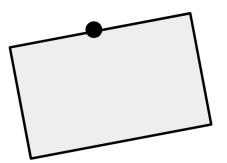
CS 184: FOUNDATIONS OF COMPUTER GRAPHICS

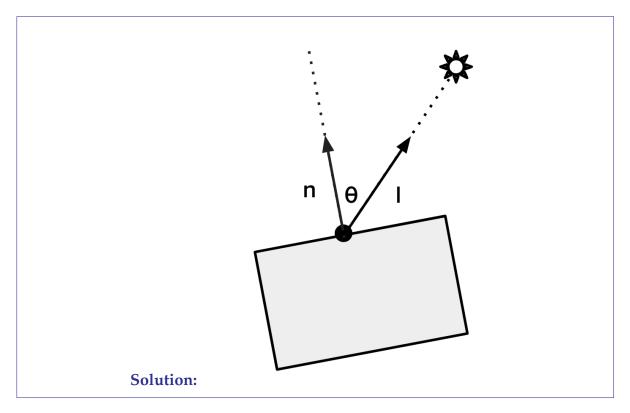
1 Graphics Pipeline — Lightning Round!

1. Name and describe the three terms in the Blinn-Phong Reflection Model.

Solution: Ambient, Diffuse, Specular lighting.

2. A light source shines on a tilted surface. Draw the light direction vector, **1**, and the normal vector, **n**, at the given point.





3. What is the light per unit area on this surface proportional to, according to Lambert's cosine law?

Solution: $\mathbf{l} \cdot \mathbf{n} = \cos \theta$. Light intensity reflected off surface is proportional to cosine of the angle between the light direction vector and surface normal vector.

4. Complete the following implementation of the Z-Buffer Algorithm in C++.

```
const int WIDTH = 800; // Width of framebuffer
const int HEIGHT = 600; // Height of framebuffer
struct Color {
    float r, g, b;
};
struct Sample {
    int x, y;
    float z;
    Color color;
};
struct Triangle {
    std::vector<Sample> samples;
};
```

```
void zBufferAlgorithm(const _____
                                                  _____ triangles,
                                                    _____ framebuffer,
                                                   _____ zbuffer) {
   for (const Triangle& T : triangles) {
       for (const Sample& sample : T.samples) {
          int x = sample.x;
          int y = sample.y;
          float z = sample.z;
          if (x >= 0 && x < WIDTH && y >= 0 && y < HEIGHT) {
              ____) {
                 framebuffer[x][y] = sample.color;
                                                    _;
              }
          }
      }
  }
}
```

Solution:

This solution uses **std::vector**, but other C++ implementations of lists could also work.

std::vector<Triangle>&
std::vector<std::vector<Color>>&

- std::vector<std::vector<float>>&
- z < zbuffer[x][y]
- z < zbuffer[x][y]
 zbuffer[x][y] = z
- 5. Prior to running this algorithm, what should the Z-buffer values be initialized to?

Solution:	
<pre>std::numeric_limits<float>::infinity()</float></pre>	
(Infinity.)	

Our goal is to fit a polynomial to given points and derivatives of a curve. We can solve this problem by formulating it as a system of linear equations in the coefficients of the polynomial.

1. List all degree 2 polynomials satisfying: f(0) = 1, f(1) = 2, f(2) = 5.

Solution: Let $f(t) = at^2 + bt + c$ be the degree 2 polynomial. The system of constraints we get is:

$$c = 1$$
$$a + b + c = 2$$
$$4a + 2b + c = 5$$

Solving the system yields a unique solution a = 1, b = 0, c = 1, so $f(t) = t^2 + 1$ is the only degree 2 polynomial.

2. How many degree 3 polynomials satisfy the given constraints?

Solution: There are infinitely many such polynomials. Intuitively, three constraints on four coefficients (since a degree 3 polynomial has four coefficients) leave one free parameter, so there must be infinitely many degree 3 solutions.

