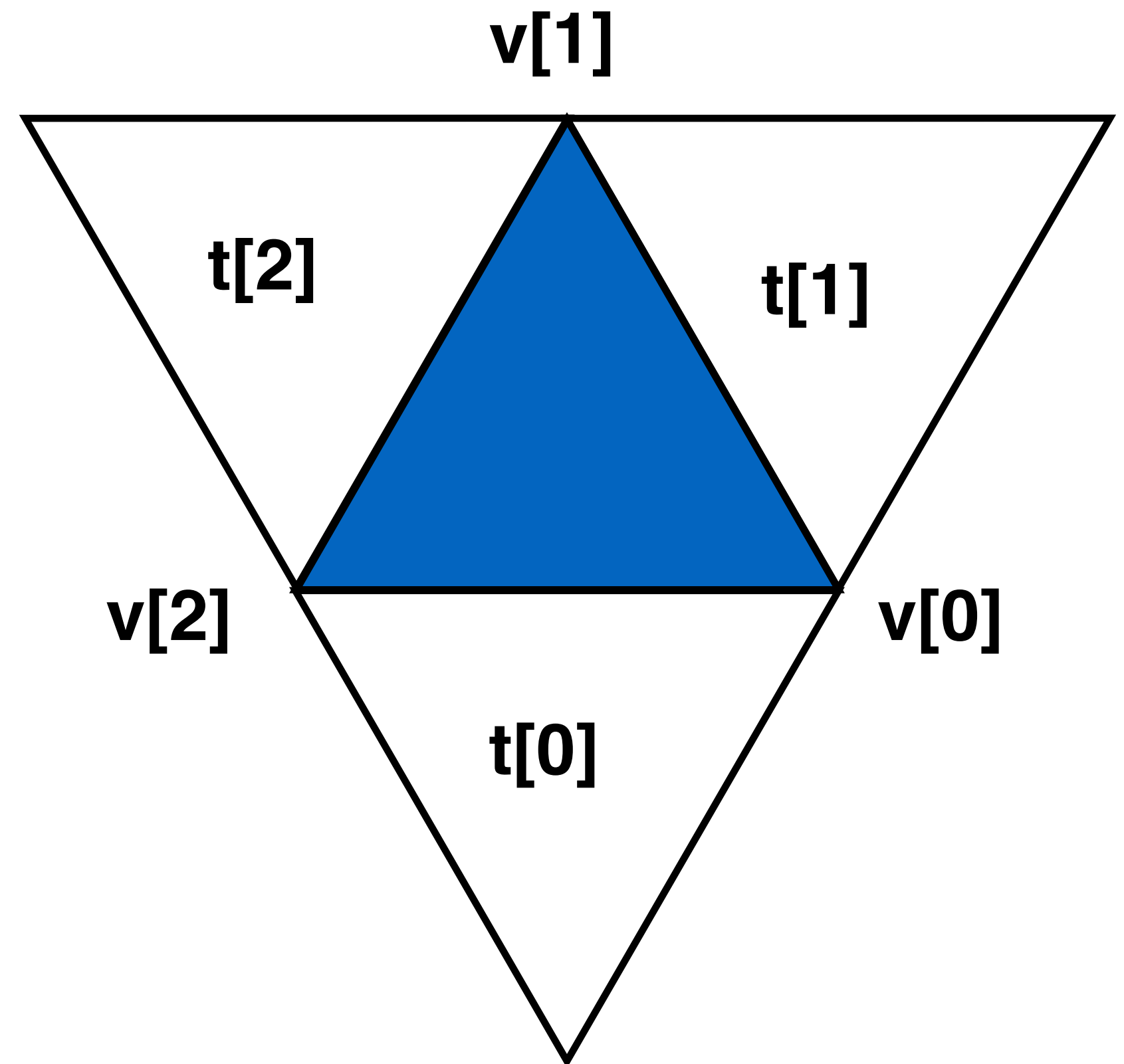
Non-orientable



# Triangle-Neighbor Data Structure

```
struct Tri {  
    Vert    * v[3];  
    Tri    * t[3];  
}
```

```
struct Vert {  
    Point   pt;  
    Tri    *t;  
}
```

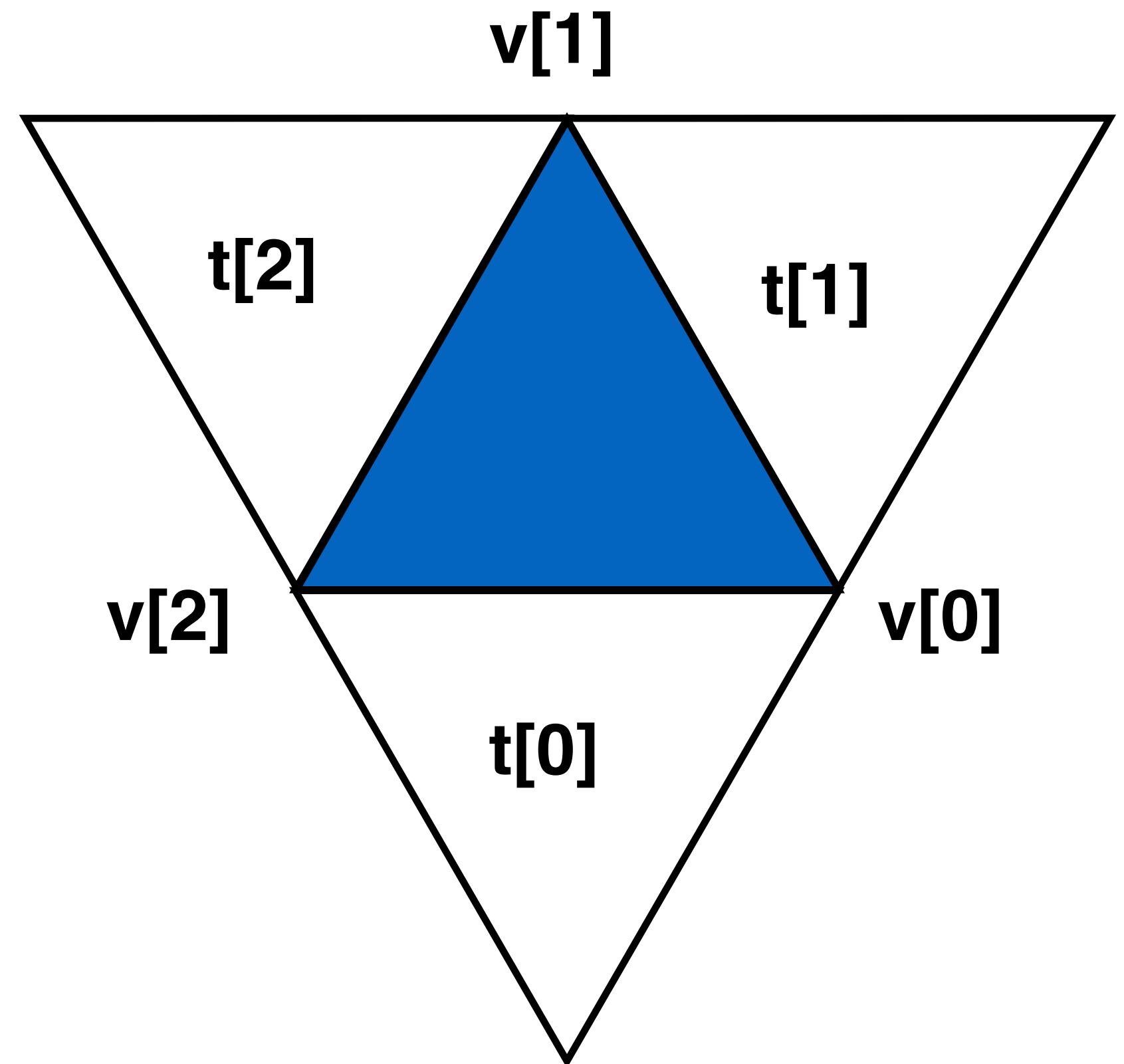




# Triangle-Neighbor – Mesh Traversal

Find next triangle counter-clockwise around vertex  $v$  from triangle  $t$

```
Tri *tccwvt(Vert *v, Tri *t)
{
    if (v == t->v[0])
        return t[0];
    if (v == t->v[1])
        return t[1];
    if (v == t->v[2])
        return t[2];
}
```

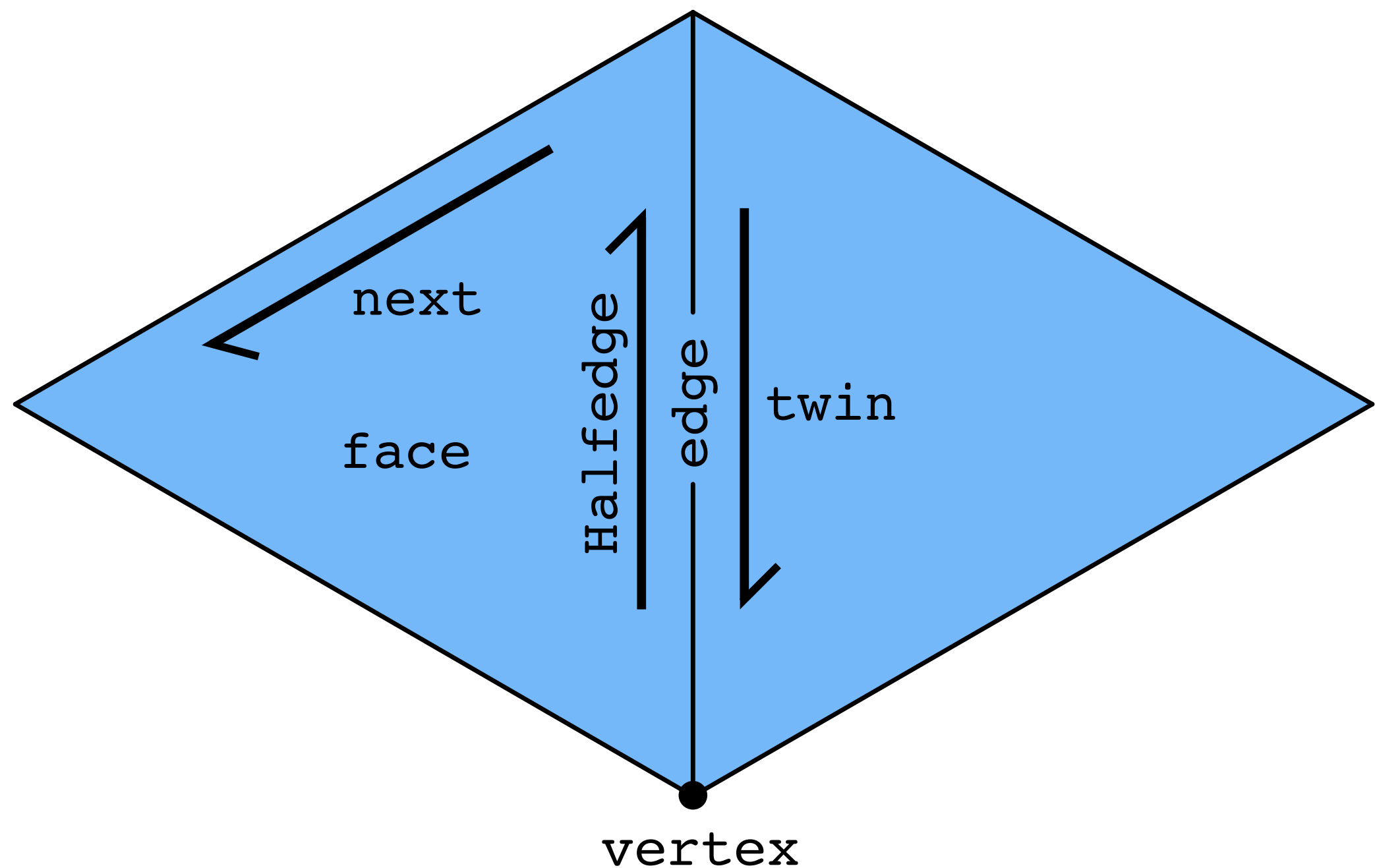


# Half-Edge Data Structure

```
struct Halfedge {  
    Halfedge *twin,  
    Halfedge *next;  
    Vertex *vertex;  
    Edge *edge;  
    Face *face;  
}  
struct Vertex {  
    Point pt;  
    Halfedge *halfedge;  
}  
struct Edge {  
    Halfedge *halfedge;  
}  
struct Face {  
    Halfedge *halfedge;  
}
```

CS184/284A

Key idea: two half-edges act as  
"glue" between mesh elements



Each vertex, edge and face points  
to one of its half edges

Ng & Kanazawa



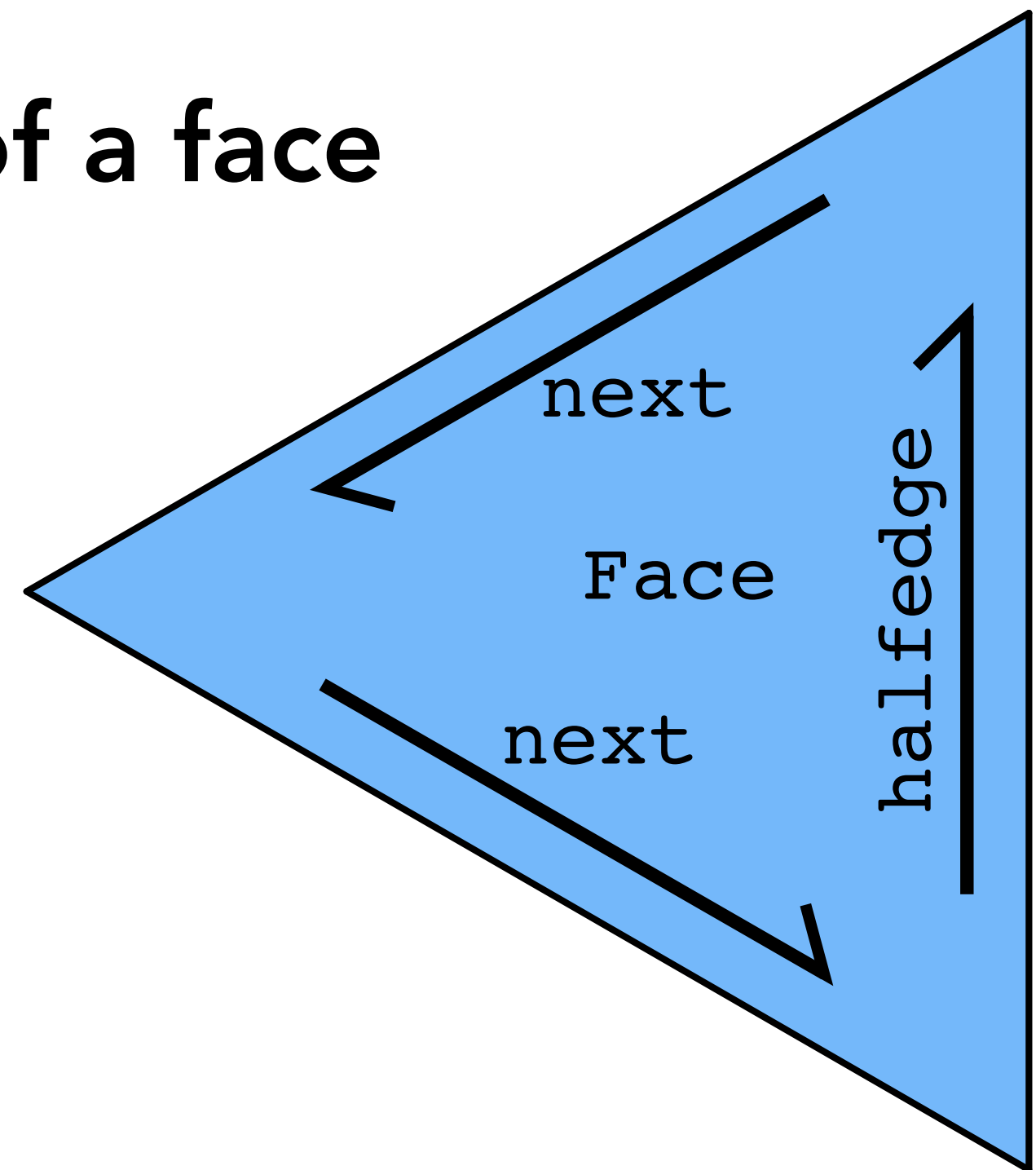
# Half-Edge Facilitates Mesh Traversal

Use twin and next pointers to move around mesh

Process vertex, edge and/or face pointers

Example 1: process all vertices of a face

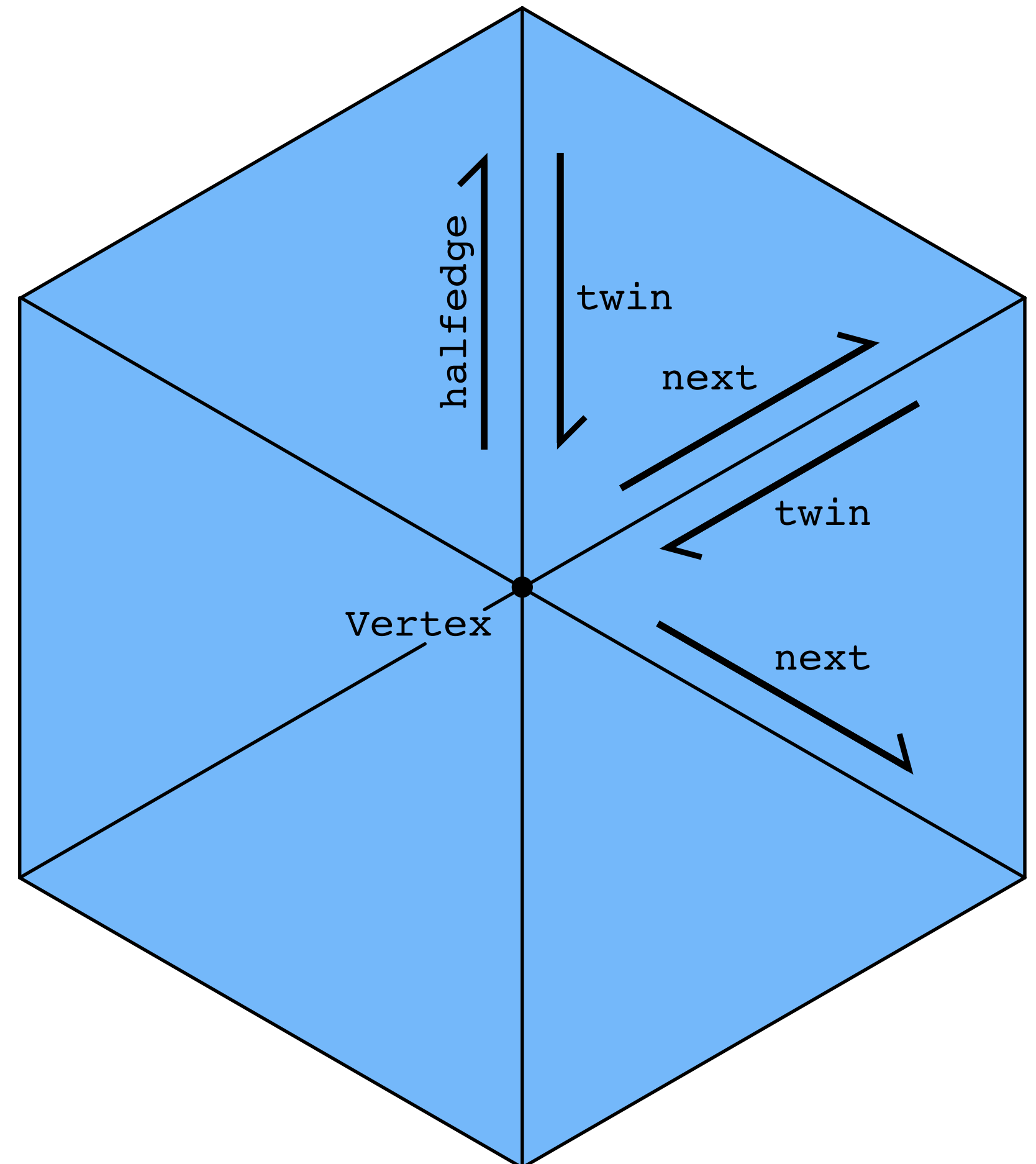
```
Halfedge* h = f->halfedge;  
do {  
    process(h->vertex);  
    h = h->next;  
}
```



# Half-Edge Facilitates Mesh Traversal

Example 2: process all edges around a vertex

```
Halfedge* h = v->halfedge;  
do {  
    process(h->edge);  
    h = h->twin->next;  
}
```