Detailed explanation (optional). We know from the first question that $f(t) = t^2 + 1$ is the unique degree 2 polynomial satisfying

$$f(0) = 1, \quad f(1) = 2, \quad f(2) = 5.$$

Now let f(t) be a degree 3 polynomial. We can write

$$f(t) = \underbrace{(t^2 + 1)}_{f_1(t)} + \underbrace{(at^3 + bt^2 + ct + d)}_{f_2(t)},$$

where $f_1(t)$ already satisfies the constraints. Therefore, the polynomial

$$f_2(t) := at^3 + bt^2 + ct + d$$

must vanish at t = 0, 1, 2:

$$f_2(0) = 0, \quad f_2(1) = 0, \quad f_2(2) = 0.$$

From $f_2(0) = 0$, we get d = 0. From $f_2(1) = 0$, a + b + c = 0. From $f_2(2) = 0$, 8a + 4b + 2c = 0.

To solve this system, subtract $\left(\frac{1}{2}\right) f_2(2)$ from $f_2(1)$

 $(4a+2b+c) - (a+b+c) = 0 \implies 3a+b = 0 \implies b = -3a.$

Next, substitute b = -3a into a + b + c = 0:

$$a - 3a + c = 0 \implies -2a + c = 0 \implies c = 2a.$$

Hence,

$$f_2(t) = a \left(t^3 - 3t^2 + 2t \right).$$

Therefore, *every* degree-3 polynomial satisfying the original constraints can be written as

$$f(t) = t^{2} + 1 + a(t^{3} - 3t^{2} + 2t),$$

for any real *a*. Since *a* is a free parameter that can take any real value, there are infinitely many solutions.

3. Suppose we have a list of constraints:

$$f(0) = p_0, f'(0) = d_0, f(1) = p_1, f'(1) = d_1, \dots, f(k) = p_k, f'(k) = d_k.$$

Since we have 2(k + 1) constraints (one function and one derivative condition per point), the unique interpolating polynomial must have degree 2k + 1.

For a function f, what are the tradeoffs when either

- solving for a single degree 2k + 1 polynomial, versus
- taking the point and derivative constraints at *i* and *i* − 1 for *i* = 1,..., *k* and using them to fit *k* cubic Hermite splines?

Solution: If we solve for a single high-degree polynomial, it will have infinitely many continuous derivatives. However, it may exhibit unpredictable behavior between the control points, and furthermore, changing a single constraint will affect the entire curve.

If we solve for k cubic Hermite splines, then the resulting curve will be continuous and have a continuous first derivative. The curve will have infinitely many continuous derivatives within each segment but may have a discontinuous second derivative at the control points. However, changing a single constraint will only affect two of the cubic splines: those that have the control point as an endpoint.

4. Consider a cubic polynomial $f(t) = at^3 + bt^2 + ct + d$ that satisfies the following

conditions:

$$f(0) = f_0,$$

$$f(1) = f_1,$$

$$f''(0) = f''_0,$$

$$f''(1) = f''_1.$$

Write the matrix that, when inverted and applied to the vector $(f_0, f_1, f_0'', f_1'')^T$, allows you to recover the coefficients a, b, c, and d of the polynomial.

Solution:

$\begin{pmatrix} a \\ b \end{pmatrix}$		$\begin{pmatrix} 0 \\ 1 \\ 0 \\ 6 \end{pmatrix}$	0	0	1	-1	(f_0)
b		1	1	1	1		$\begin{pmatrix} f_0 \\ f_1 \\ f_0'' \\ f_1'' \end{pmatrix}$
$\begin{pmatrix} c \\ c \\ d \end{pmatrix}$	=	0	2	0	0		f_0''
d		$\setminus 6$	2	0	0/		f_1''

5. Given the numerical inverse of the matrix:

$$\begin{pmatrix} a \\ b \\ c \\ d \end{pmatrix} = \begin{pmatrix} 0 & 0 & -1/6 & 1/6 \\ 0 & 0 & 1/2 & 0 \\ -1 & 1 & -1/3 & -1/6 \\ 1 & 0 & 0 & 0 \end{pmatrix} \begin{pmatrix} f_0 \\ f_1 \\ f_0'' \\ f_1'' \end{pmatrix}$$

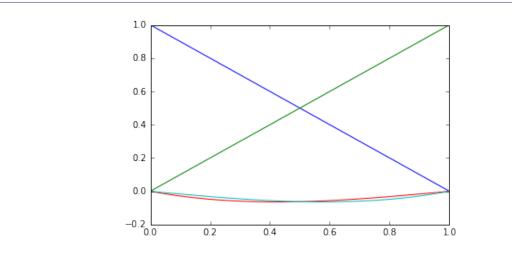
what are the **basis polynomials** that allow expressing f(t) in terms of f_0, f_1, f''_0, f''_1 ?

Solution:

The basis polynomials for this problem are:

$$G_0(t) = -t + 1, \quad G_1(t) = t, \quad G_2(t) = -\frac{t^3}{6} + \frac{t^2}{2} - \frac{t}{3}, \quad G_3(t) = \frac{t^3}{6} - \frac{t}{6}.$$

Here's a plot of these functions:



These basis polynomials form a set of functions that allow us to express any cubic polynomial satisfying the given constraints. The interpolating function is given by:

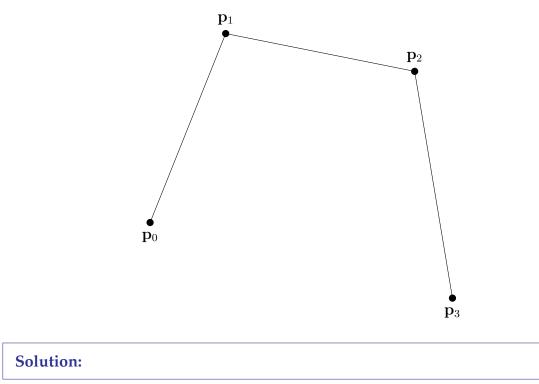
$$f(t) = f_0 G_0(t) + f_1 G_1(t) + f_0'' G_2(t) + f_1'' G_3(t).$$

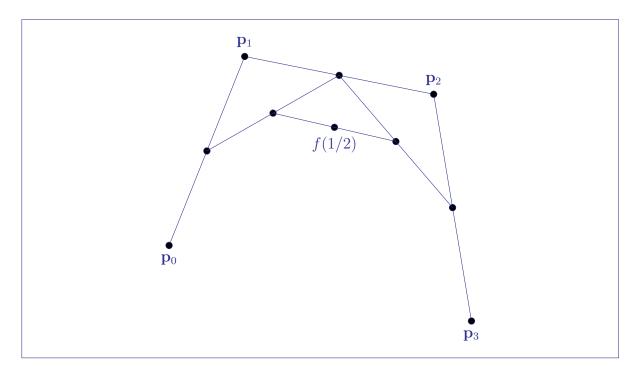
f(t) matches the given function values and second derivatives at t = 0 and t = 1.

Given k + 1 points $\mathbf{p}_0, \dots, \mathbf{p}_k$, create a new set of k points $\mathbf{p}'_0, \dots, \mathbf{p}'_{k-1}$ by computing $\mathbf{p}'_i = \operatorname{lerp}(\mathbf{p}_i, \mathbf{p}_{i+1}, t)$, where $\operatorname{lerp}(\mathbf{p}_i, \mathbf{p}_{i+1}, t) = (1 - t)\mathbf{p}_i + t\mathbf{p}_{i+1}$.