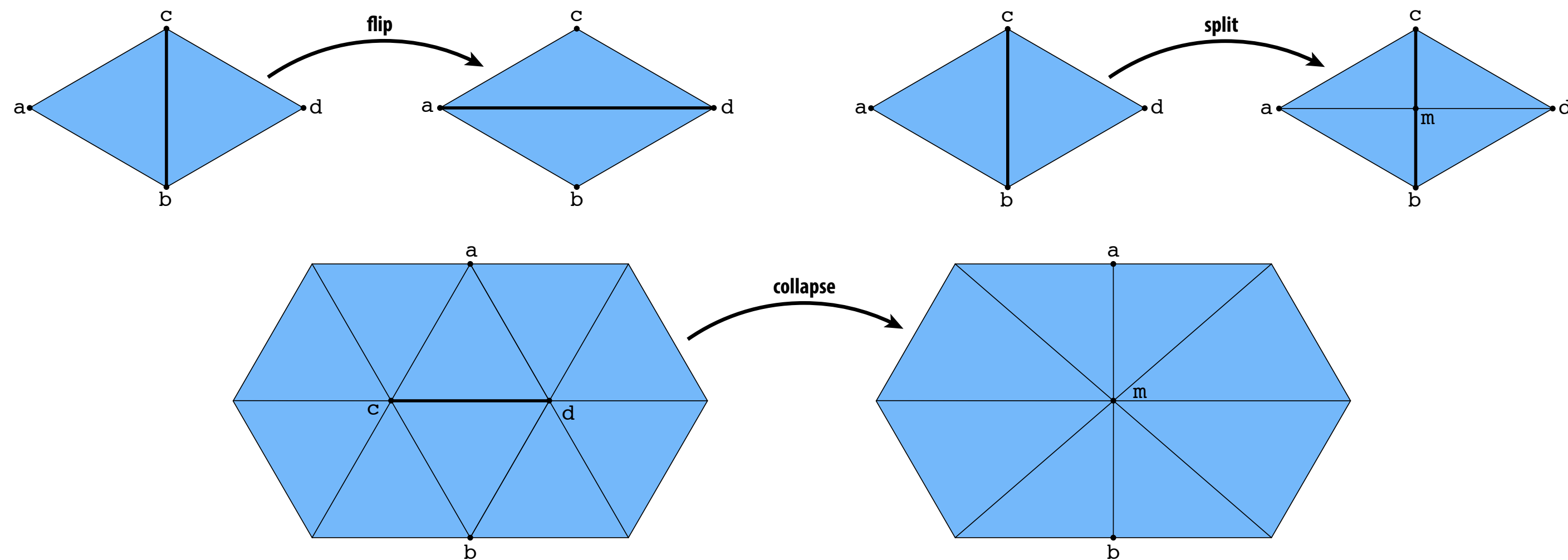


# **Local Mesh Operations**

# Half-Edge – Local Mesh Editing

Basic operations for linked list: insert, delete

Basic ops for half-edge mesh: flip, split, collapse edges



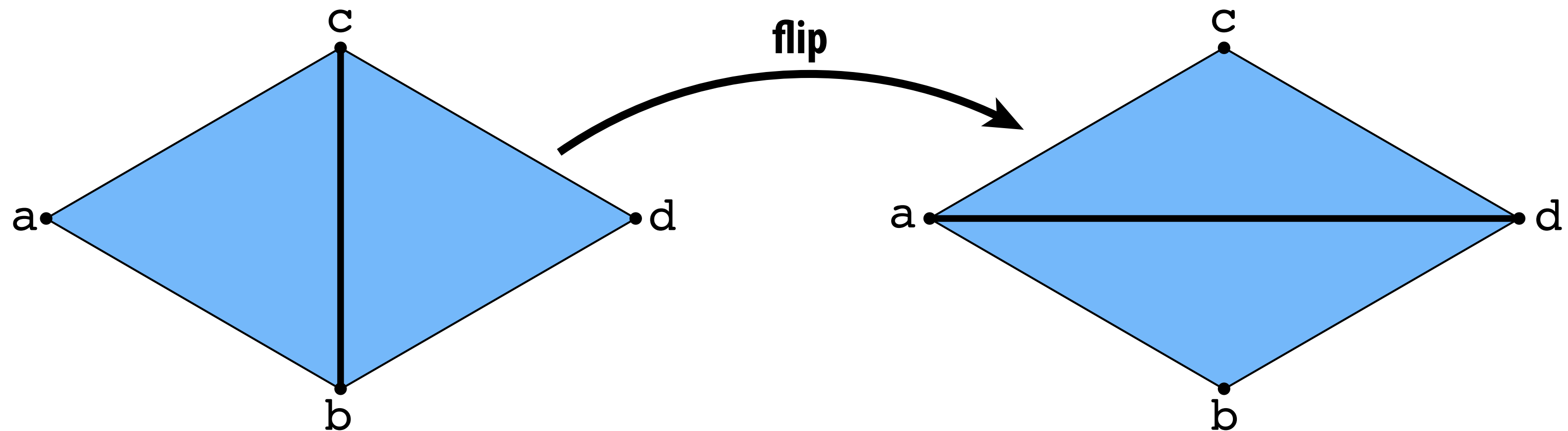
Allocate / delete elements; reassign pointers

(Care needed to preserve mesh manifold property)



# Half-Edge – Edge Flip

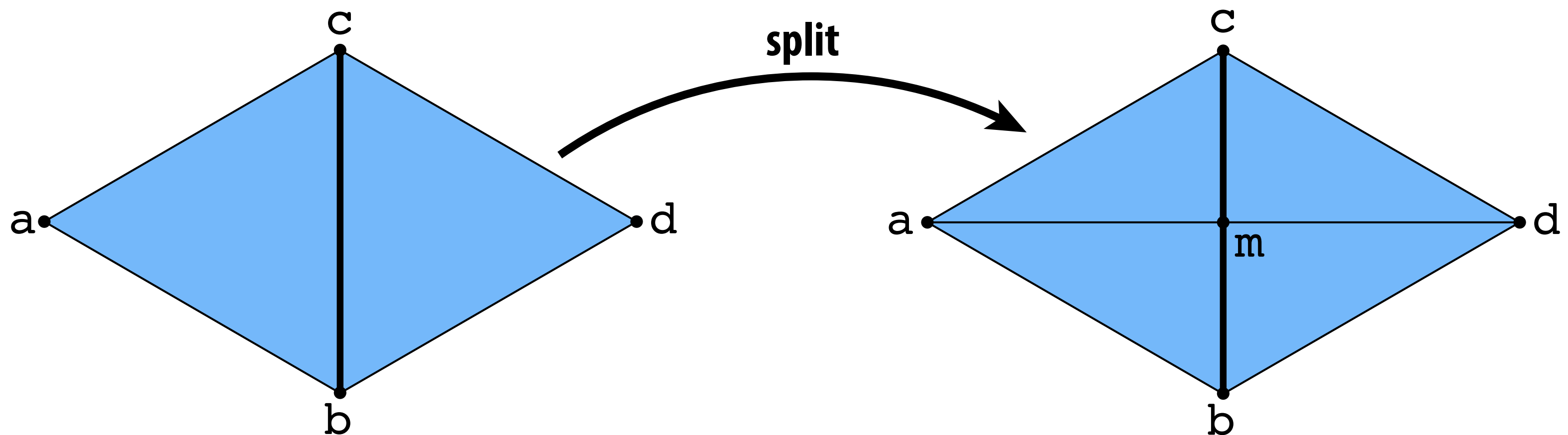
- Triangles  $(a,b,c)$ ,  $(b,d,c)$  become  $(a,d,c)$ ,  $(a,b,d)$ :



- Long list of pointer reassignments
- However, no elements created/destroyed.

# Half-Edge – Edge Split

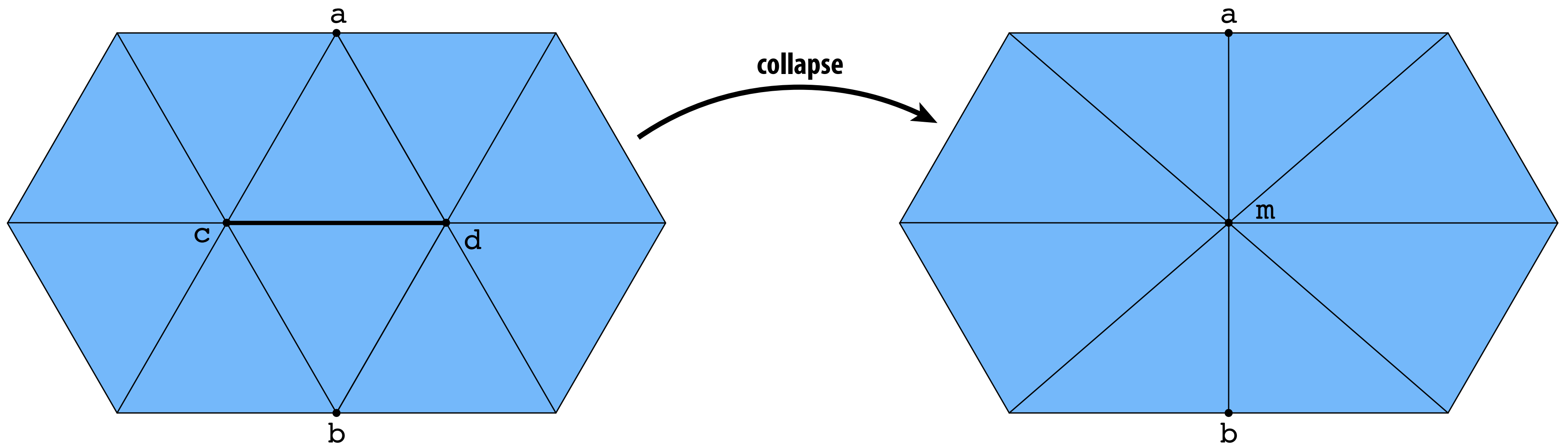
- Insert midpoint  $m$  of edge  $(c,b)$ , connect to get four triangles:



- This time have to add elements
- Again, many pointer reassignments

# Half-Edge – Edge Collapse

- Replace edge (c,d) with a single vertex m:

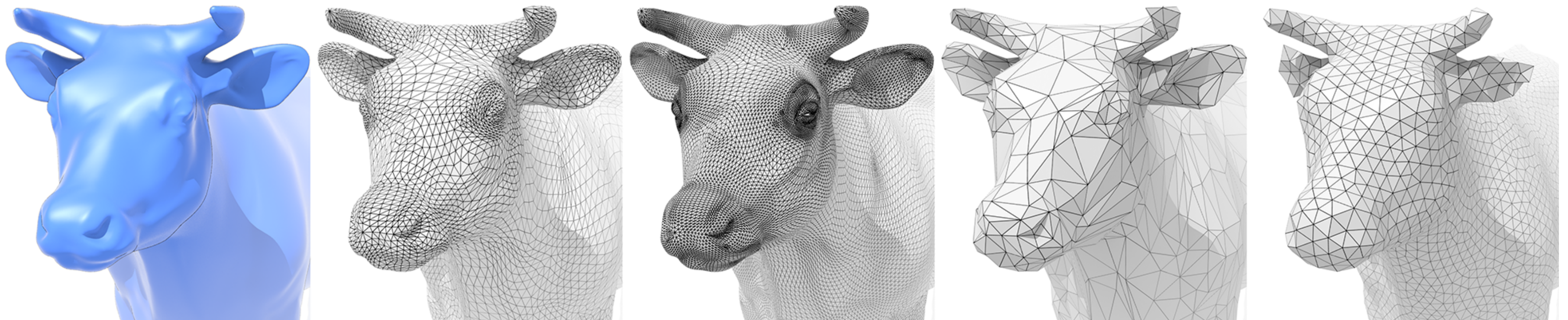


- This time have to delete elements
- Again, many pointer reassignments



# Global Mesh Operations: Geometry Processing

- Mesh subdivision
- Mesh simplification
- Mesh regularization



# **Subdivision Surfaces**



# Subdivision Surfaces

Start with coarse polygon mesh ("control cage")

- Subdivide each element
- Update vertices via local averaging

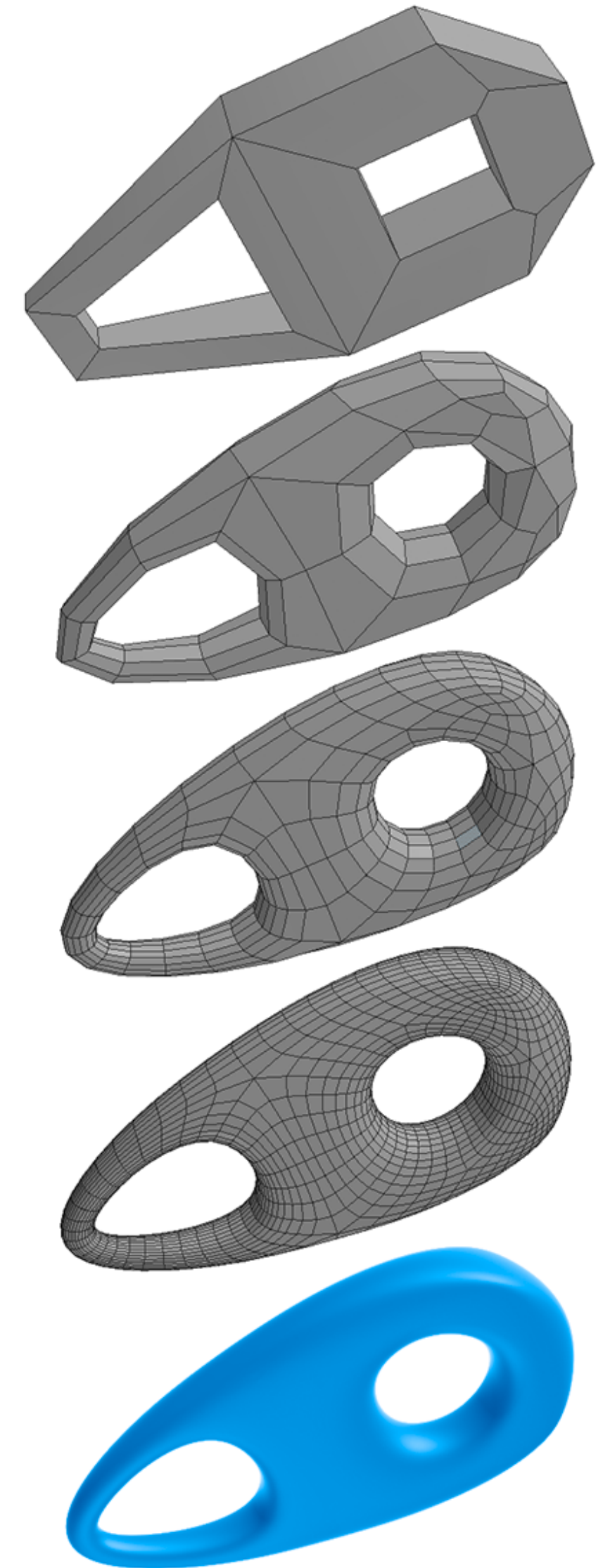
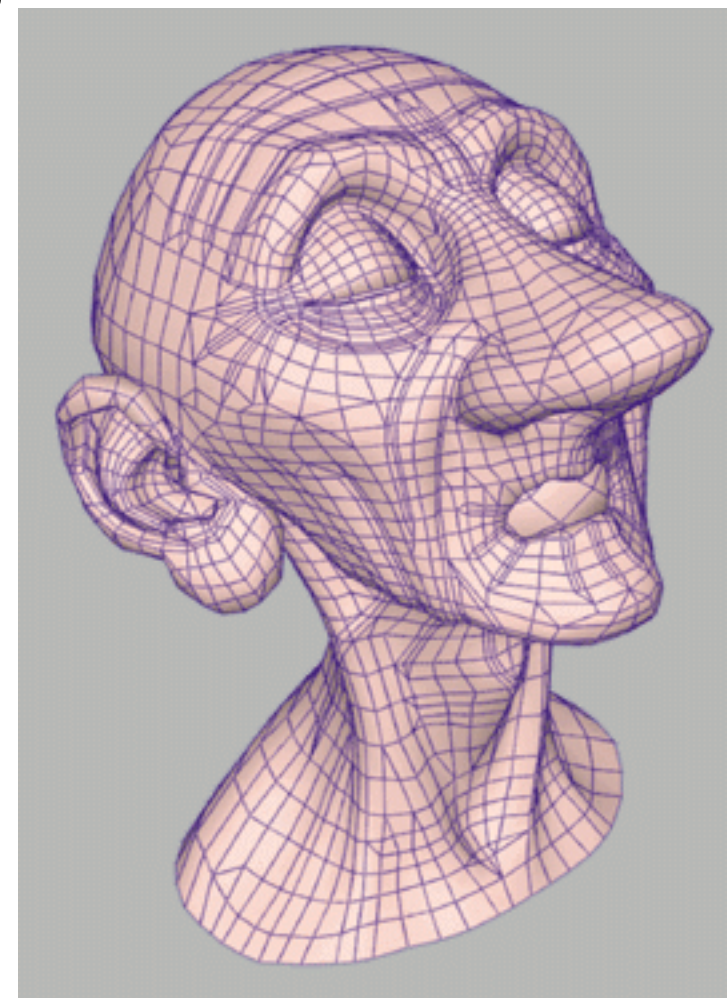
Many possible rule:

- Catmull-Clark (quads)
- Loop (triangles)
- ...

Common issues:

- interpolating or approximating?
- continuity at vertices?

Relatively easy for modeling; harder to guarantee continuity





# Core Idea: Let Subdivision Define The Surface

In Bezier curves, we saw:

- Evaluation by subdivision (de Casteljau algorithm)
- Or evaluation by algebra (Bernstein polynomials)

Insight that leads to subdivision surfaces:

- Free ourselves from the algebraic evaluation
- Let subdivision fully define the surface

Many possible subdivision rules – different surfaces

- Technical challenge shifts to designing rules and proving properties (e.g. convergence and continuity)
- Applying rules to compute surface is procedural

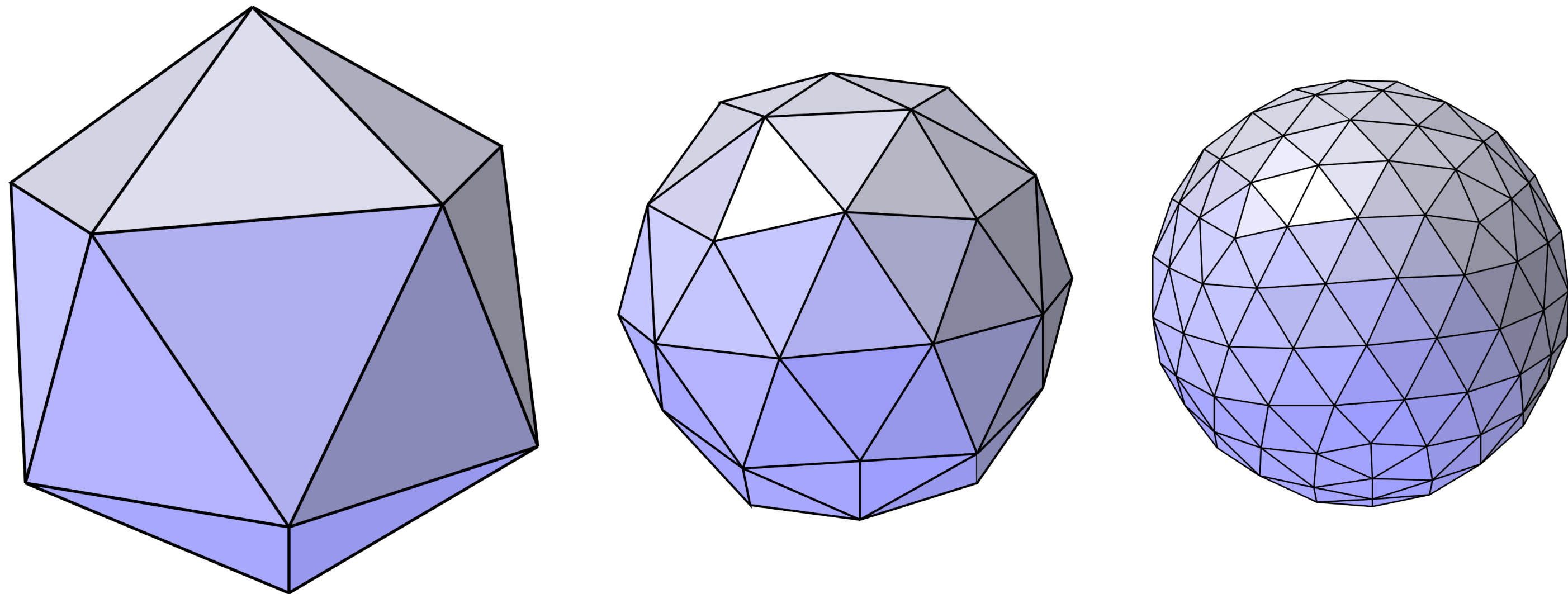
# Loop Subdivision

# Loop Subdivision

Common subdivision rule for triangle meshes

"C2" smoothness away from extraordinary vertices

Approximating, not interpolating

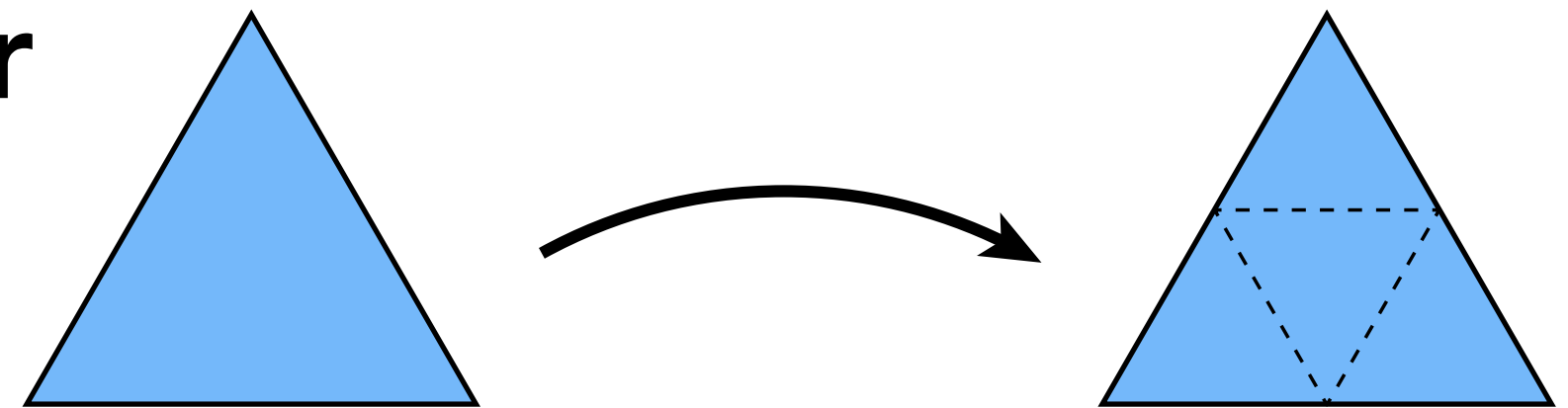


Simon Fuhrman

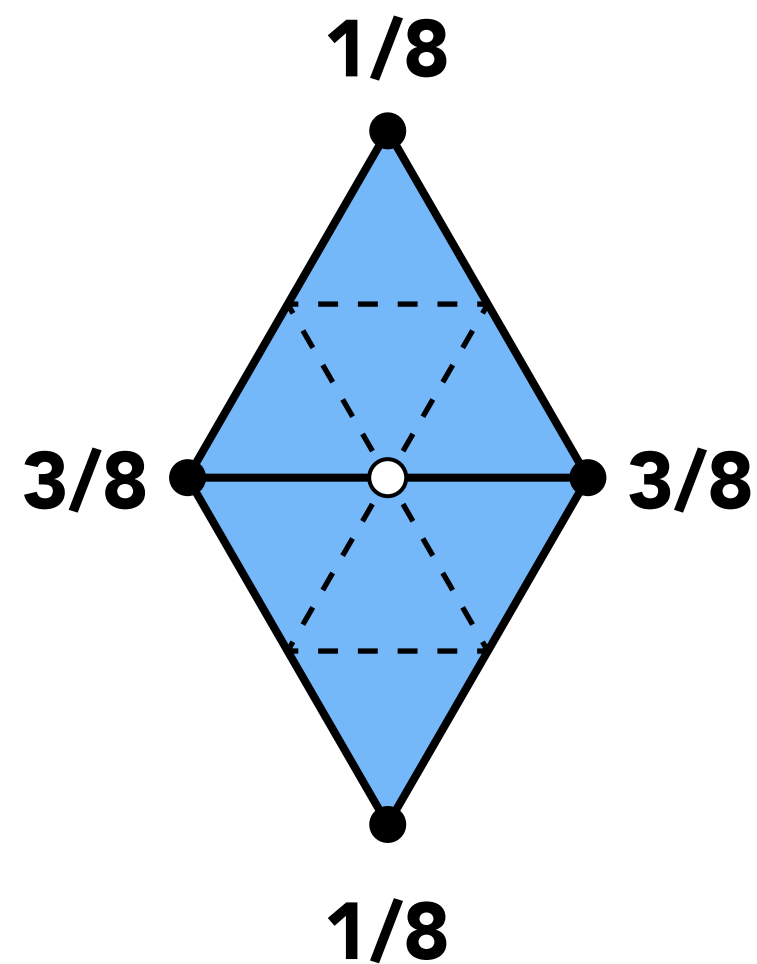


# Loop Subdivision Algorithm

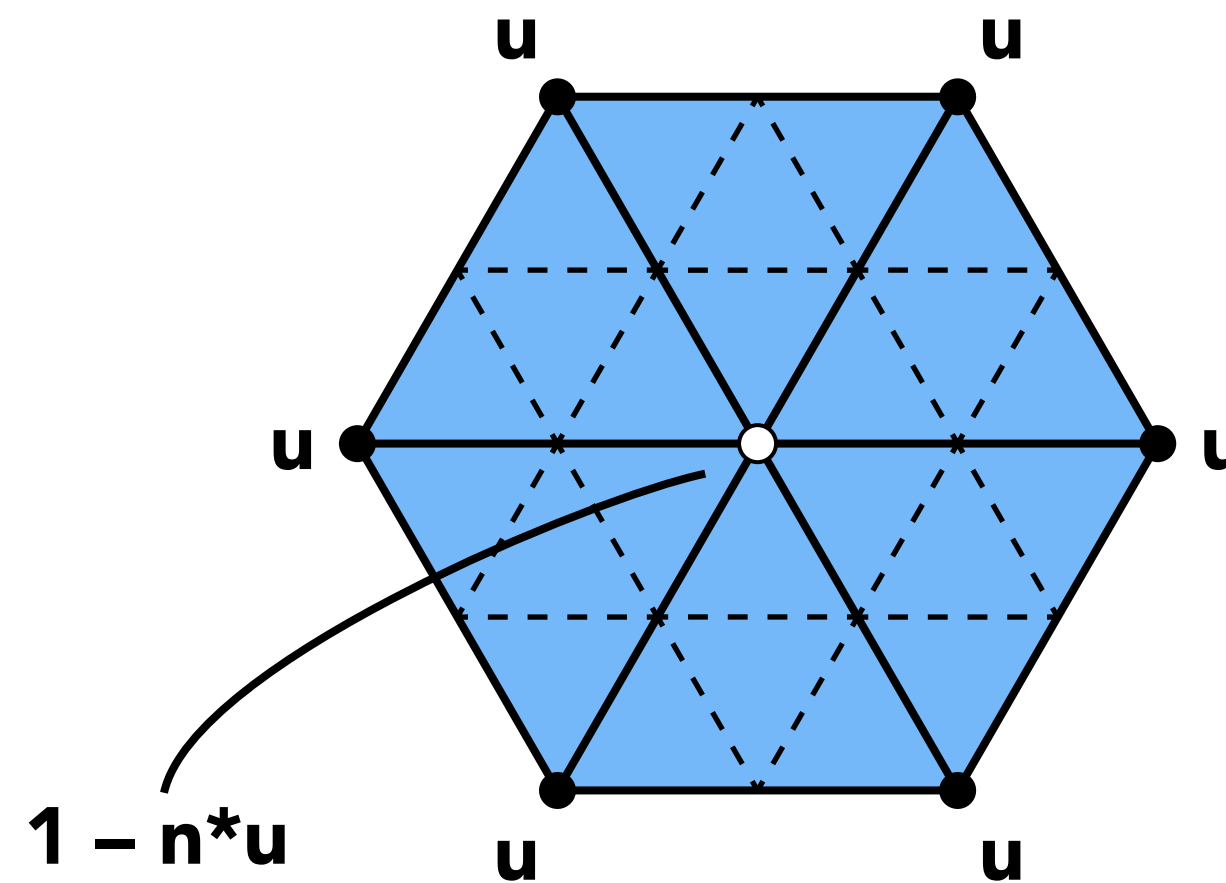
- Split each triangle into four



- Assign new vertex positions according to weights:



New vertices



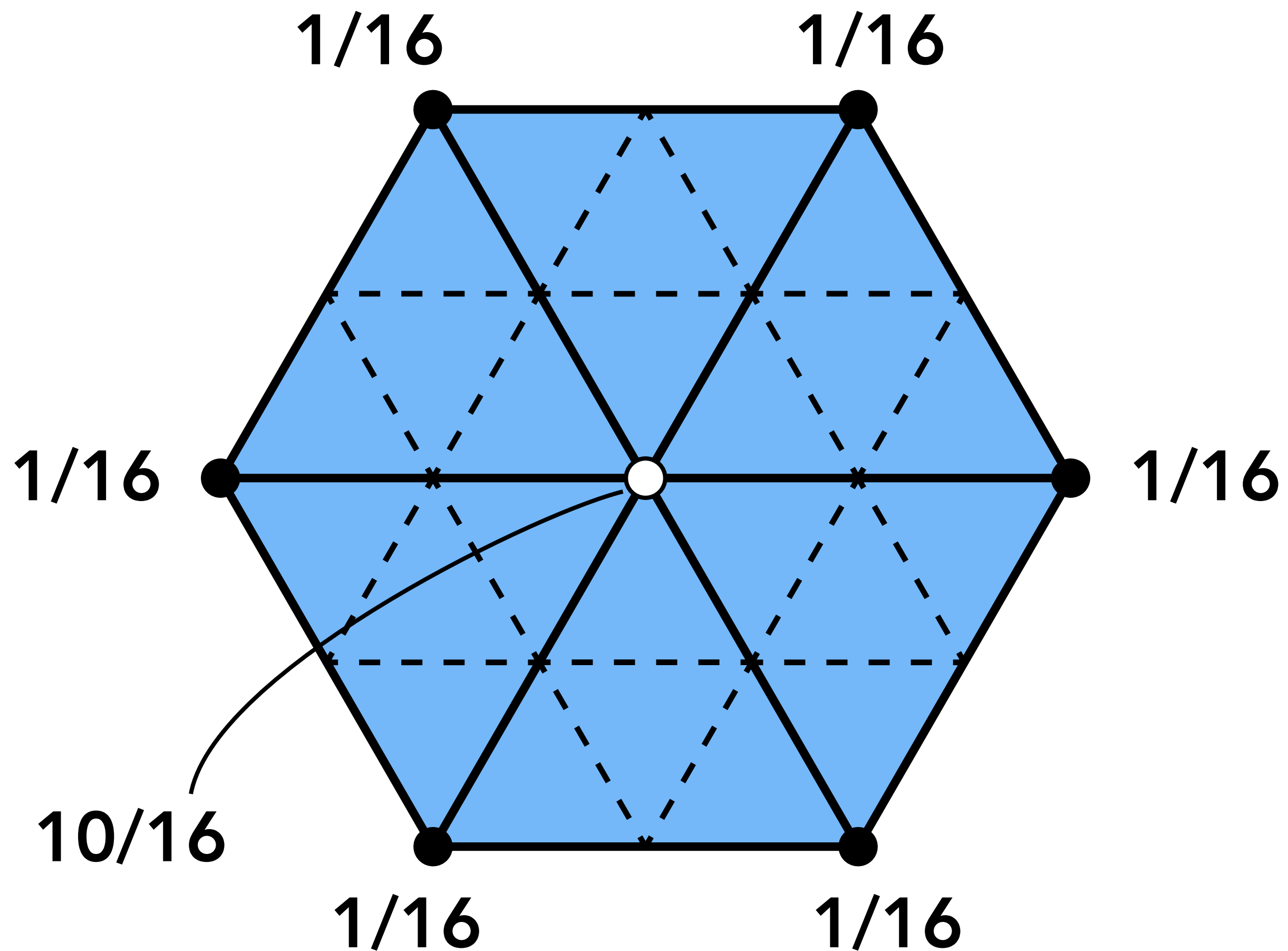
Old vertices

$n$ : vertex degree

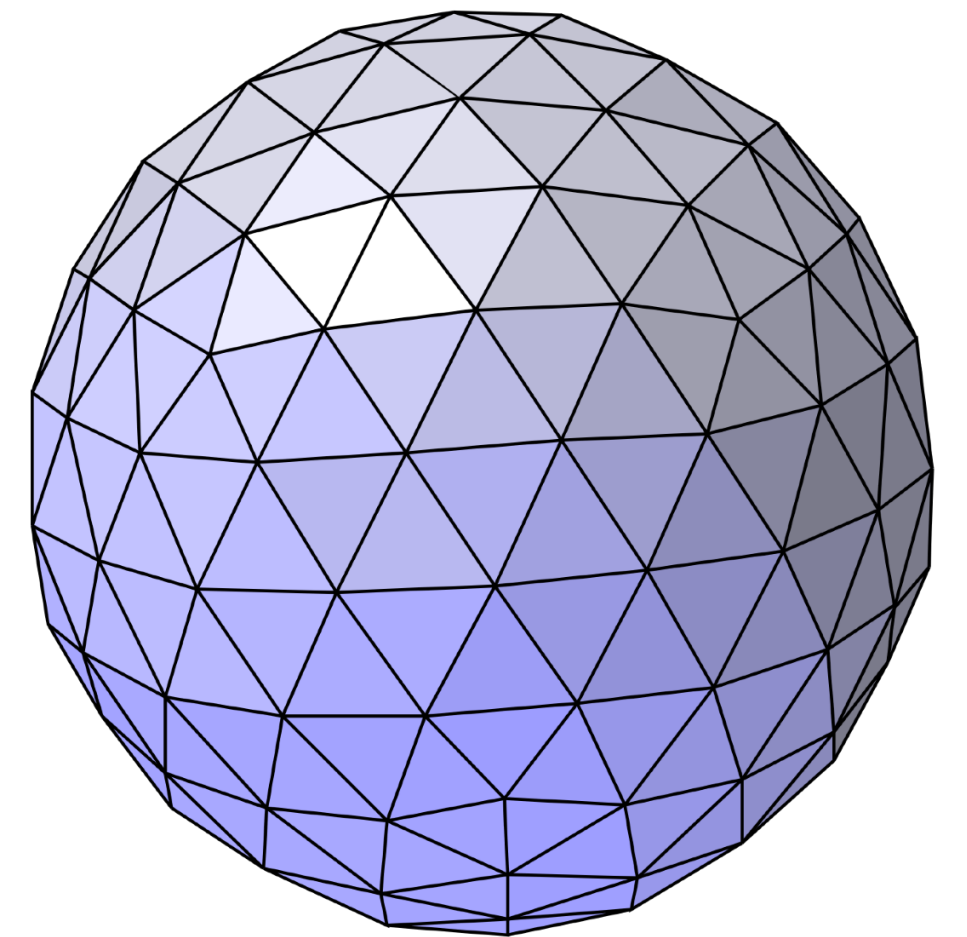
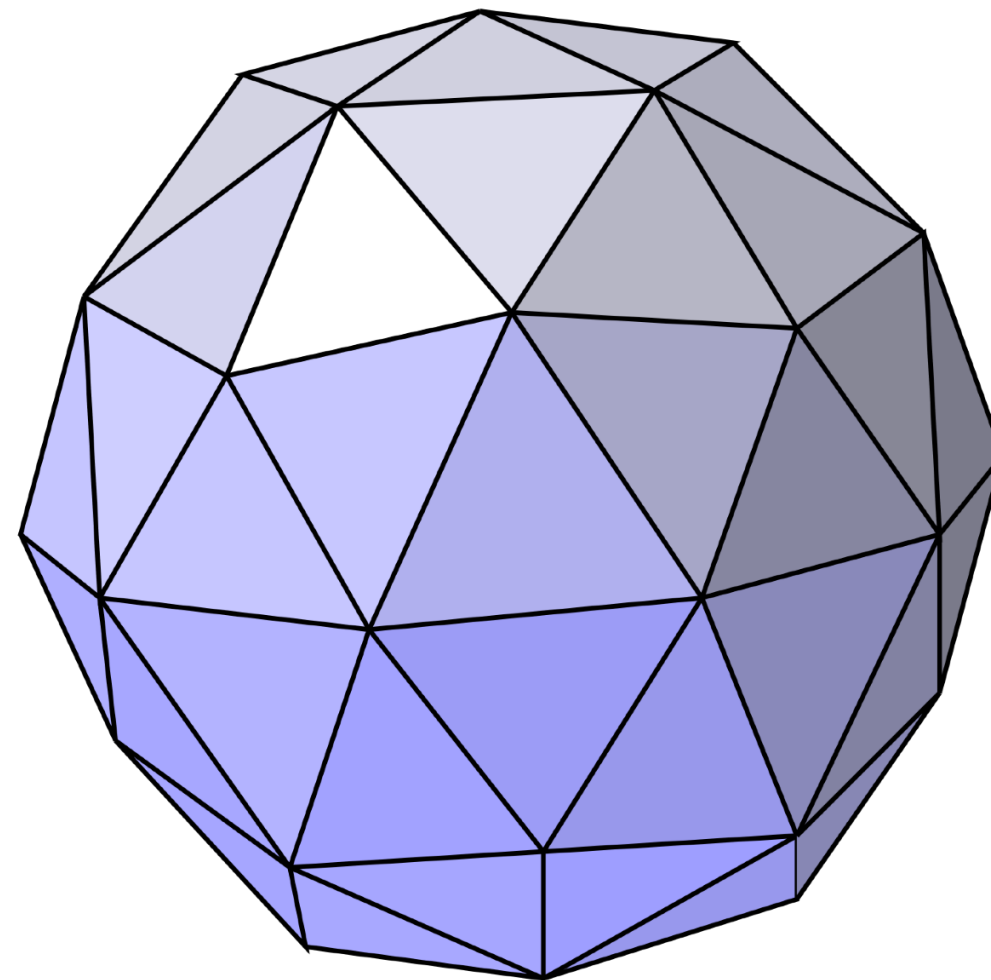
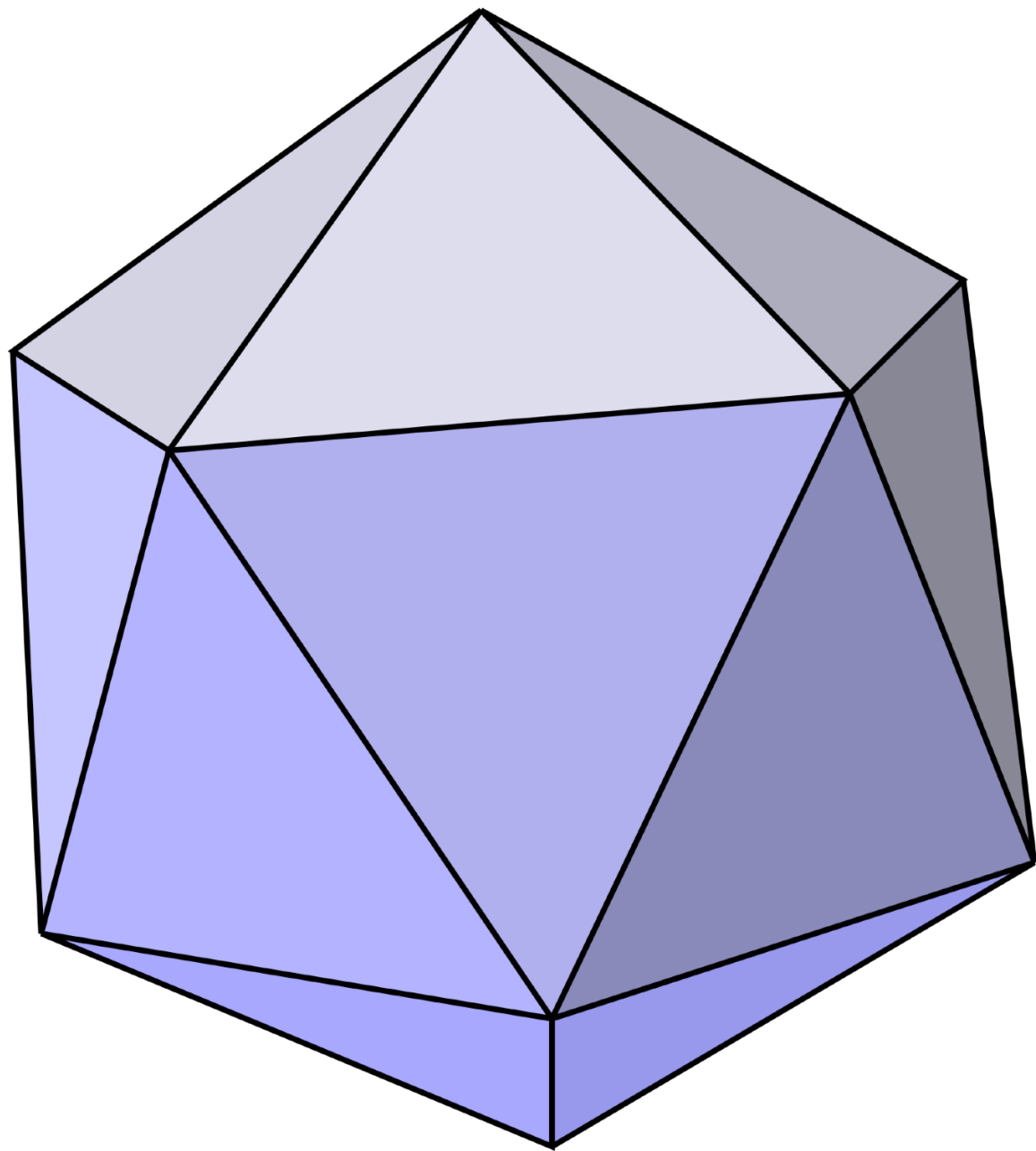
$u$ :  $3/16$  if  $n=3$ ,  $3/(8n)$  otherwise

# Loop Subdivision Algorithm

Example, for degree 6 vertices



# Loop Subdivision Algorithm



Simon Fuhrman



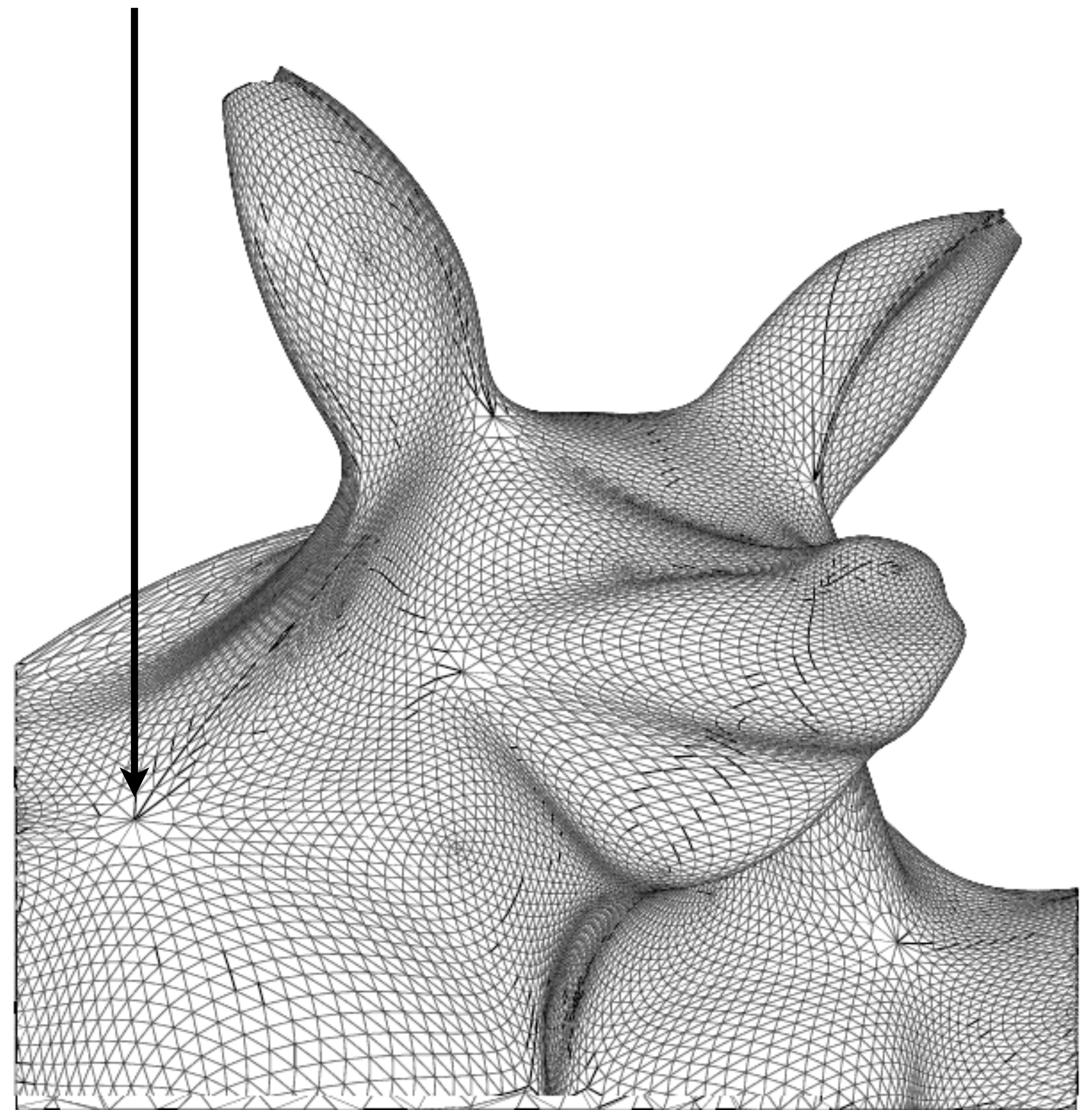
# Semi-Regular Meshes

Most of the mesh has vertices with degree 6

But if the mesh is topologically equivalent to a sphere, then not all the vertices can have degree 6

Must have a few extraordinary points (degree not equal to 6)

Extraordinary point



# Proof: Always an Extraordinary Vertex

Our mesh (topologically equivalent to sphere) has  $V$  vertices,  $E$  edges, and  $T$  triangles

$$E = \frac{3}{2} T$$

- There are 3 edges per triangle, and each edge is part of 2 triangles
- Therefore  $E = \frac{3}{2}T$

$$T = 2V - 4$$

- Euler Convex Polyhedron Formula:  $T - E + V = 2$
- $\Rightarrow V = \frac{3}{2}T - T + 2 \Rightarrow T = 2V - 4$