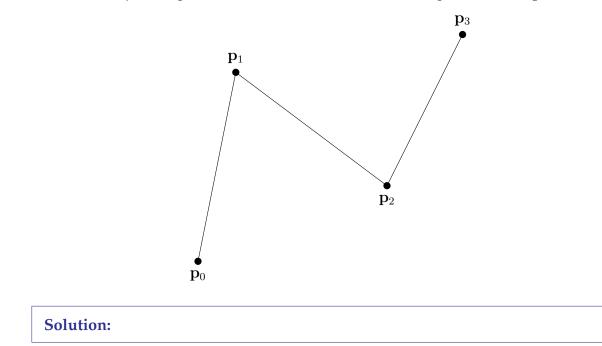
Perform k times to yield a single point, f(t).

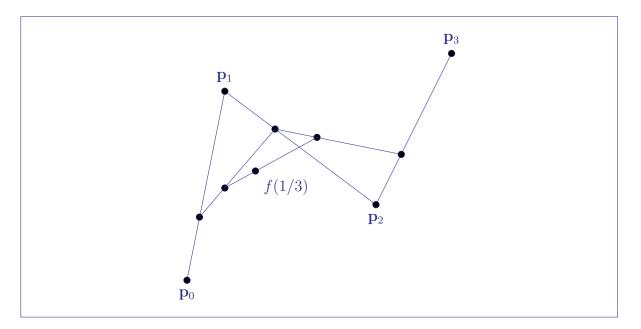
1. Use de Casteljau's algorithm to construct f(1/2) with the given control points.





2. Use de Casteljau's algorithm to construct f(1/3) with the given control points.





3. Show that the point with parameter *t* on the Bézier curve with control points \mathbf{p}_0 , \mathbf{p}_1 , \mathbf{p}_2 , \mathbf{p}_3 is given by $s^3\mathbf{p}_0 + 3s^2t\mathbf{p}_1 + 3st^2\mathbf{p}_2 + t^3\mathbf{p}_3$, where s = 1 - t. (Hint: apply de Casteljau's algorithm algebraically to the control points. With this setup, linear interpolation between two points \mathbf{q}_0 and \mathbf{q}_1 looks like $s\mathbf{q}_0 + t\mathbf{q}_1$.)

Solution:

Level 0: $\mathbf{p}_0, \mathbf{p}_1, \mathbf{p}_2, \mathbf{p}_3$ Level 1: $\mathbf{b}_0^1(t) = s\mathbf{p}_0 + t\mathbf{p}_1$ $\mathbf{b}_1^1(t) = s\mathbf{p}_1 + t\mathbf{p}_2$ $\mathbf{b}_2^1(t) = s\mathbf{p}_2 + t\mathbf{p}_3$ Level 2: $\mathbf{b}_0^2(t) = s(s\mathbf{p}_0 + t\mathbf{p}_1) + t(s\mathbf{p}_1 + t\mathbf{p}_2) = s^2\mathbf{p}_0 + 2st\mathbf{p}_1 + t^2\mathbf{p}_2$ $\mathbf{b}_1^2(t) = s(s\mathbf{p}_1 + t\mathbf{p}_2) + t(s\mathbf{p}_2 + t\mathbf{p}_3) = s^2\mathbf{p}_1 + 2st\mathbf{p}_2 + t^2\mathbf{p}_3$ Level 3: $\mathbf{b}_0^3(t) = s(s^2\mathbf{p}_0 + 2st\mathbf{p}_1 + t^2\mathbf{p}_2) + t(s^2\mathbf{p}_1 + 2st\mathbf{p}_2 + t^2\mathbf{p}_3)$ $= s^3\mathbf{p}_0 + 3s^2t\mathbf{p}_1 + 3st^2\mathbf{p}_2 + t^3\mathbf{p}_3$

4. What is this matrix product? (Hint: *don't* expand it. Instead, think about what each matrix in the product does. How are they related to de Casteljau's algorithm?)

$$\begin{pmatrix} s & t \end{pmatrix} \begin{pmatrix} s & t & 0 \\ 0 & s & t \end{pmatrix} \begin{pmatrix} s & t & 0 & 0 \\ 0 & s & t & 0 \\ 0 & 0 & s & t \end{pmatrix}$$

Solution: Without doing any additional computation, the matrix product must be $(s^3 \ 3s^2t \ 3st^2 \ t^3)$.