If all vertices had 6 triangles,  $T = 2V$

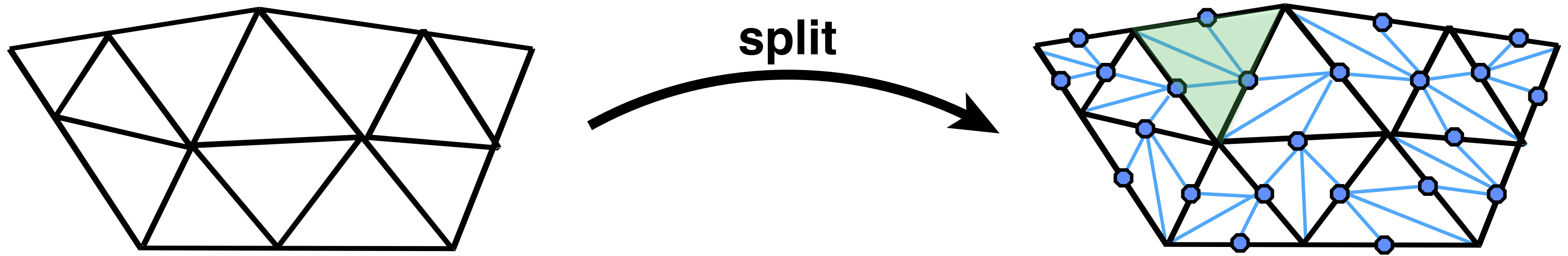
- There are 6 edges per vertex, and every edge connects 2 vertices
- Therefore,  $E = \frac{6}{2}V \Rightarrow \frac{3}{2}T = \frac{6}{2}V \Rightarrow T = 2V$

$T$  cannot equal both  $2V - 4$  and  $2V$ , a contradiction

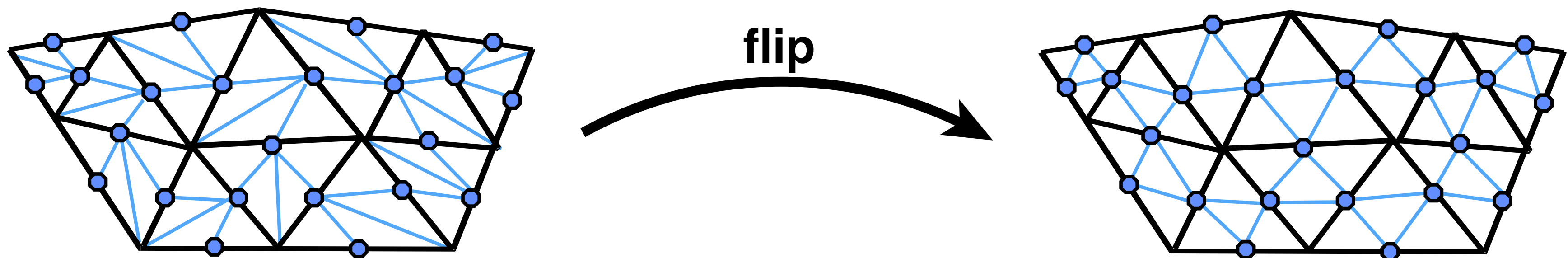
- Therefore, the mesh cannot have 6 triangles for every vertex

# Loop Subdivision via Edge Operations

First, split edges of original mesh in any order:



Next, flip new edges that touch a new & old vertex:



**(Don't forget to update vertex positions!)**

# Continuity of Loop Subdivision Surface

At extraordinary points

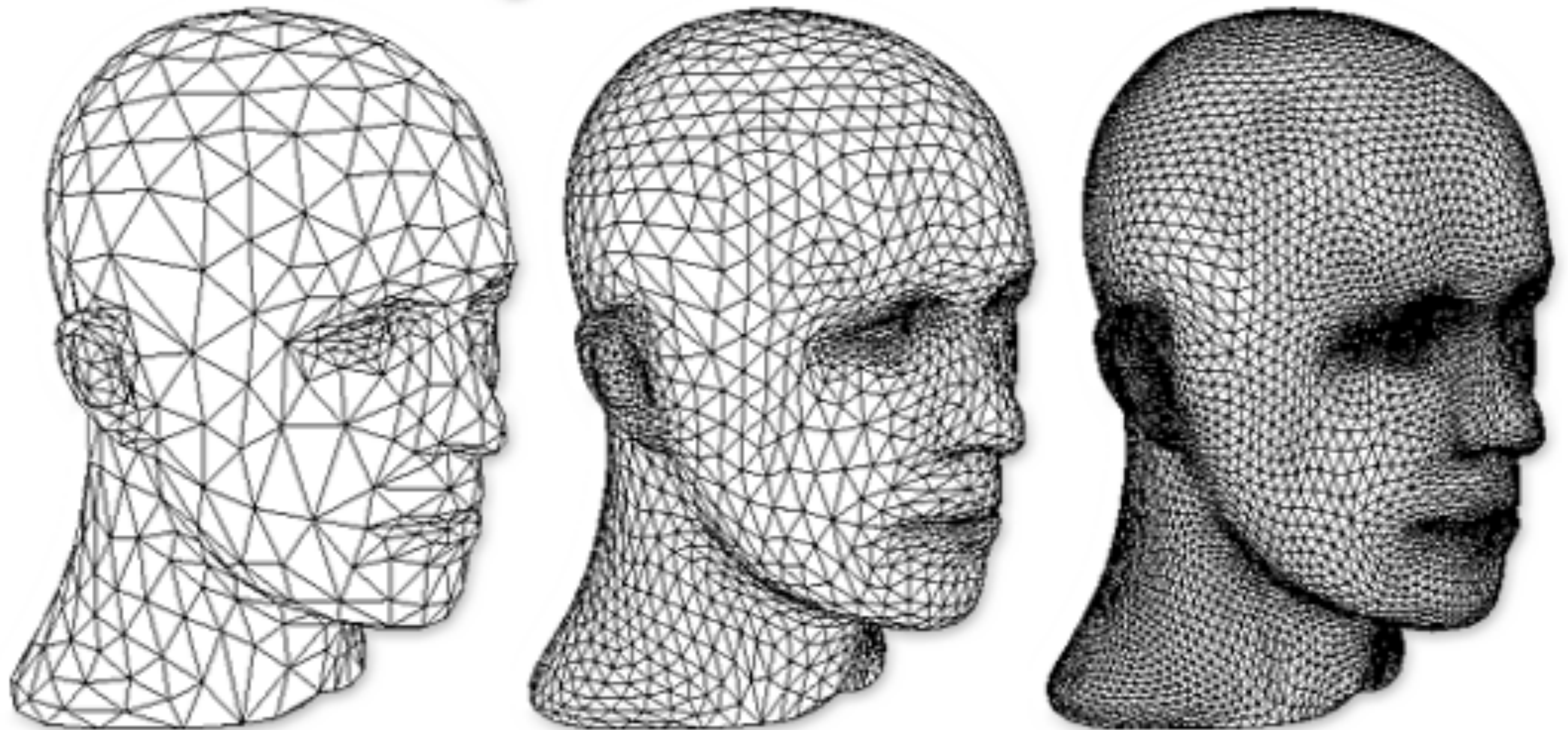
- Surface is at least  $C^1$  continuous

Everywhere else ("ordinary" regions)

- Surface is  $C^2$  continuous

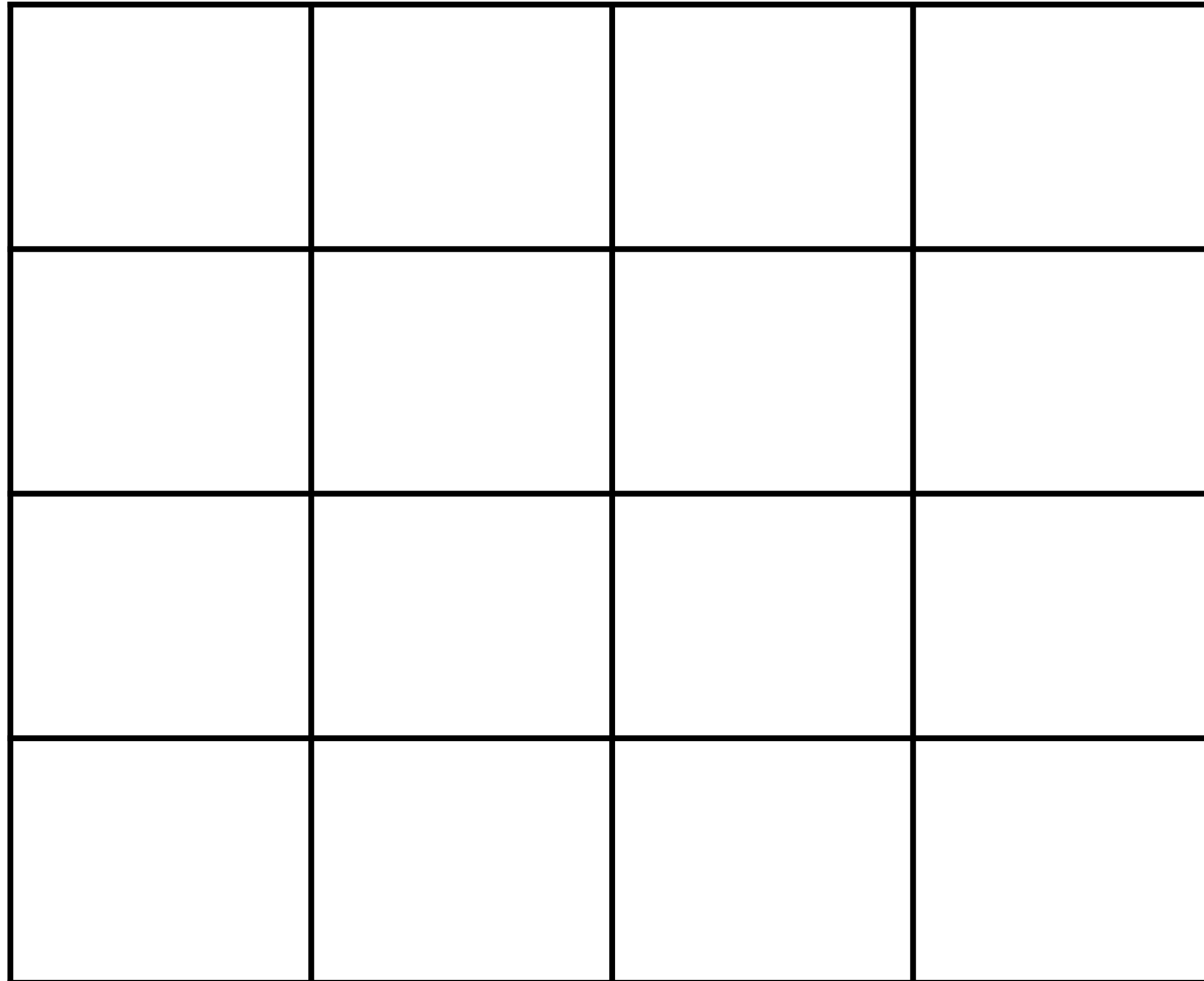


# Loop Subdivision Results

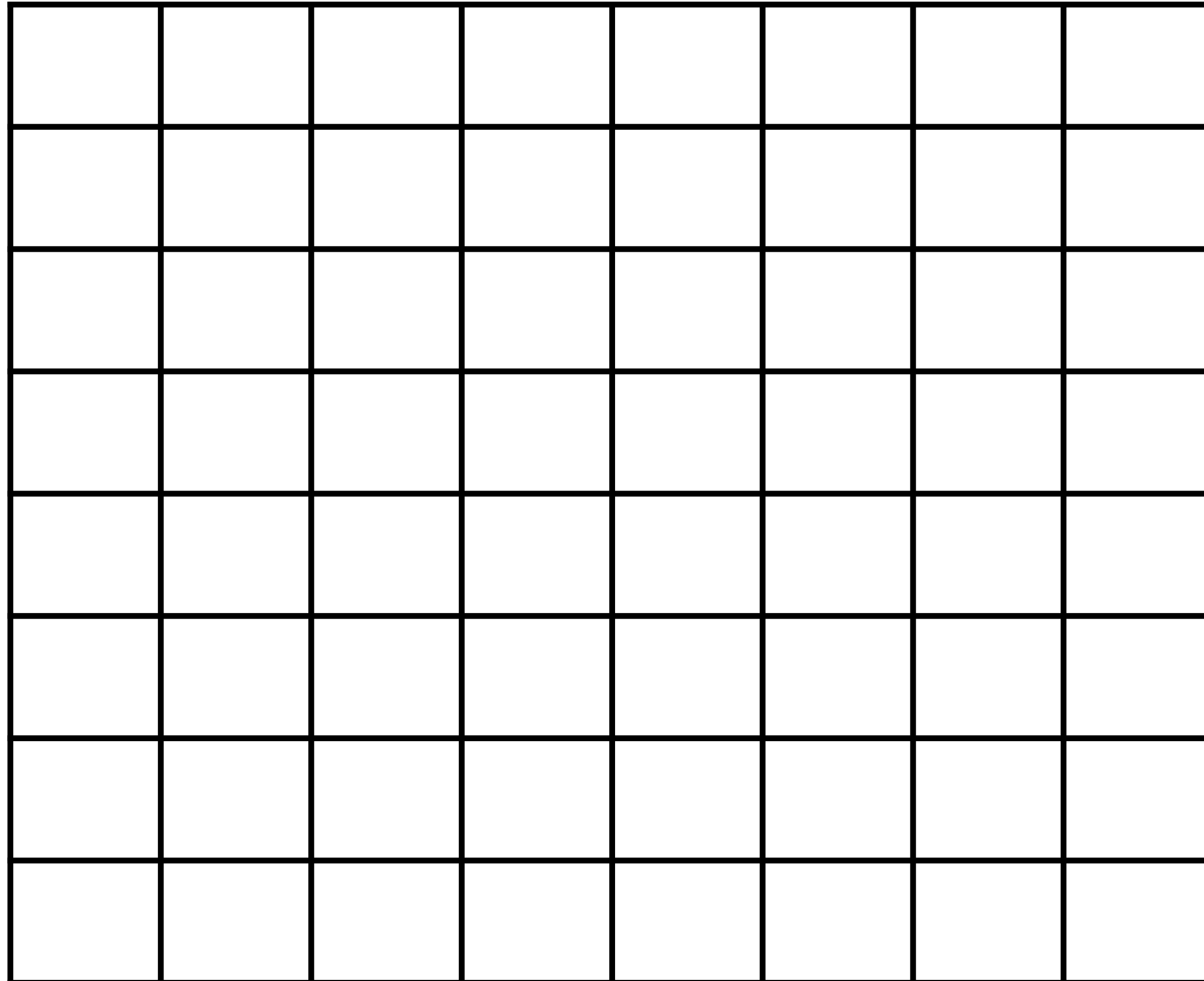


# **Catmull-Clark Subdivision**

# Catmull-Clark Subdivision (Regular Quad Mesh)

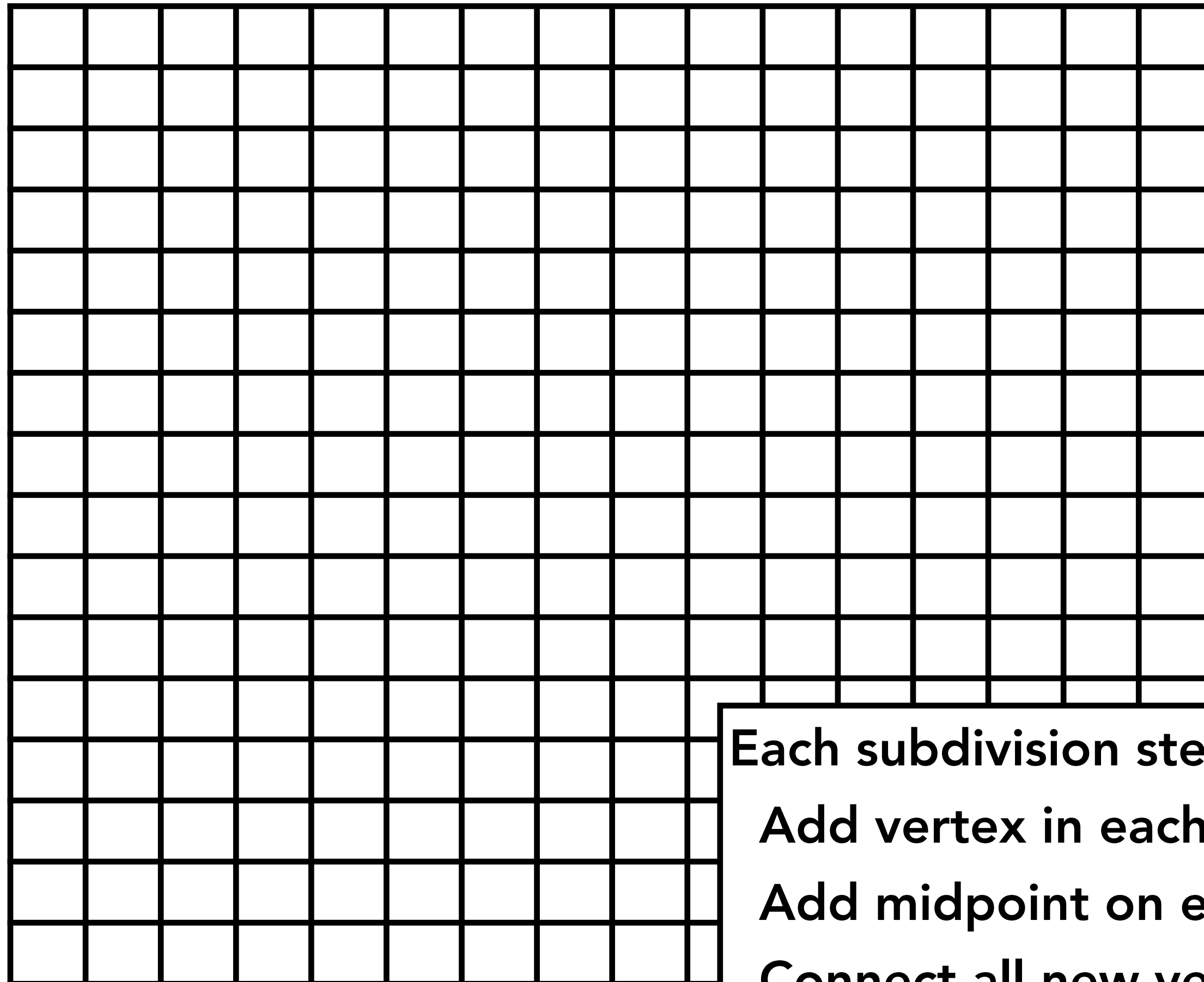


# Catmull-Clark Subdivision (Regular Quad Mesh)





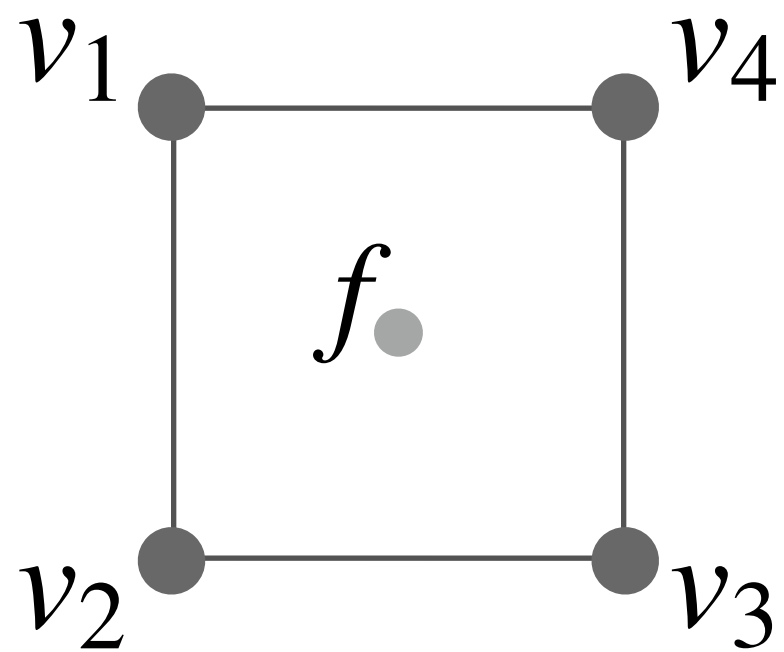
# Catmull-Clark Subdivision (Regular Quad Mesh)



Each subdivision step:  
Add vertex in each face  
Add midpoint on each edge  
Connect all new vertices

# Catmull-Clark Vertex Update Rules (Quad Mesh)

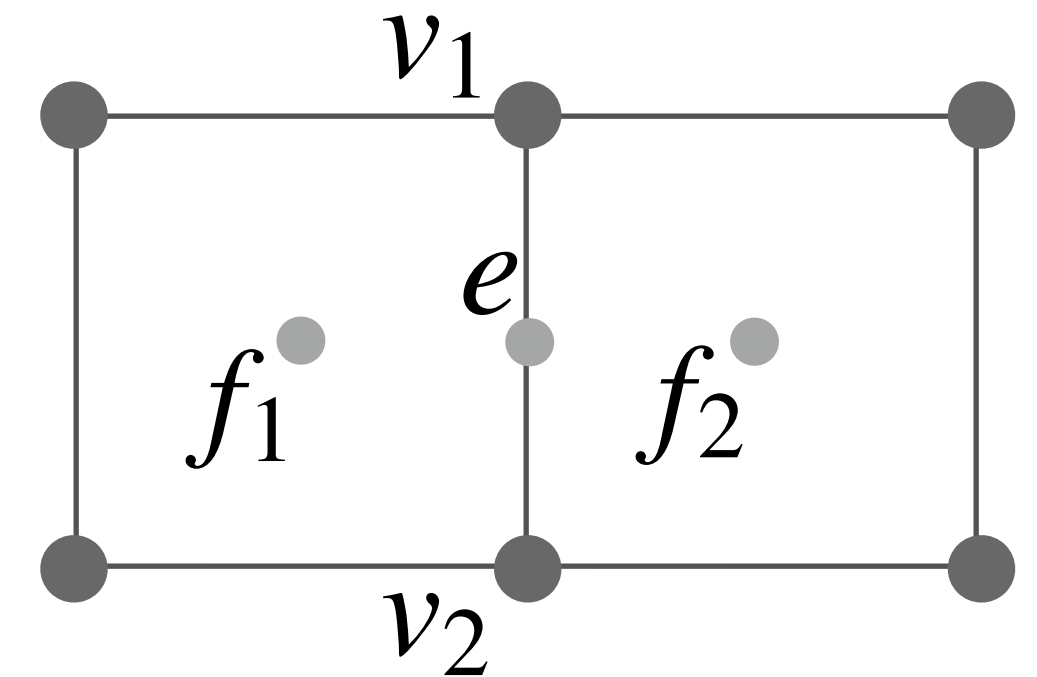
Face point



$$f = \frac{v_1 + v_2 + v_3 + v_4}{4}$$

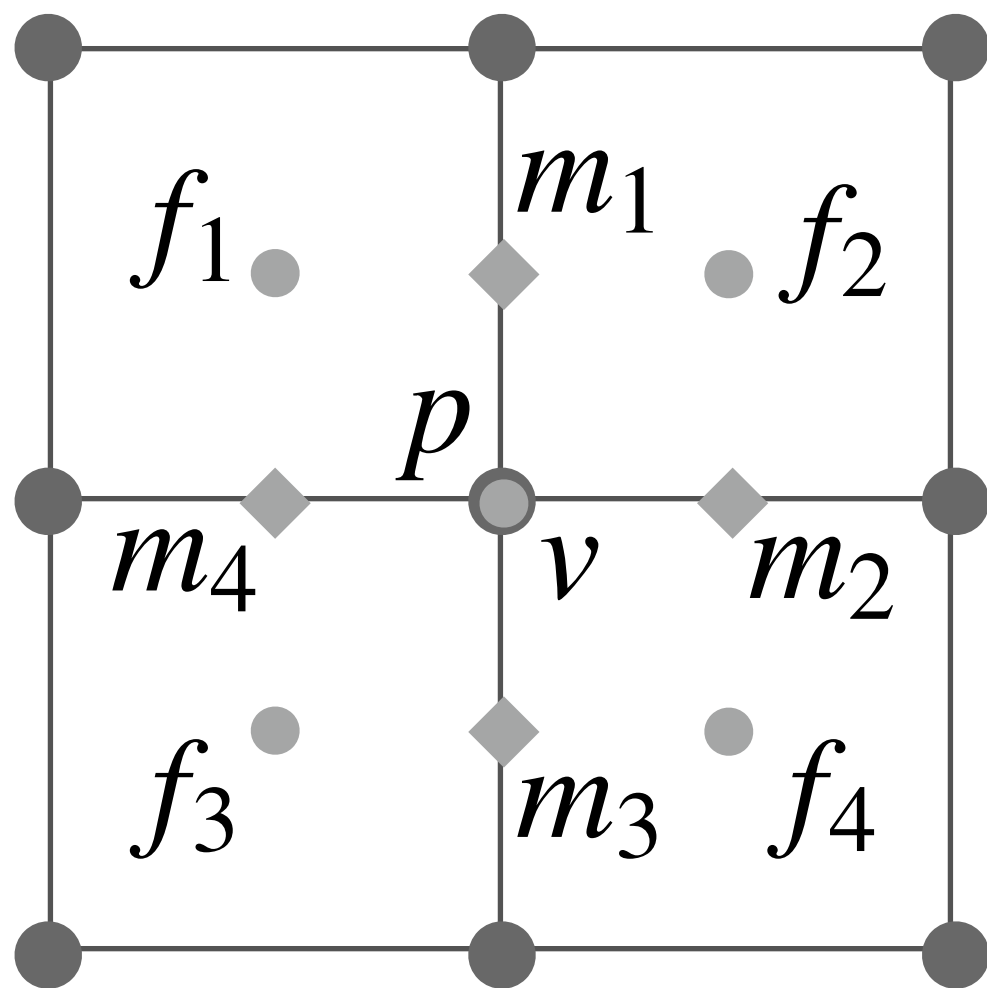
$$e = \frac{v_1 + v_2 + f_1 + f_2}{4}$$

Edge point



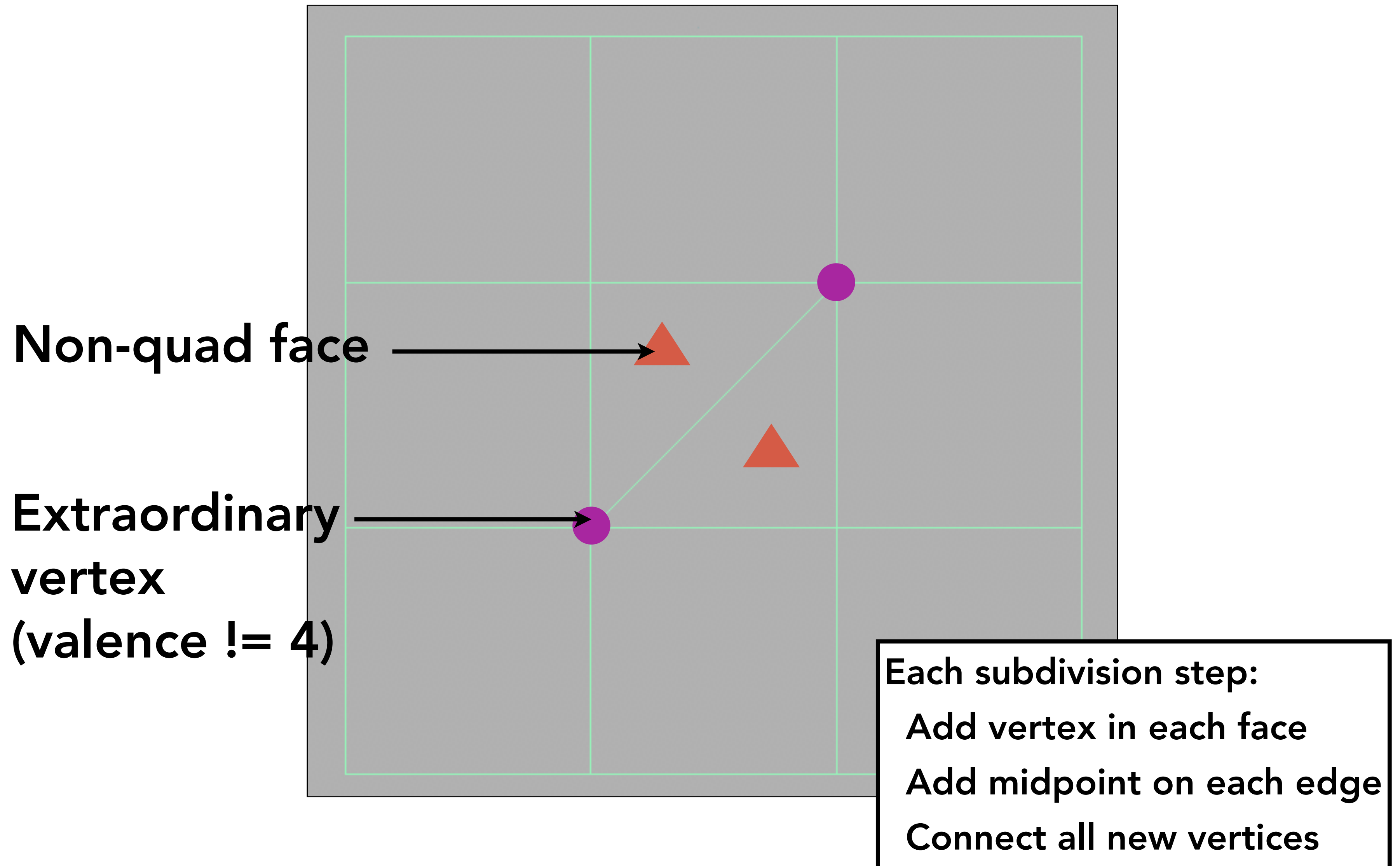
Vertex point

$$v = \frac{f_1 + f_2 + f_3 + f_4 + 2(m_1 + m_2 + m_3 + m_4) + 4p}{16}$$



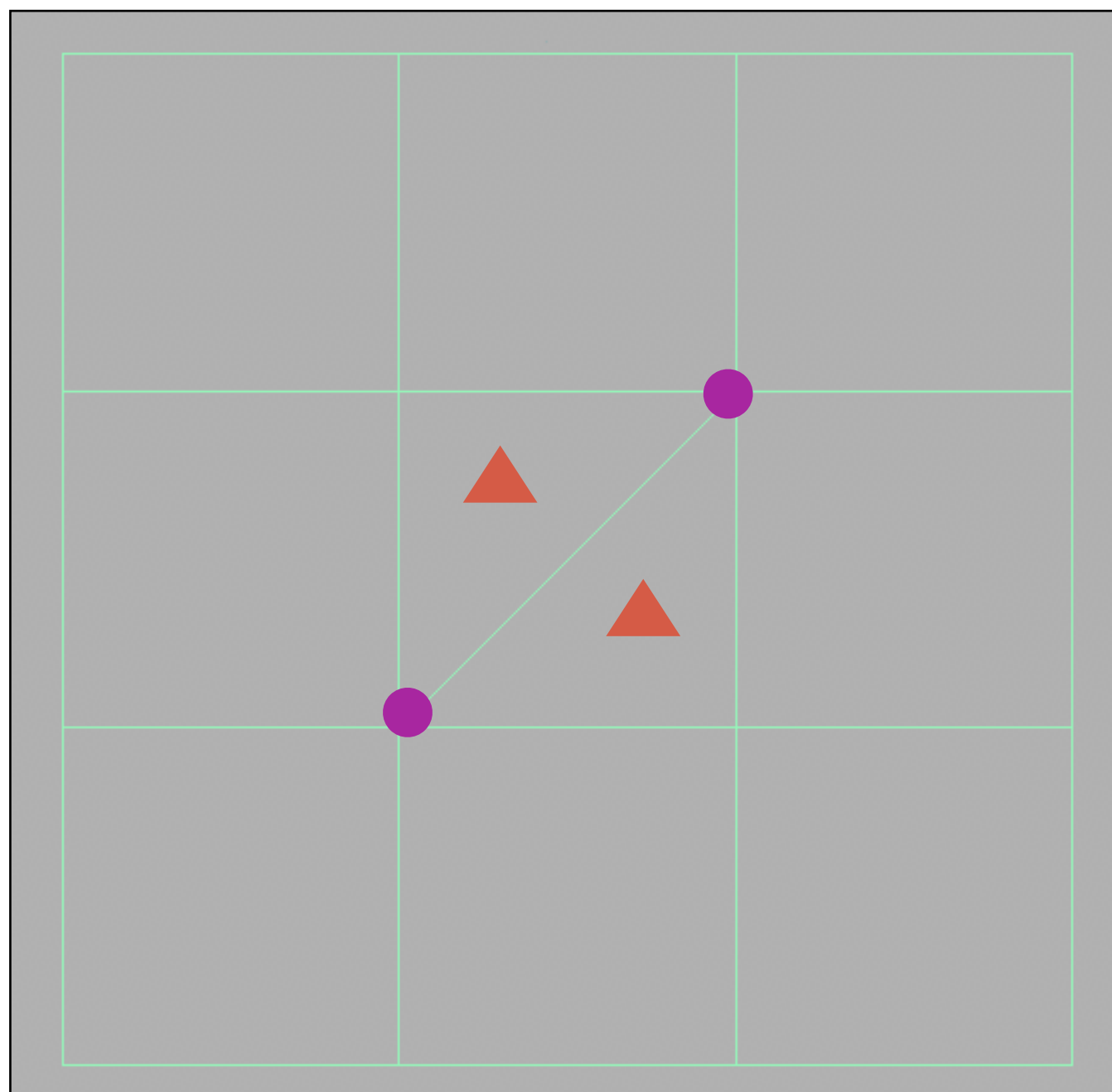
$m$  midpoint of edge, not "edge point"  
 $p$  old "vertex point"

# Catmull-Clark Subdivision (General Mesh)

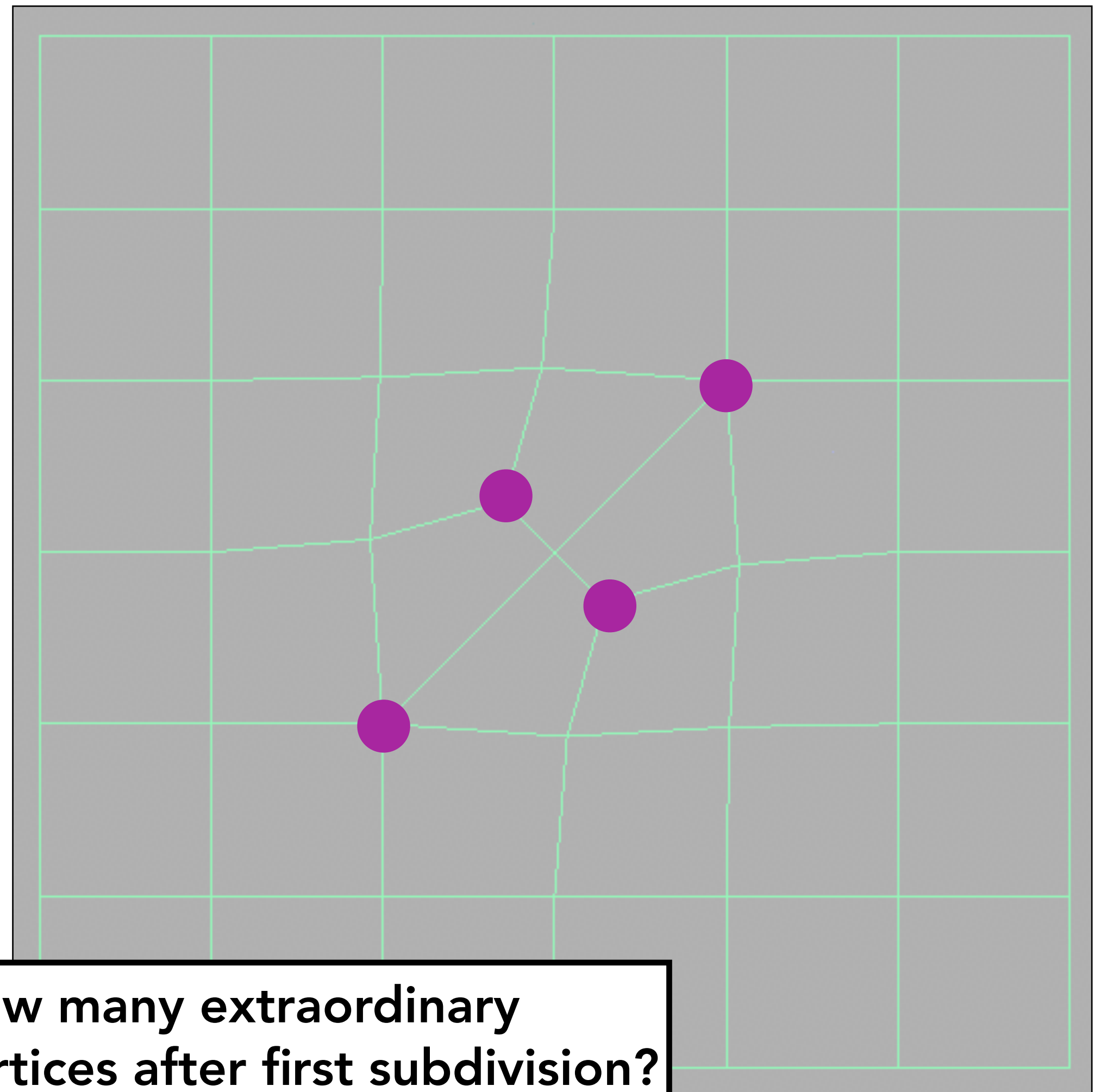


# Catmull-Clark Subdivision (General Mesh)

Before



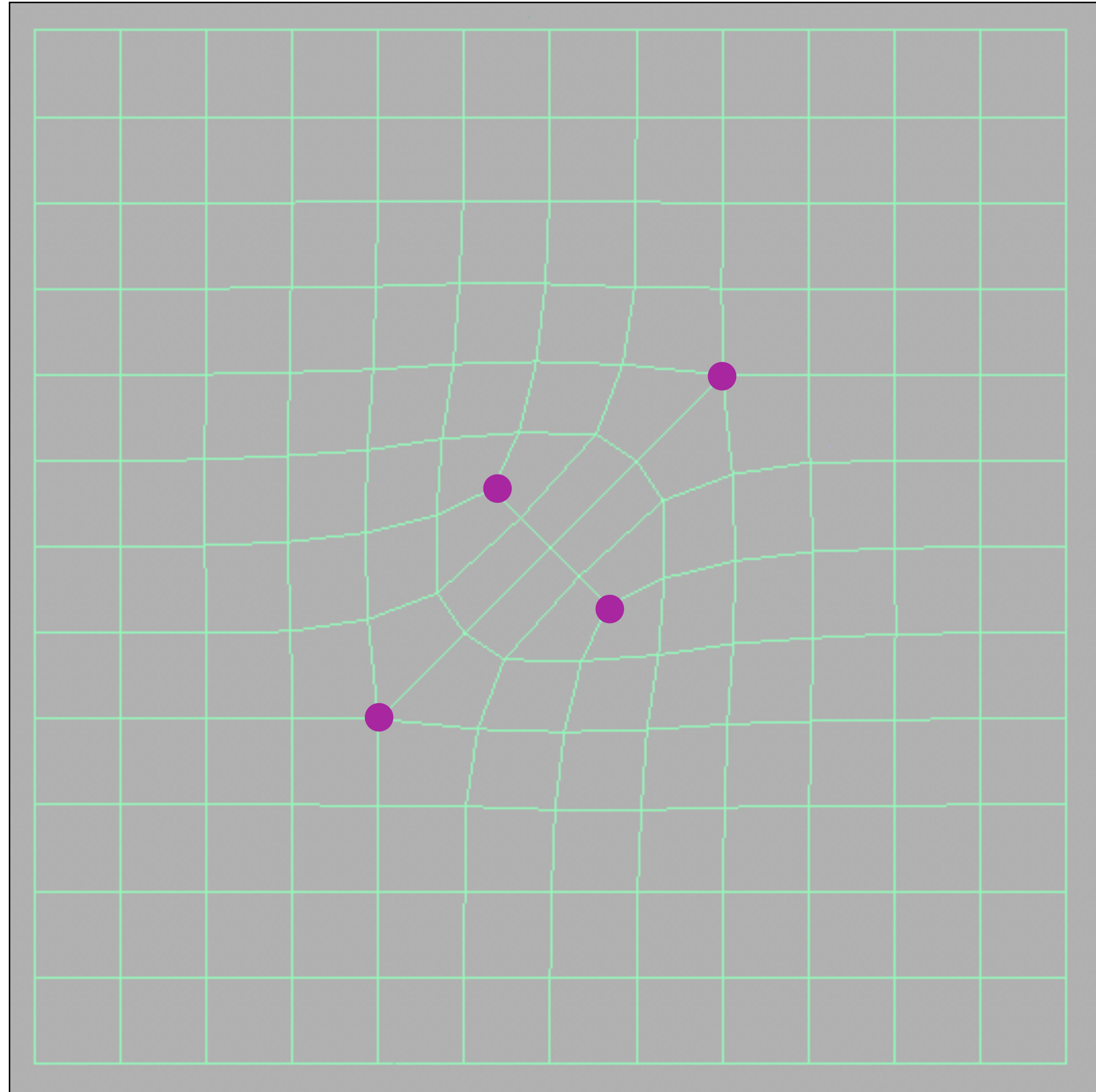
After



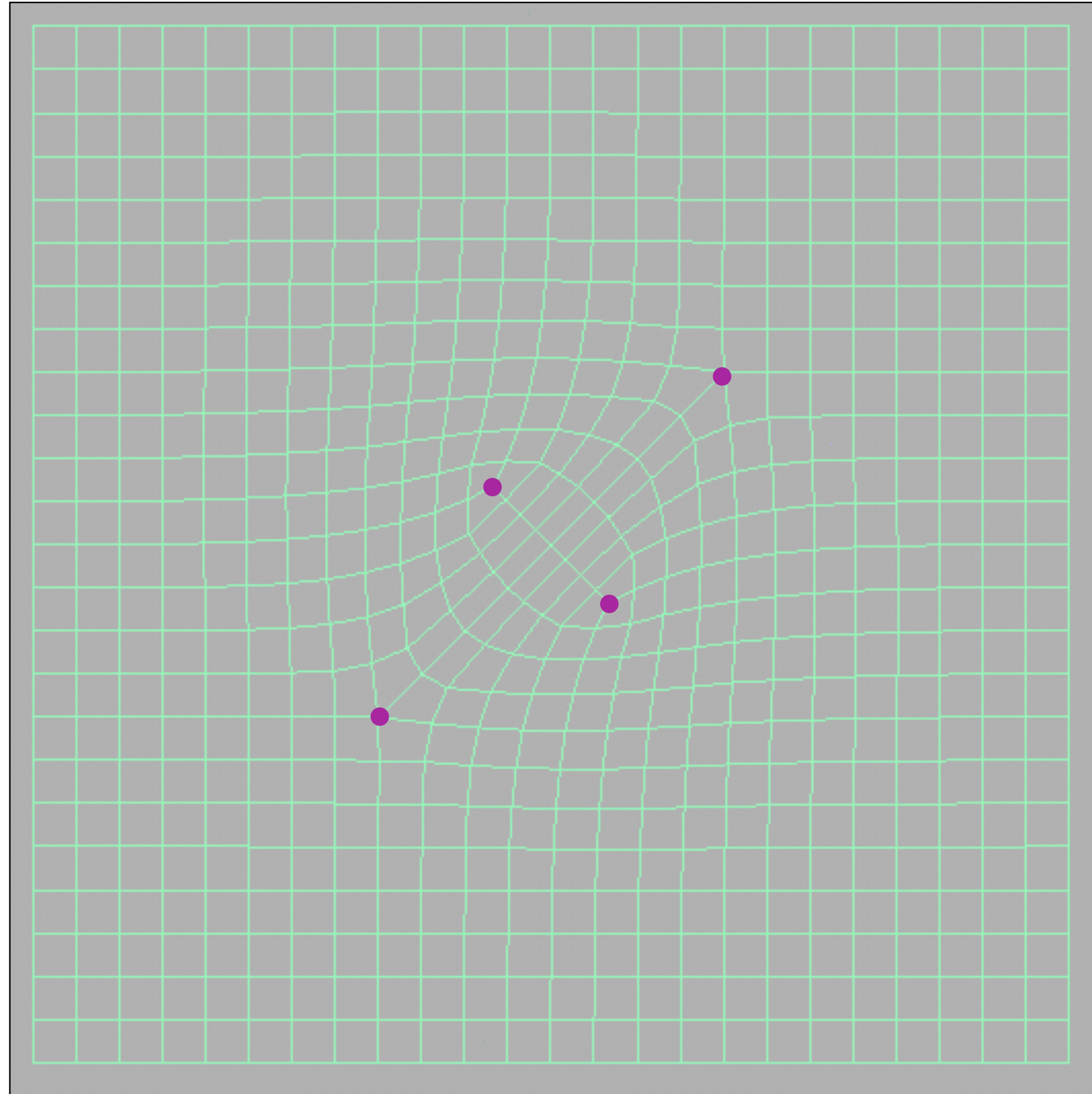
How many extraordinary vertices after first subdivision?  
What are their valences?  
How many non-quad faces?