The rest of the solution assumes that s = 1 - t, but the same logic applies to any *s* (in Question 3, we made no assumptions on *s*). Let's consider the result of applying this matrix product to a set of points,

$$\begin{pmatrix} \mathbf{p}_0^T \\ \mathbf{p}_1^T \\ \mathbf{p}_2^T \\ \mathbf{p}_3^T \end{pmatrix}.$$

Starting from the right, the first matrix maps this input to the three points that result in applying one iteration of de Casteljau's algorithm,

$$\begin{pmatrix} s\mathbf{p}_0^T + t\mathbf{p}_1^T \\ s\mathbf{p}_1^T + t\mathbf{p}_2^T \\ s\mathbf{p}_2^T + t\mathbf{p}_3^T \end{pmatrix} = \begin{pmatrix} \mathbf{b}_0^1(s,t)^T \\ \mathbf{b}_1^1(s,t)^T \\ \mathbf{b}_2^1(s,t)^T \end{pmatrix}.$$

The other two matrices perform the remaining 2 iterations of the algorithm. As we saw in Question 3, applying this matrix product is equivalent to the map

$$\begin{pmatrix} \mathbf{p}_0^T \\ \mathbf{p}_1^T \\ \mathbf{p}_2^T \\ \mathbf{p}_3^T \end{pmatrix} \mapsto s^3 \mathbf{p}_0^T + 3s^2 t \mathbf{p}_1^T + 3st^2 \mathbf{p}_2^T + t^3 \mathbf{p}_3^T.$$

Hence, the matrix product must be $\begin{pmatrix} s^3 & 3s^2t & 3st^2 & t^3 \end{pmatrix}$.

Alternatively, you could just do the computation... but why would you do that?

$$(s \ t) \begin{pmatrix} s \ t \ 0 \\ 0 \ s \ t \end{pmatrix} \begin{pmatrix} s \ t \ 0 \\ 0 \ s \ t \end{pmatrix} \begin{pmatrix} s \ t \ 0 \ 0 \\ 0 \ s \ t \ 0 \\ 0 \ 0 \ s \ t \end{pmatrix}$$
$$= (s^2 \ 2st \ t^2) \begin{pmatrix} s \ t \ 0 \ 0 \\ 0 \ s \ t \ 0 \\ 0 \ 0 \ s \ t \end{pmatrix}$$
$$= (s^3 \ 3s^2t \ 3st^2 \ t^3)$$

5. For a Bézier curve defined by 3 control points, what is the degree of the polynomial that results from de Casteljau's algorithm? What about for *k* points?

Solution: Three points yields a quadratic polynomial, and k points yields a polynomial of degree k - 1.¹

In Question 3, we saw that

$$\mathbf{b}_0^3(t) = \binom{3}{0}(1-t)^3\mathbf{p}_0 + \binom{3}{1}(1-t)^2t\mathbf{p}_1 + \binom{3}{2}(1-t)t^2\mathbf{p}_2 + \binom{3}{3}t^3\mathbf{p}_3.$$

More generally, we can define Bézier curves in Bernstein form, i.e.

$$\mathbf{b}_0^n(t) = \sum_{j=0}^n \mathbf{p}_j \binom{n}{i} (1-t)^{n-i} t^i.$$

This is a polynomial with degree n. If we have 3 points, then n = 2, and the polynomial is quadratic in t. If we have k points, then n = k - 1, and the polynomial has degree k.

¹Here, we consider the control points \mathbf{p}_j to be variables and the resulting function to be a polynomial in *t*. If you can fix \mathbf{p}_j 's, it is possible to obtain a polynomial with lower degree. For example, if every $\mathbf{p}_j = \mathbf{0}$, then the resulting polynomial is $\mathbf{0}$.