# Catmull-Clark Subdivision (General Mesh)



# Catmull-Clark Subdivision (General Mesh)



# Catmull-Clark Vertex Update Rules (General Mesh)

$f$  = average of surrounding vertices

$$e = \frac{f_1 + f_2 + v_1 + v_2}{4}$$

$$v = \frac{\bar{f}}{n} + \frac{2\bar{m}}{n} + \frac{p(n-3)}{n}$$

**These rules reduce to earlier quad rules for ordinary vertices / faces**

$\bar{m}$  = average of adjacent midpoints

$\bar{f}$  = average of adjacent face points

$n$  = valence of vertex

$p$  = old "vertex" point

# Continuity of Catmull-Clark Surface

At extraordinary points

- Surface is at least  $C^1$  continuous

Everywhere else ("ordinary" regions)

- Surface is  $C^2$  continuous



# What About Sharp Creases?



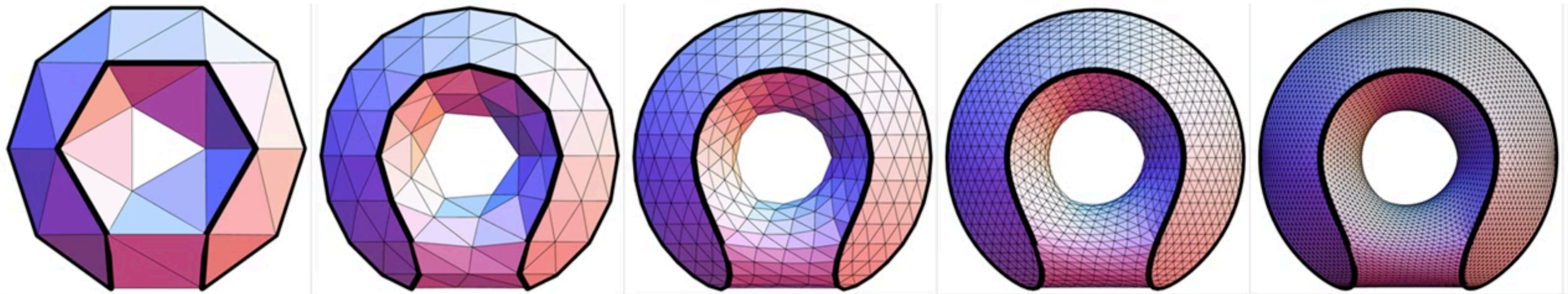
**From Pixar Short, "Geri's Game"**

**Hand is modeled as a Catmull Clark surface with creases between skin and fingernail**



# What About Sharp Creases?

Loop with Sharp Creases



Catmull-Clark with Sharp Creases

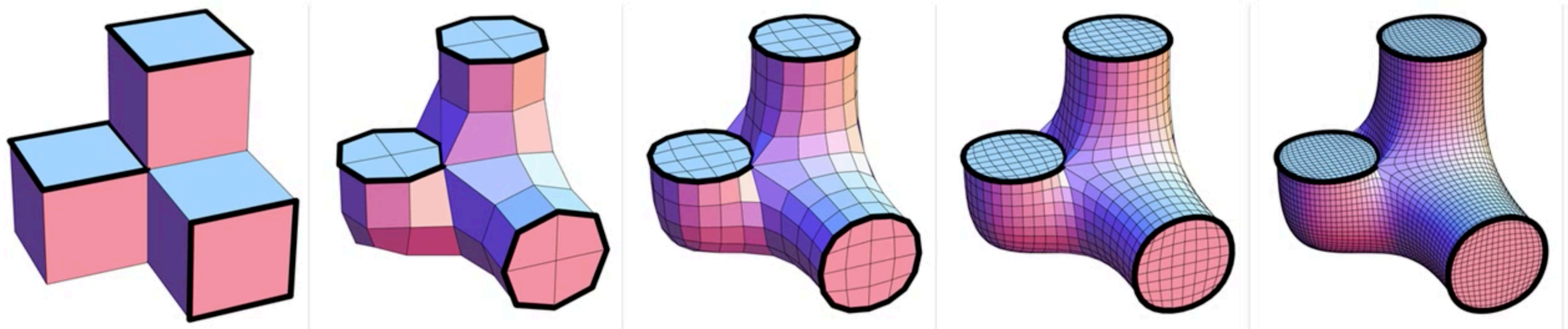
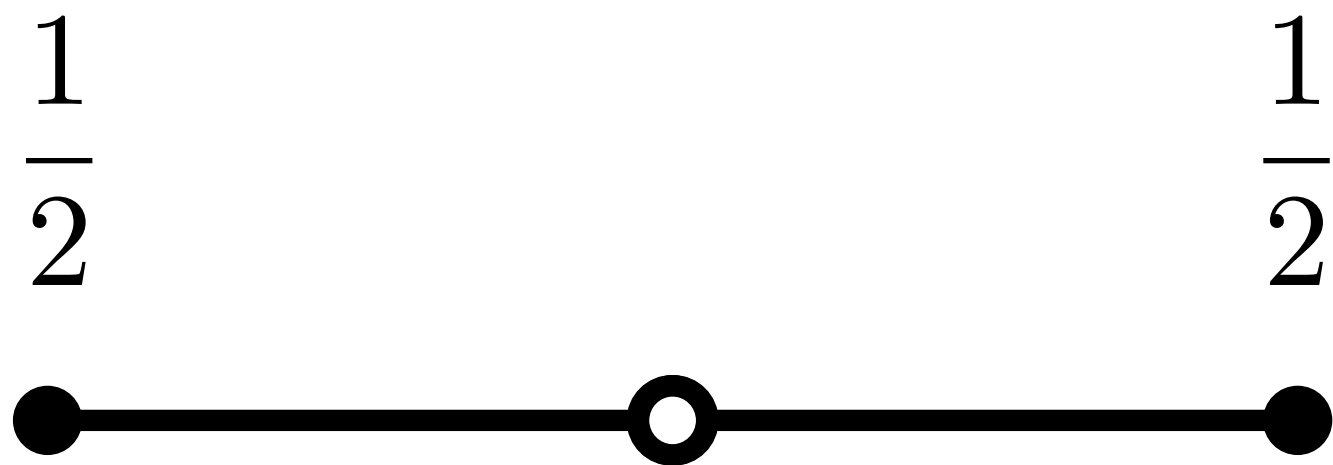


Figure from: Hakenberg et al. Volume Enclosed by Subdivision Surfaces with Sharp Creases

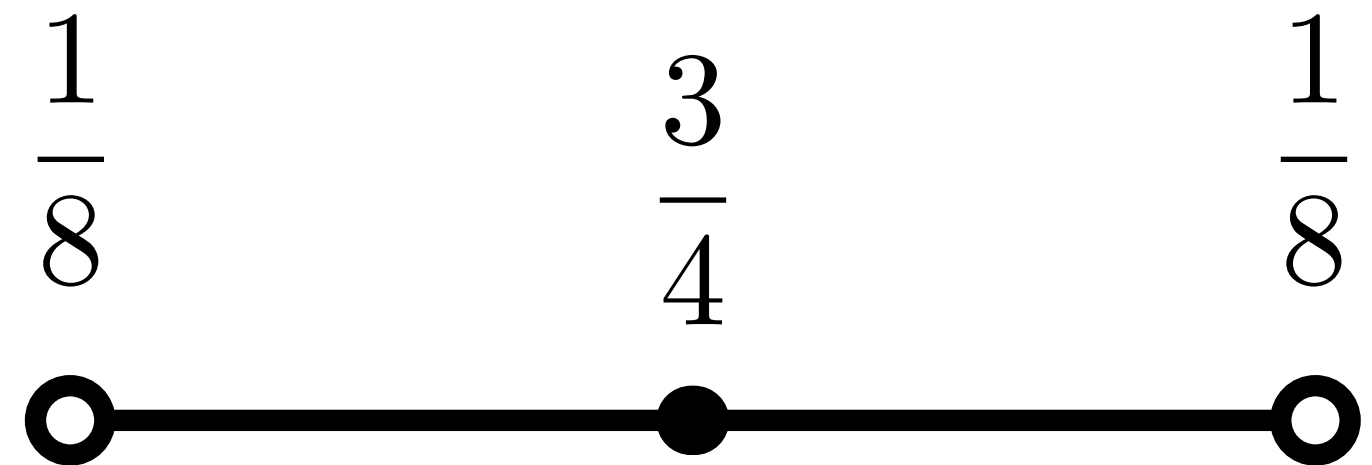
# Creases + Boundaries

Can create creases in subdivision surfaces by marking certain edges as "sharp". Boundary edges can be handled the same way

- Use different subdivision rules for vertices along these "sharp" edges



Insert new midpoint vertex,  
weights as shown



Update existing vertices,  
weights as shown



# Subdivision in Action (“Geri’s Game”, Pixar)

Subdivision used for entire character:

- Hands and head
- Clothing, tie, shoes





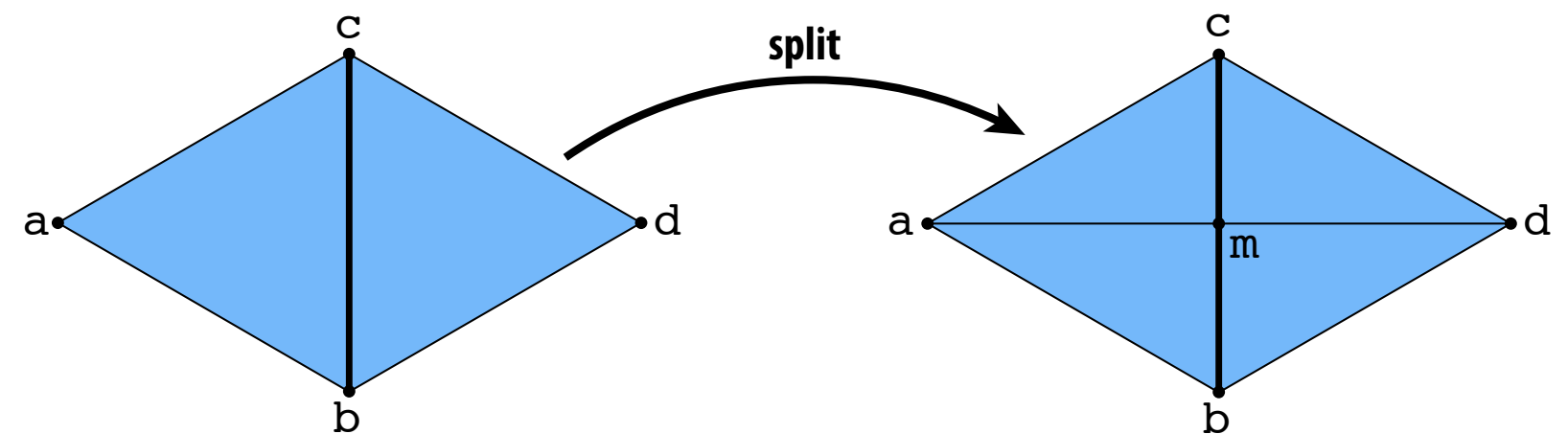
# Subdivision in Action (Pixar's "Geri's Game")



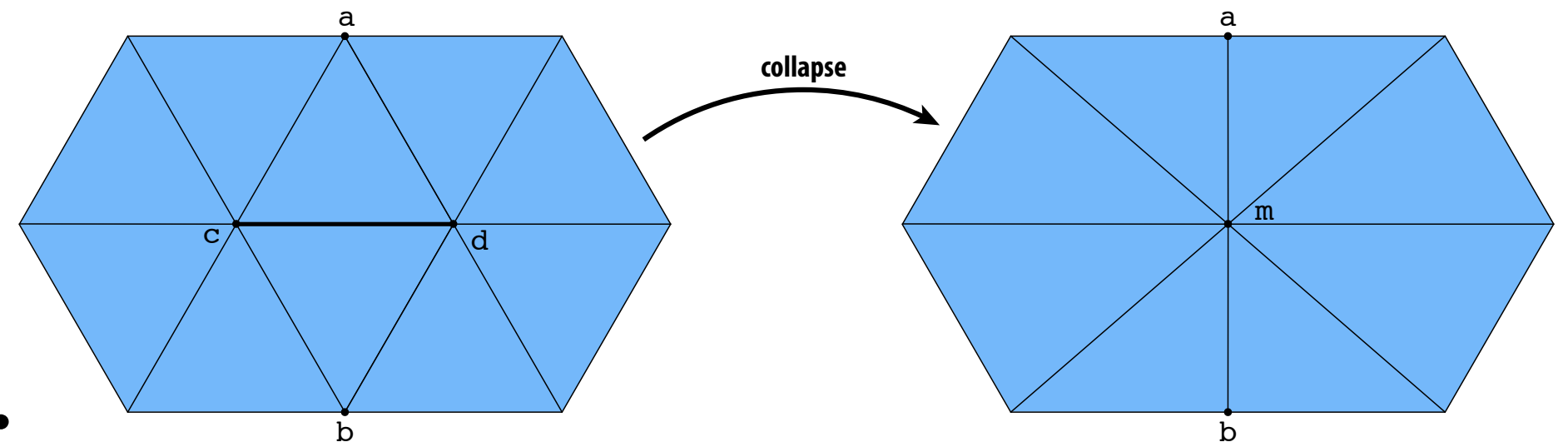
# **Mesh Simplification**

# How Do We Resample Meshes? (Reminder)

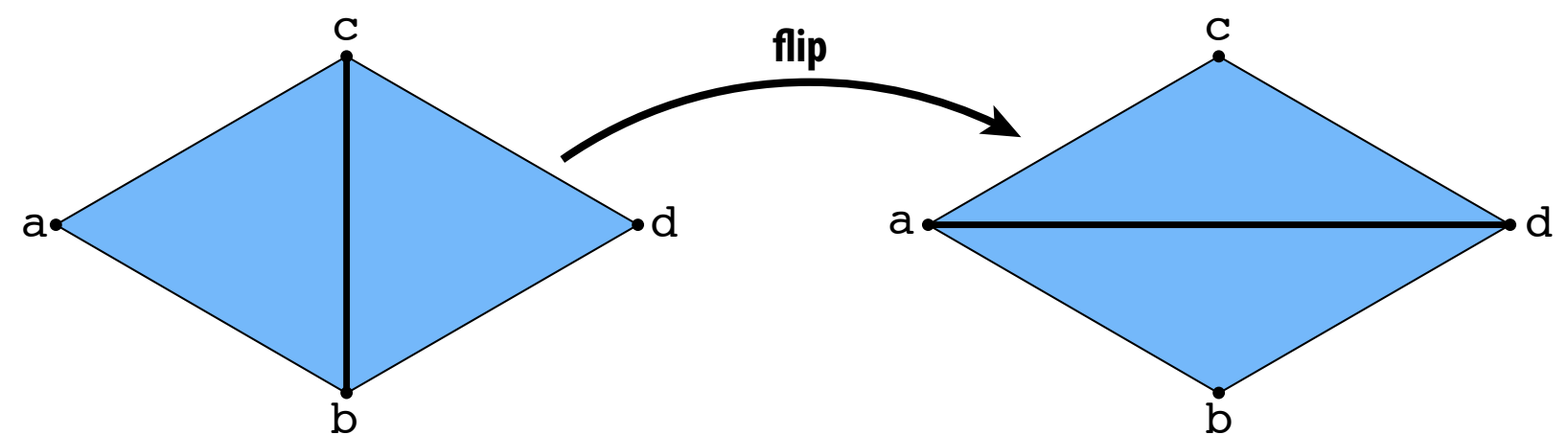
Edge split is (local) upsampling:



Edge collapse is (local) downsampling:



Edge flip is (local) resampling:



Still need to intelligently decide which edges to modify!



# Mesh Simplification

Goal: reduce number of mesh elements while maintaining overall shape



30,000 triangles



3,000



300



30

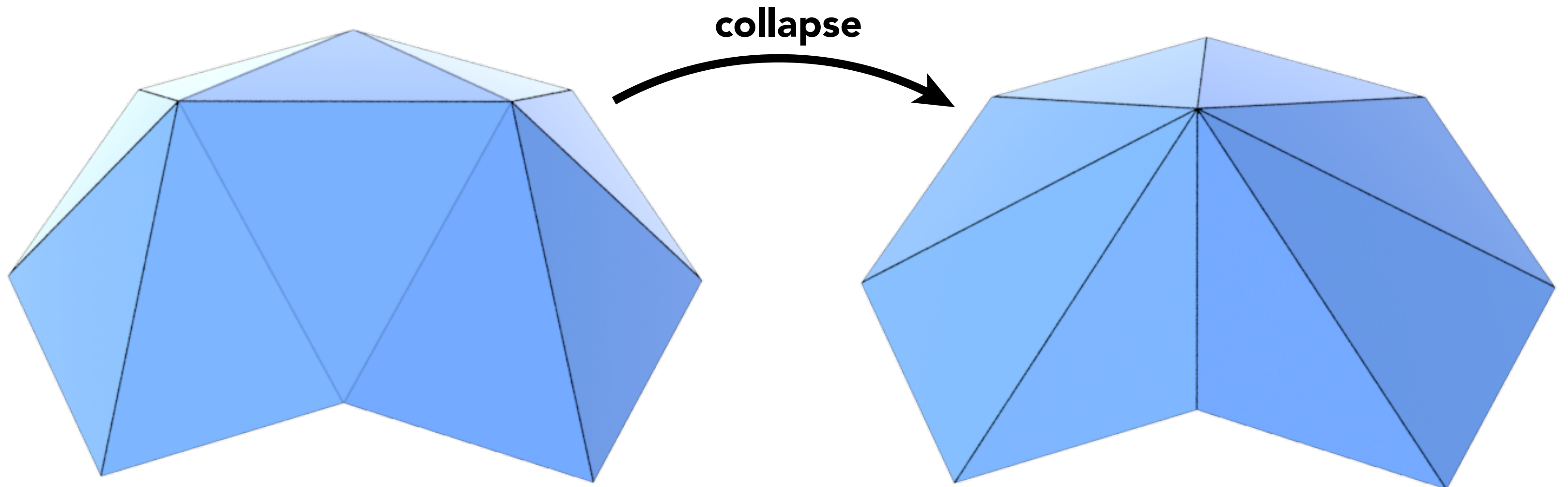


How to compute?



# Estimate: Error Introduced by Collapsing An Edge?

- How much geometric error for collapsing an edge?



# **Sketch of Quadric Error Mesh Simplification**

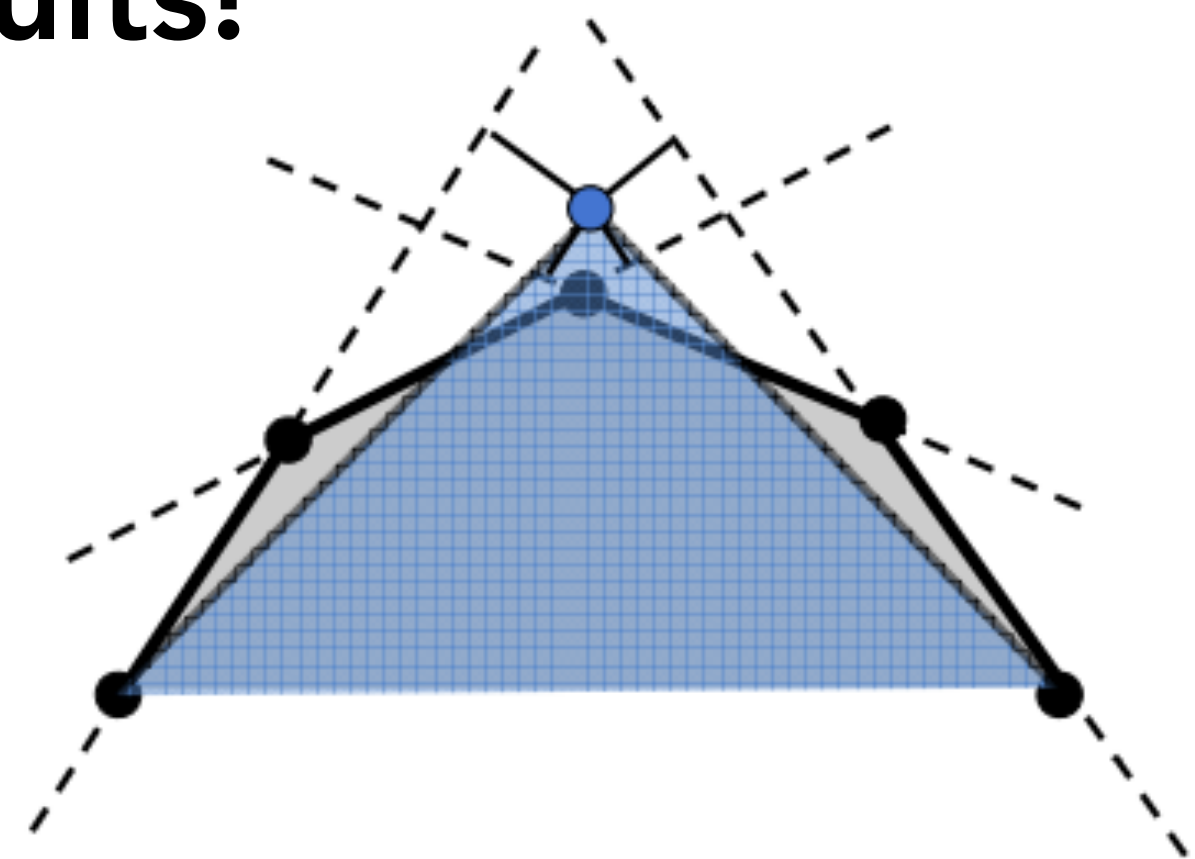
# Simplification via Quadric Error

Iteratively collapse edges

Which edges? Assign score with quadric error metric\*

- approximate distance to surface as sum of distances to planes containing triangles
- iteratively collapse edge with smallest score
- greedy algorithm... great results!

\* (Garland & Heckbert 1997)



# Quadric Error Matrix

Key idea:

- 4x4 ("quadric") symmetric matrix encodes distance to plane

For plane  $ax + by + cz + d = 0$

- Distance of query point  $(x, y, z)$  from plane is given by  $u^T Q u$ :
- $u := (x, y, z, 1)^T$  is the query point in homogeneous coordinates
- And  $Q$  is a symmetric matrix as follows:

$$Q = \begin{bmatrix} a^2 & ab & ac & ad \\ ab & b^2 & bc & bd \\ ac & bc & c^2 & cd \\ ad & bd & cd & d^2 \end{bmatrix}$$

- $Q$  contains 10 unique coefficients (small storage)



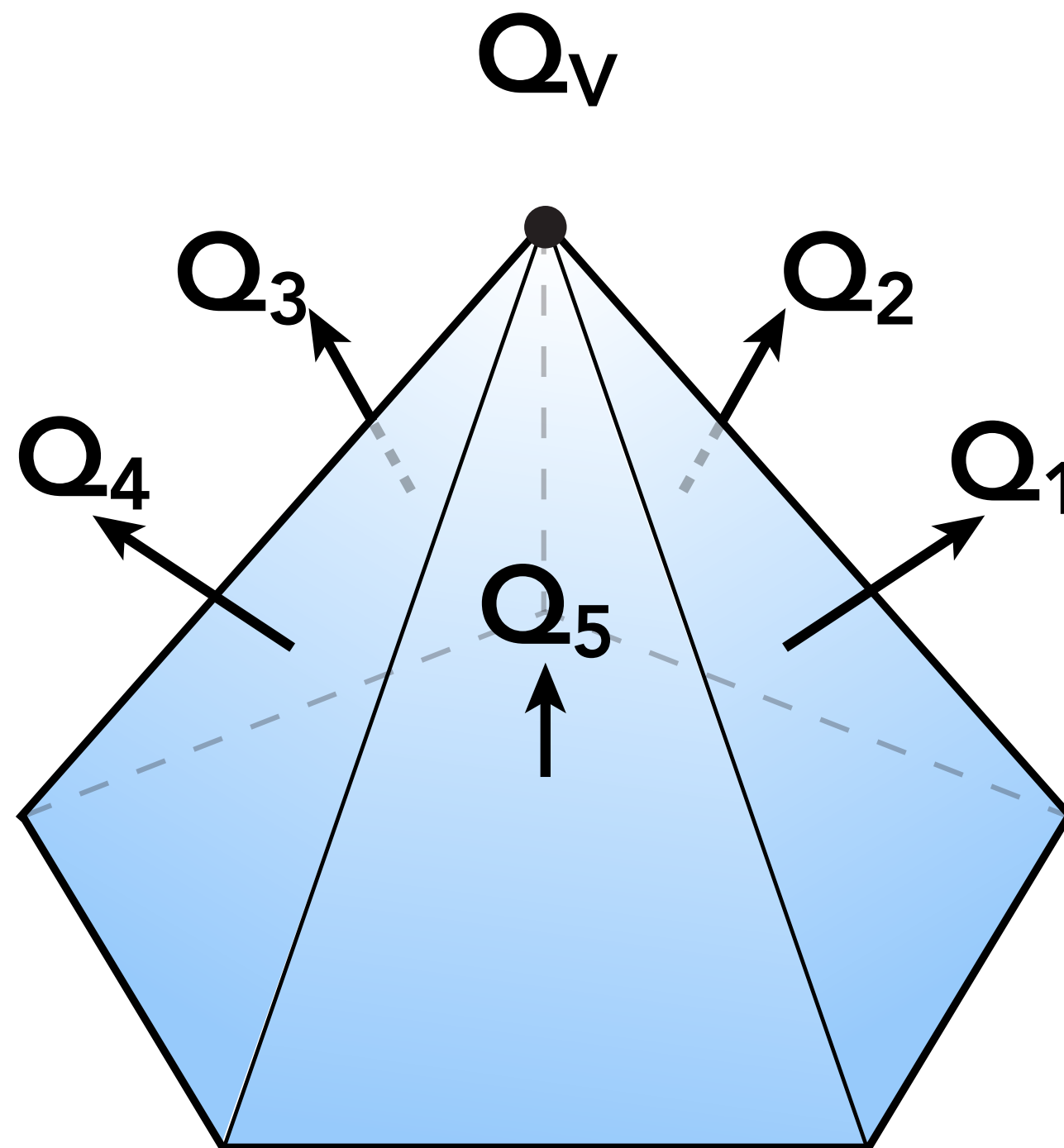
# Quadric Error Matrix: Derivation

- Suppose in coordinates we have
    - a query point  $(x,y,z)$
    - a normal  $(a,b,c)$
    - an offset  $d := -(x_p, y_p, z_p) \cdot (a, b, c)$
- $$Q = \begin{bmatrix} a^2 & ab & ac & ad \\ ab & b^2 & bc & bd \\ ac & bc & c^2 & cd \\ ad & bd & cd & d^2 \end{bmatrix}$$
- Then in homogeneous coordinates, let
    - $u := (x,y,z,1)$
    - $v := (a,b,c,d)$
  - Signed distance to plane is then  
 $D = uv^T = vu^T = ax+by+cz+d$
  - Squared distance is  $D^2 = (uv^T)(vu^T) = u (v^T v) u^T := u^T Q u$

# Quadric Error At Vertex

Approximate distance to vertex's triangles as sum of distances to each triangle's plane.

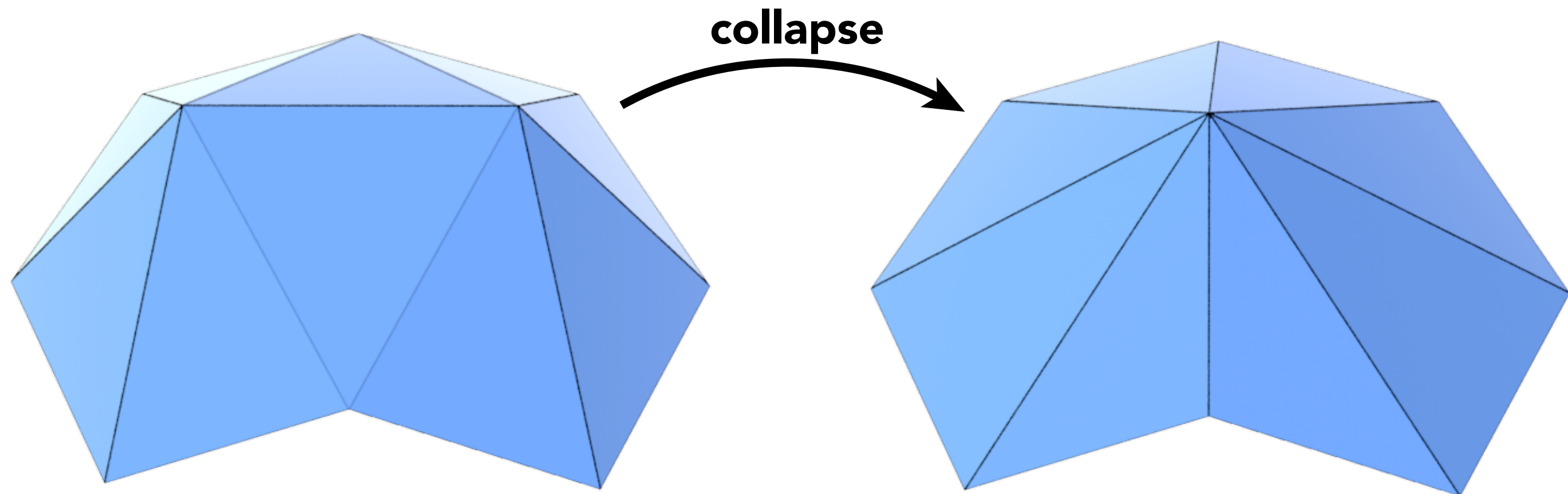
Encode this as a single quadric matrix for the vertex that is the sum of quadric error matrices for all triangles



$$Q_V = \sum_{i=1}^N Q_i$$

# Quadric Error of Edge Collapse

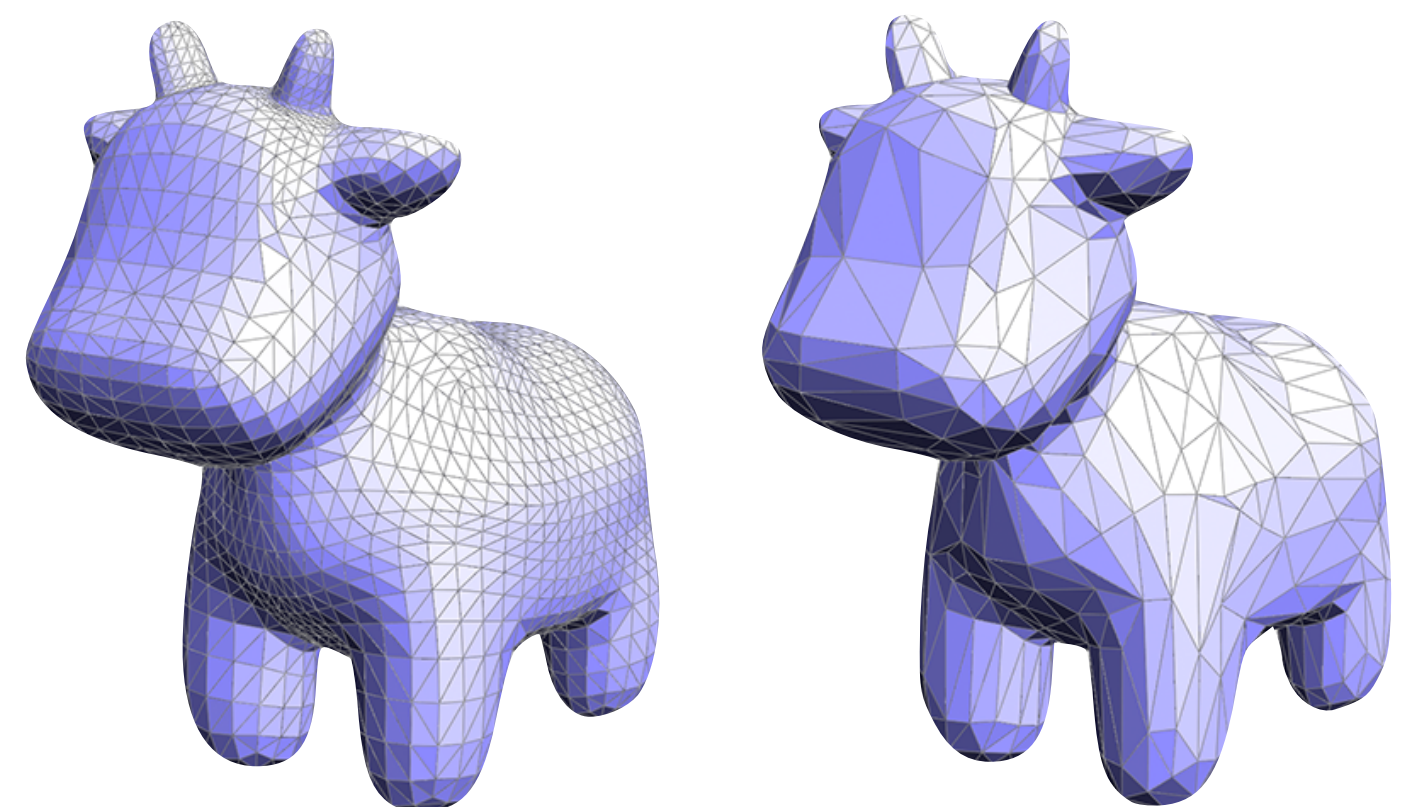
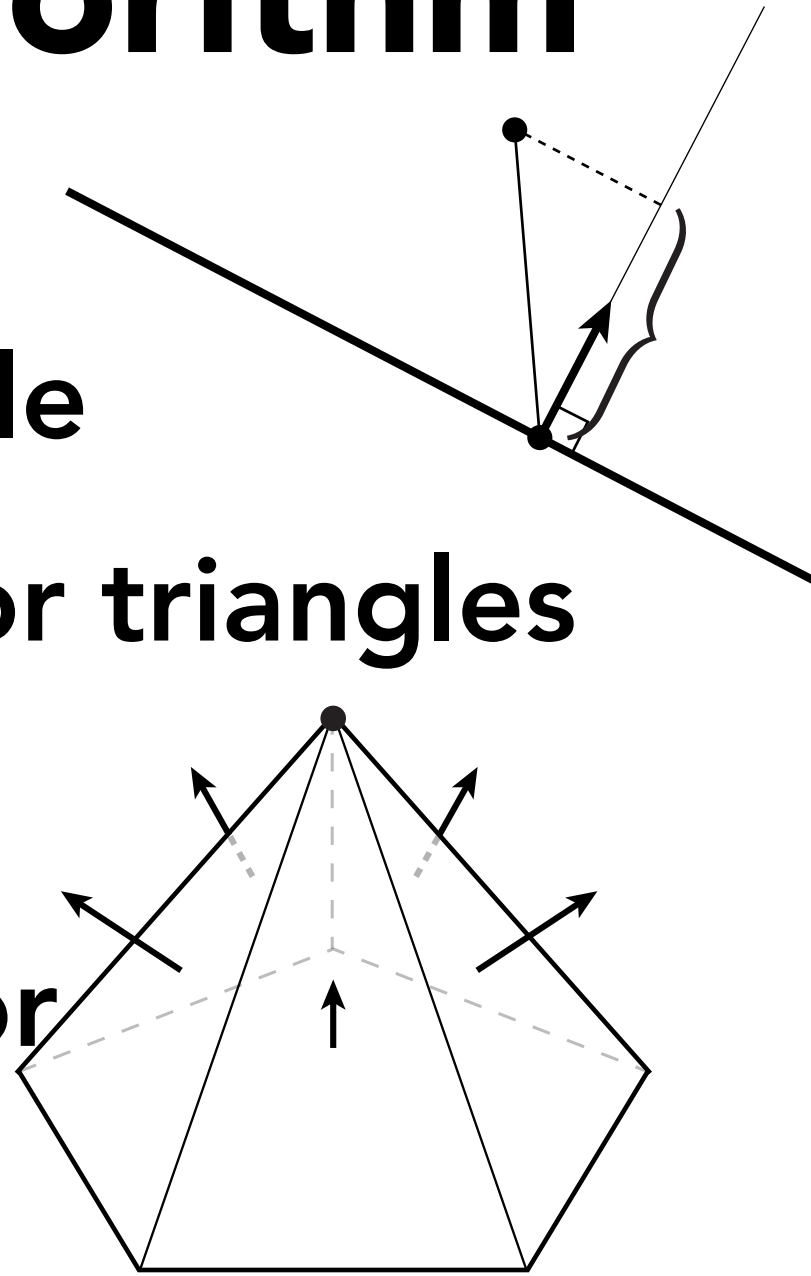
- How much does it cost to collapse an edge?
- Idea: compute edge midpoint, measure quadric error



- Better idea: choose point that minimizes quadric error
- More details: Garland & Heckbert 1997.

# Quadric Error Simplification: Algorithm

- Compute quadric error matrix  $Q$  for each triangle
- Set  $Q$  at each vertex to sum of  $Q$ s from neighbor triangles
- Set  $Q$  at each edge to sum of  $Q$ s at endpoints
- Find point at each edge minimizing quadric error
- Until we reach target # of triangles:
  - collapse edge  $(i,j)$  with smallest cost to get new vertex  $m$
  - add  $Q_i$  and  $Q_j$  to get quadric  $Q_m$  at vertex  $m$
  - update cost of edges touching vertex  $m$





# Quadric Error Mesh Simplification



5,804

994

532

248

64

Garland and Heckbert '97



30,000 triangles

3,000

300

30

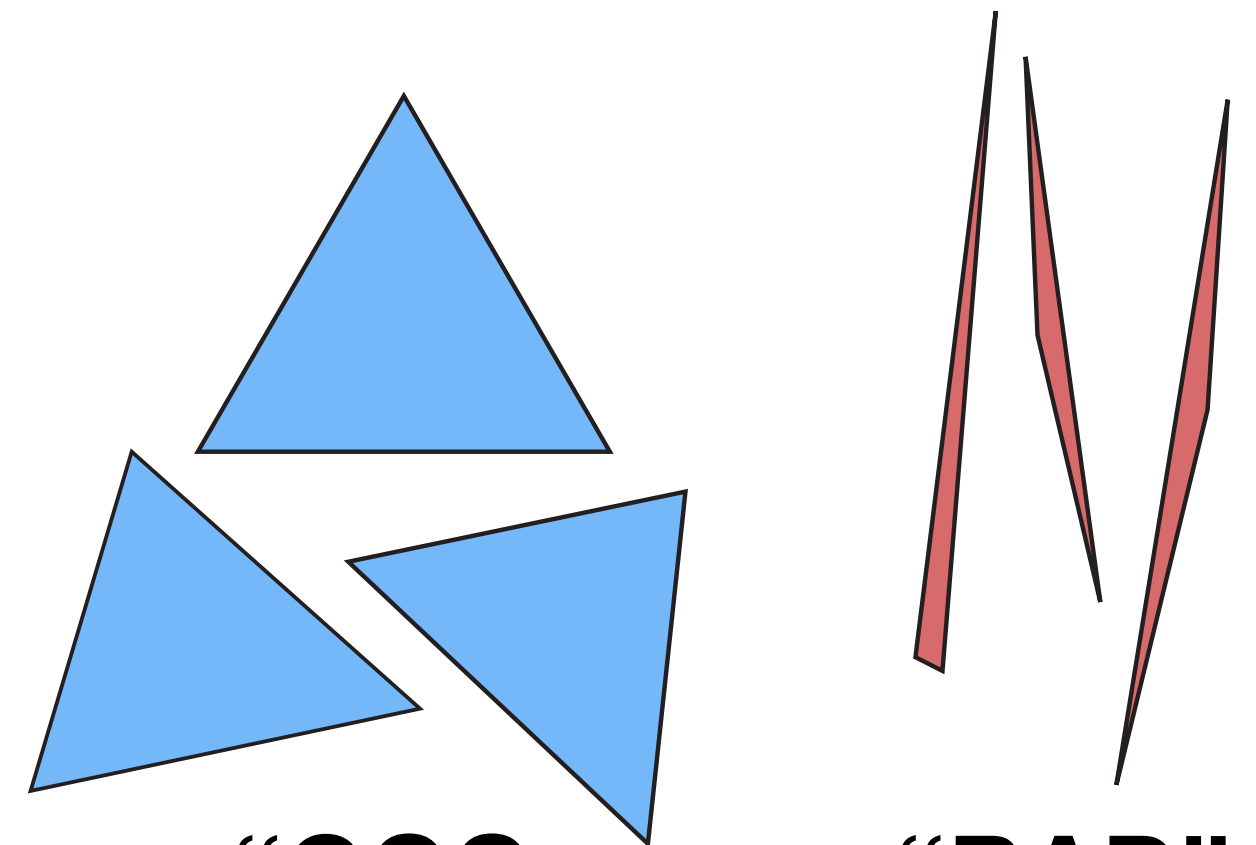
CS184/284A

Ng & Kanazawa

# **Mesh Regularization**

# What Makes a "Good" Triangle Mesh?

One rule of thumb: triangle shape

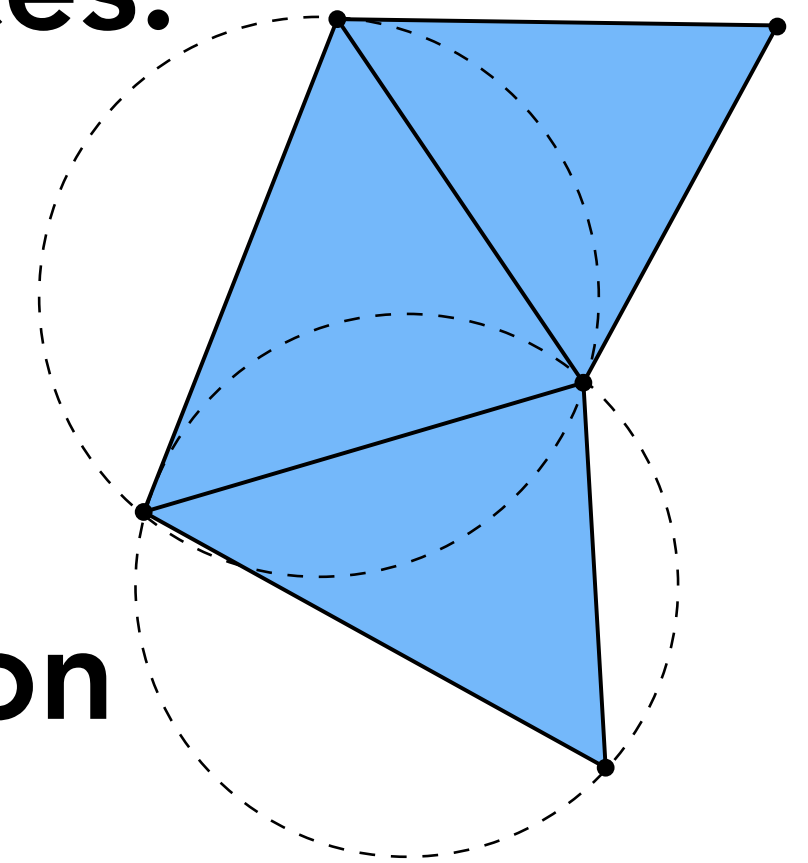


More specific condition: Delaunay

- "Circumcircle interiors contain no vertices."

Not always a good condition, but often\*

- Good for simulation
- Not always best for shape approximation

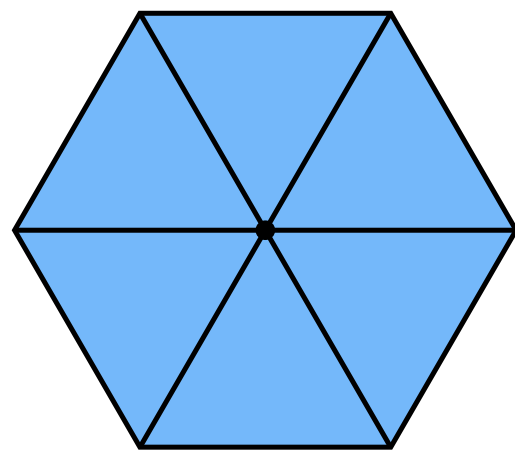


\*See Shewchuk, "What is a Good Linear Element"

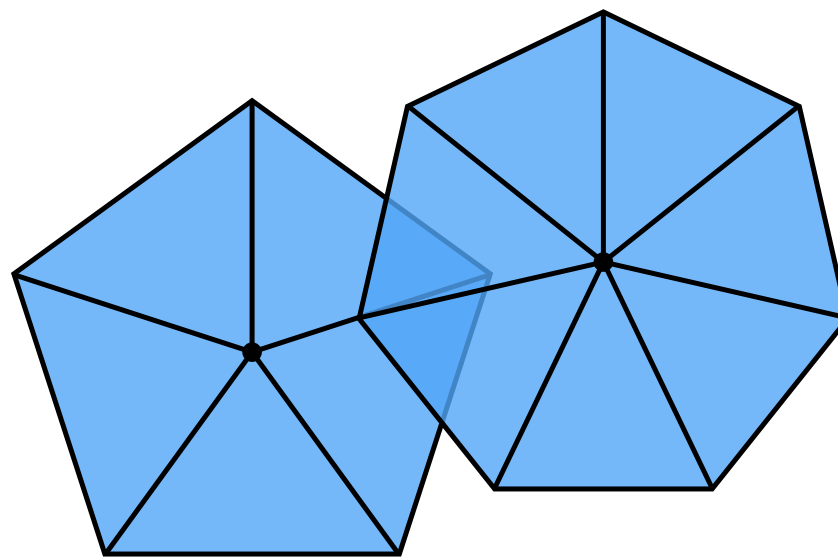
# What Else Constitutes a Good Mesh?

Rule of thumb: regular vertex degree

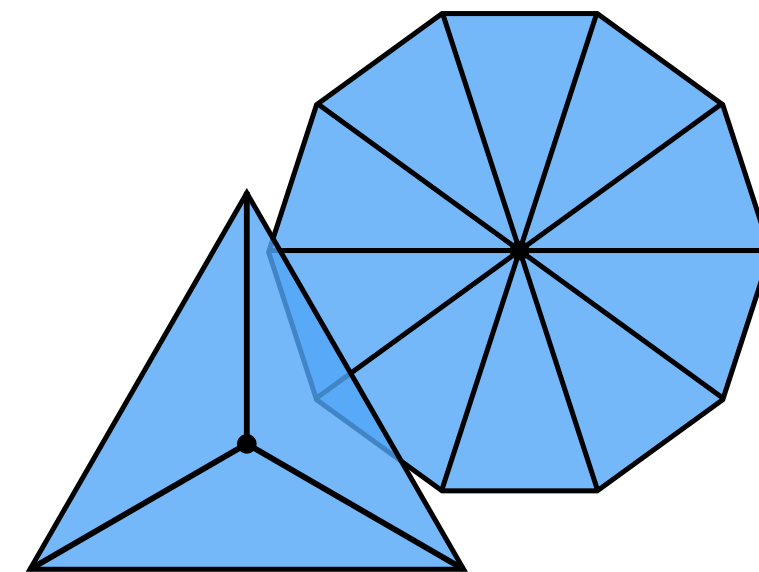
Triangle meshes: ideal is every vertex with valence 6:



“GOOD”

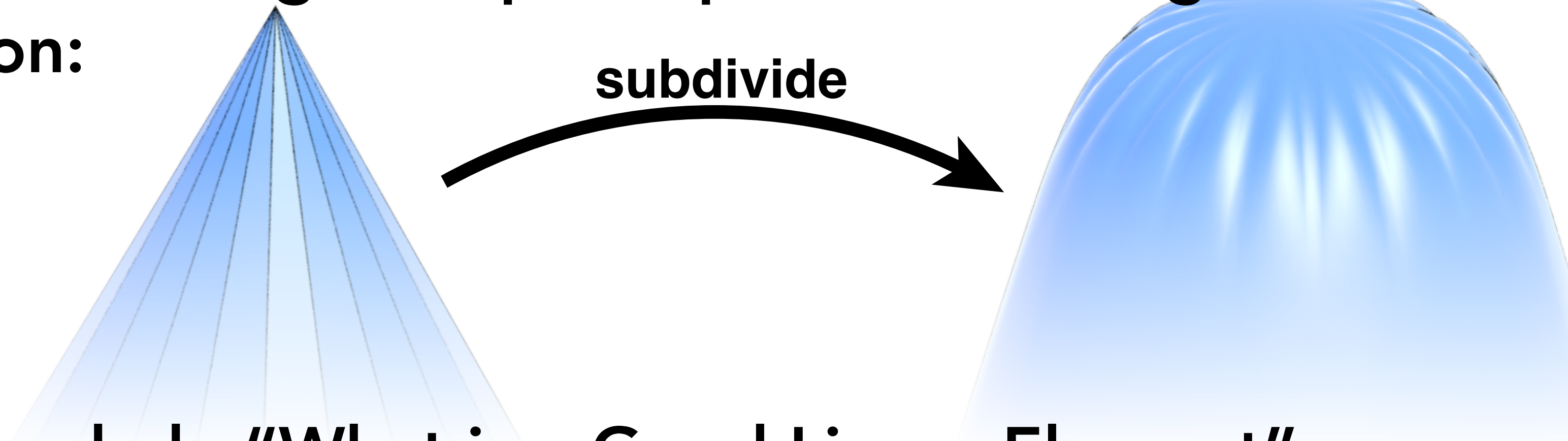


“OK”



“BAD”

Why? Better triangle shape, important for (e.g.) subdivision:

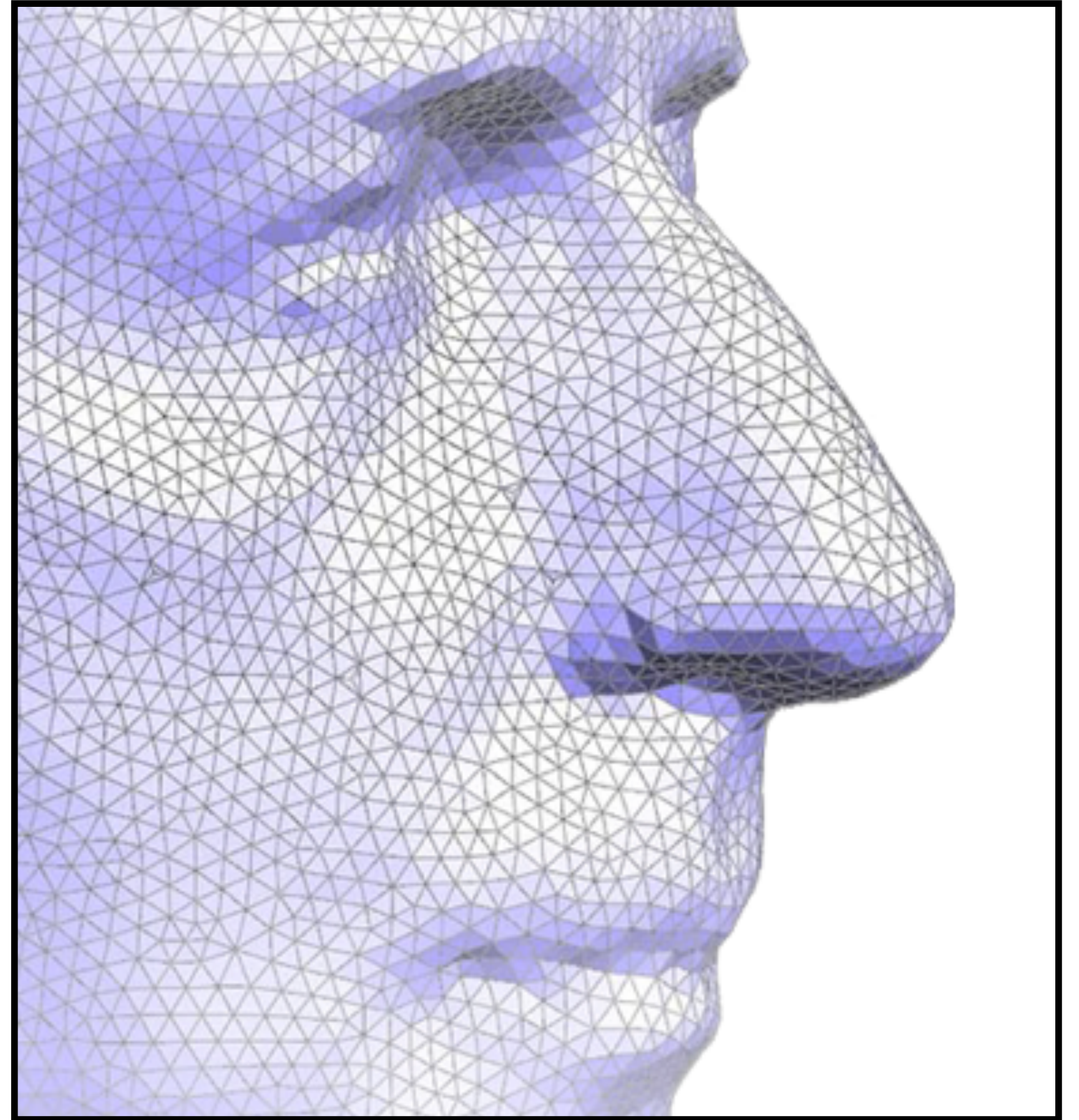
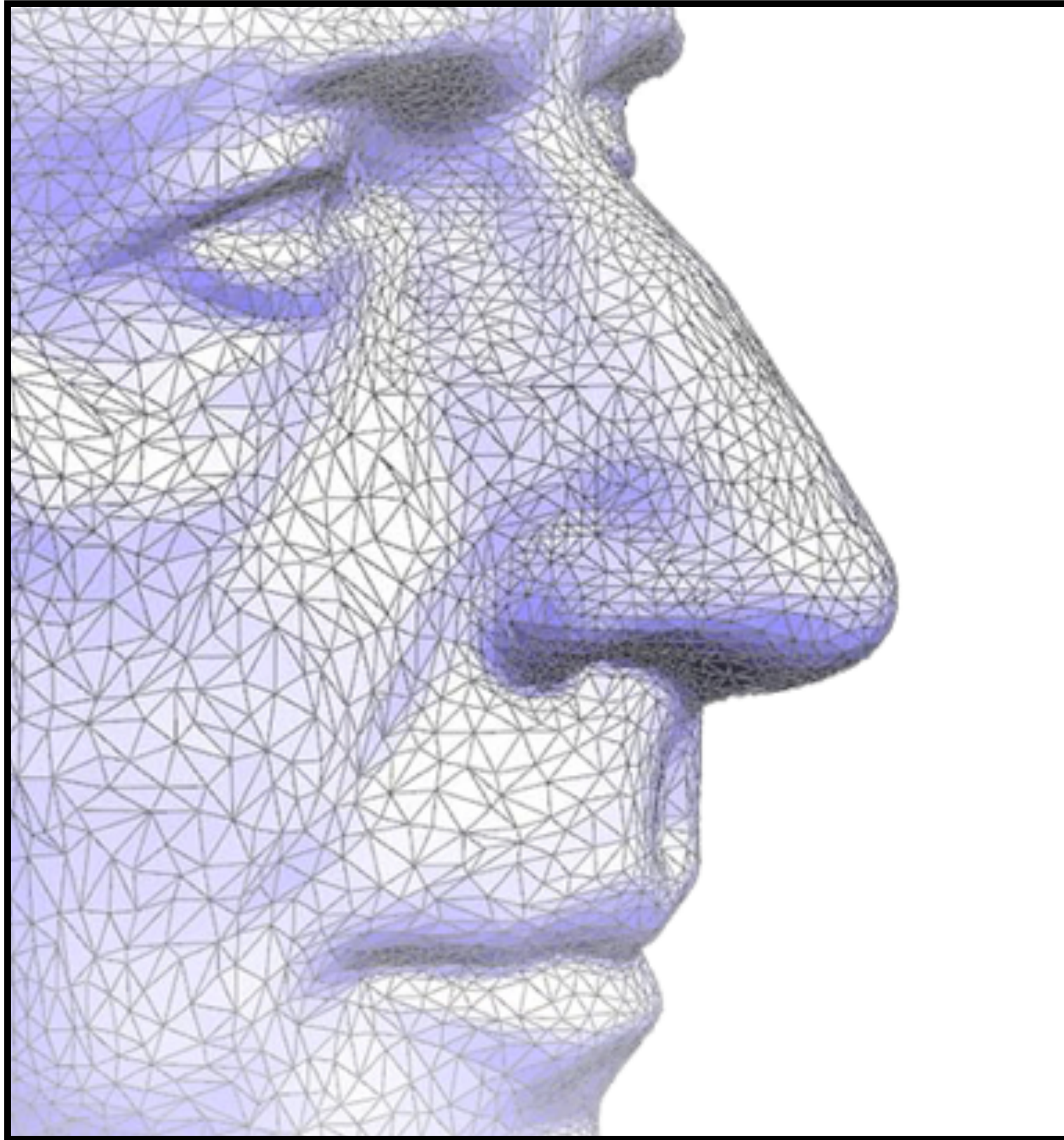


\*See Shewchuk, “What is a Good Linear Element”



# Isotropic Remeshing

Try to make triangles uniform in shape and size

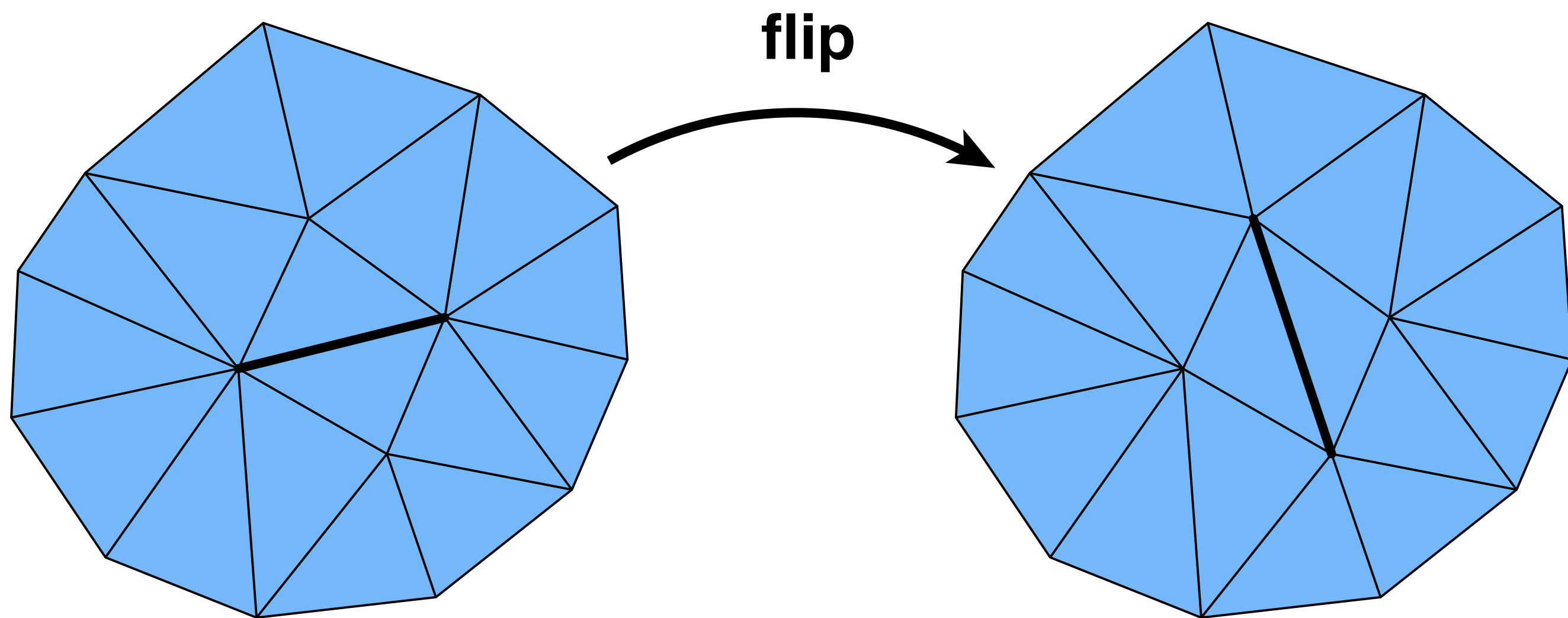




# How Do We Improve Degree?

Edge flips!

If total deviation from degree 6 gets smaller, flip it!



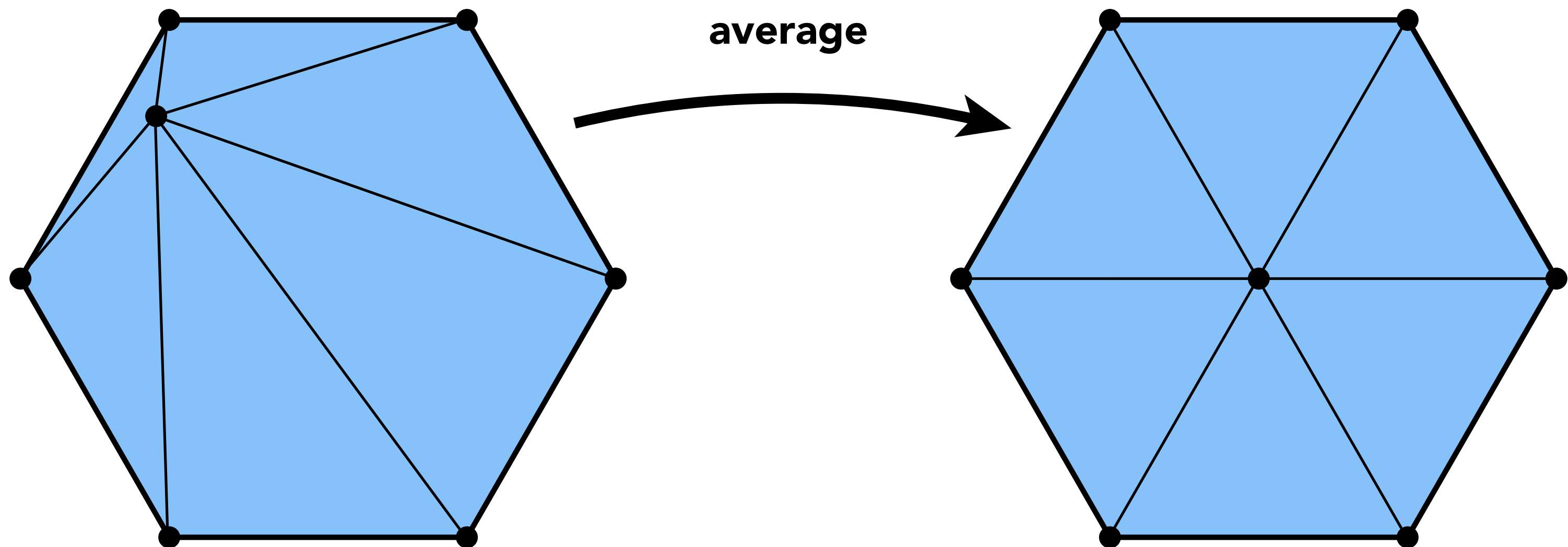
Iterative edge flipping acts like "discrete diffusion" of degree

No (known) guarantees; works well in practice

# How Do We Make Triangles “More Round”?

Delaunay doesn't mean equilateral triangles

Can often improve shape by centering vertices:



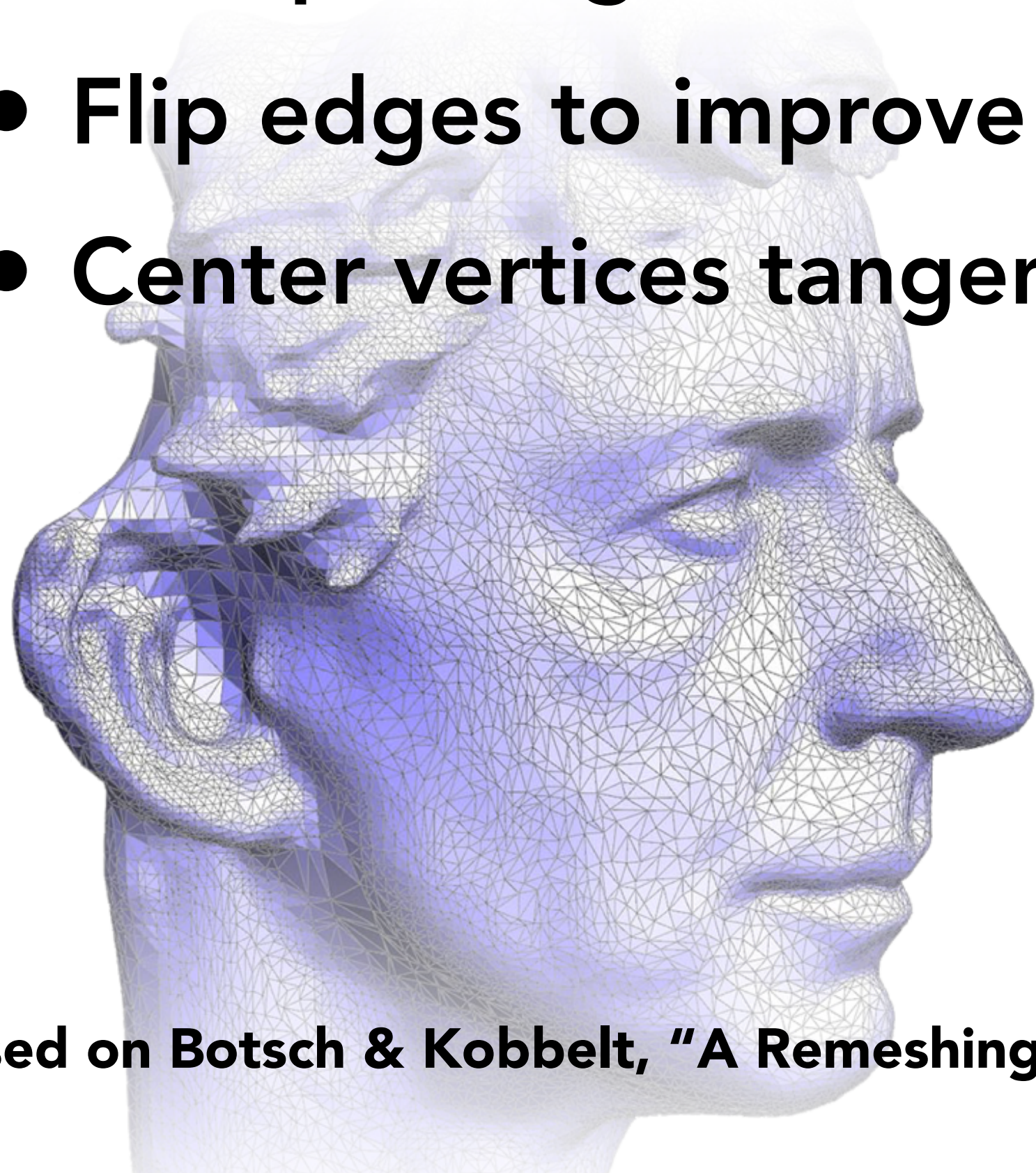
[Crane, "Digital Geometry Processing with Discrete Exterior Calculus"]



# Isotropic Remeshing Algorithm\*

Repeat four steps:

- Split edges over  $\frac{4}{3}$  mean edge length
- Collapse edges less than  $\frac{4}{5}$  mean edge length
- Flip edges to improve vertex degree
- Center vertices tangentially



\*Based on Botsch & Kobbelt, "A Remeshing Approach to Multiresolution Modeling"



# **MeshLab Demo**

# Things to Remember

## Triangle mesh representations

- Triangles vs points+triangles
- Half-edge structure for mesh traversal and editing

## Geometry processing basics

- Local operations: flip, split, and collapse edges
- Upsampling by subdivision (Loop, Catmull-Clark)
- Downsampling by simplification (Quadric error)
- Regularization by isotropic remeshing

# Acknowledgments

Thanks to Keenan Crane, Pat Hanrahan, Steve Marschner and James O'Brien for presentation resources.